# Structural Variant Detection 

A Novel Approach
by

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# A Dissertation Presented in Partial Fulfillment of the Requirements for the Degree <br> Doctor of Philosophy 

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#### Abstract

Genomic structural variation (SV) is defined as gross alterations in the genome broadly classified as insertions/duplications, deletions inversions and translocations. DNA sequencing ushered structural variant discovery beyond laboratory detection techniques to high resolution informatics approaches. Bioinformatics tools for computational discovery of SVs however are still missing variants in the complex cancer genome. This study aimed to define genomic context leading to tool failure and design novel algorithm addressing this context. Methods: The study tested the widely held but unproven hypothesis that tools fail to detect variants which lie in repeat regions. Publicly available 1000-Genomes dataset with experimentally validated variants was tested with SVDetect-tool for presence of true positives (TP) SVs versus false negative (FN) SVs, expecting that FNs would be overrepresented in repeat regions. Further, the novel algorithm designed to informatically capture the biological etiology of translocations (non-allelic homologous recombination and 3-D placement of chromosomes in cells-context) was tested using simulated dataset. Translocations were created in known translocation hotspots and the novel-algorithm tool compared with SVDetect and BreakDancer. Results: 53\% of false negative (FN) deletions were within repeat structure compared to 81\% true positive (TP) deletions. Similarly, 33\% FN insertions versus $42 \%$ TP, $26 \%$ FN duplication versus $57 \%$ TP and $54 \%$ FN novel sequences versus $62 \%$ TP were within repeats. Repeat structure was not driving the tool's inability to detect variants and could not be used as context. The novel algorithm with a redefined context, when tested against SVDetect and BreakDancer was able to detect $10 / 10$ simulated translocations with 30 X coverage dataset and $100 \%$ allele frequency, while SVDetect captured $4 / 10$ and BreakDancer detected 6/10. For 15X coverage dataset with $100 \%$ allele frequency, novel algorithm was able to detect all


ten translocations albeit with fewer reads supporting the same. BreakDancer detected $4 / 10$ and SVDetect detected 2/10 Conclusion: This study showed that presence of repetitive elements in general within a structural variant did not influence the tool's ability to capture it. This context-based algorithm proved better than current tools even with half the genome coverage than accepted protocol and provides an important first step for novel translocation discovery in cancer genome.

## DEDICATION

This work is dedicated to my family without whose unwavering solid support I would not have tested the limits of by ability.

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## 1. I NTRODUCTI ON

Identifying differences at the genome level between a diseased and healthy individual has been the cornerstone of current medical genetic research with the purpose of identifying and targeting disease causing variants.

### 1.1 Overview

Structural variations in the human genome are changes in the genome, when compared to the reference human genome, that lead to gross aberrations in the physical structure of the genome. These typically include insertions, deletions, inversions, and translocations. Structural variations are typically studied by comparing the target DNA to the consensus reference genome. Insertions represent the inclusion of a segment of DNA sequence inserted into the target genome. Similarly a deletion represents a deleted segment, duplication represents a duplicated or repeat segment, and an inversion is a change in the orientation of the segment. A translocation is a change in the physical location of a segment of DNA normally present on the reference chromosome to another chromosomal location in the donor. Structural variations play a very important role in cancer development. Translocations were the earliest identified variants in cancer with the identification of Philadelphia chromosome (translocation between chromosomes 9 and 22 [t(9;22)]) as a hallmark of chronic myeloid leukemia (CML) (Nowell \& Hungerford, 1960; Rowley, 1973). These gross variants in cancer produce fusion genes as BCR-ABL1, between a breakpoint cluster region on chromosome 22 and tyrosine kinase receptor gene chromosome 9 that is a recurrent phenomenon in CML and one of the diagnostic criteria for this disease. Another commonly occurring fusion gene between a transmembrane protease gene (TMPRSS2) and one of two transcription factors (ETV1 or ERG) was first reported in 23 of 29 prostate cancers (Tomlins et al., 2005).

These commonly occurring fusion genes are fast becoming the hallmark of cancer diagnostics with researchers now aiming to define the etiology of this recurrence.

### 1.2 Significance of the Problem

Detecting structural variations in the cancer genome is the first step towards extracting disease causing mutations. Current tools using well-examined bioinformatics approaches to variant detection are still missing variants in the cancer genome. The reasons for this could be due to allele frequency of the mutant genome captured in an experiment, due to tumor heterogeneity, mapping algorithm flaws or structural variant detection algorithm flaws. In addition, these tools have been designed without accounting for the very special case of cancer genome complexity. Understanding the biological processes leading to structural variation generation in the cancer genome and using this context to define the algorithm for variant detection is an overlooked novel approach explored in this study, specifically in the study of translocations.

Chromosomal translocations are the most obvious signature for many cancers and serve as a very important biomarker. Cancer genomes have inherent genetic instability either due to various micro-environmental factors leading to somatic mutations or inherent predisposition in the genome leading to germline mutations. Normal cellular machinery has checks to curb such major overhauls using DNA repair mechanism. The cancer genome however has circumvented this repair mechanism; most evident in the resulting fusion-genes which produce protein products which further disrupt the repair infrastructure.

DNA sequencing technologies aim to identify these variants at a faster rate and with more accuracy based on the sheer volume of data exiting the machines. Clinical translation of this information helps to design specific drug targets for these specific
translocation fusion gene products. Two major hurdles to achieve this purpose, namely the presence of the human reference genome and the reducing cost of sequencing the patient genome have now been surmounted. Bioinformatics challenges to map and align this large amount of data, though not fully conquered, have made significant progress. The challenge now facing the scientific community is how to make sense of the data by designing better algorithms for detecting variants.

Every approach available in the literature has used some form of statistical/mathematical modeling or computational algorithm to solve this issue of detecting structural variants without paying much attention to the biology driving these variants. Computational methods to detect structural variants (SV) utilize a consistent algorithm (Tuzun et al., 2005) including: 1) creating a distribution of the reads length (insert size) to derive a mean and standard deviation (s.d) 2) defining SV signatures (e.g. insertion = reads mapping 3 s.d. outside the mean length), and 3) clustering all reads which support the same SV. The current study used a different approach using biological domain information which is already known and well documented in literature to help design a more effective and biologically plausible algorithm to detect translocations in the cancer genome.

### 1.3 Theoretical basis for the study

Cancer occurs due to genomic instability that leads to disruption of normal cell functions. These mutations and structural alterations can be germline or somatic. Genomic mutations in germline cells lead to their transmission to the next generation as a cancer susceptibility gene like BRCA1 (P. Kent et al., 1995). Mutations in somatic cells can cause cancer due to abnormal proliferation of somatic cells with aberrant genomic structure. Identifying these gross structural changes has been done traditionally using experimental technique like karyotyping. Karyotyping uses stains (Giemsa stains) to color the chromosome during mitosis and studying the
banding patterns (G-banding) to define structural abnormalities like duplication of a region of chromosome or presence or absence of chromosomes (McNeil, Montagna, Difilippantonio, \& Ried, 2012). However the primary issue with karyotyping remains to be the resolution of the abnormal chromosomal region, which is $5-10$ million base pairs (Mbps). A successful experiment is also dependent on quality of cell division rates in metaphase, which is when these chromosomes are captured (Bridge \& Cushman-Vokoun, 2011).

One of the major advances in cytogenetic diagnosis was development of fluorescent in-situ hybridization (FISH) techniques and its various offshoots (Volpi \& Bridger, 2008). The principle was based on hybridization of complementary DNA to specific fluorescent probes, designed for specific regions on the genome, and viewed through special cameras for detecting structural abnormalities (McNeil et al., 2012). These technologies have far reaching applications from pre-natal detection of structural abnormalities (Hastings, Nisbet, Waters, Spencer, \& Chitty, 1999), profiling gene expression during meiosis in mammalian cells (Mahadevaiah, Costa, \& Turner, 2009) to identifying novel fusion genes in leukemias (H. Lee et al., 2013). Further these techniques have been seminal in identifying specific recurrent translocations in tumors leading to targeted therapy for fusion genes (Buchdunger et al., 1996; Druker et al., 1996; Inokuchi, 2006; Lynch et al., 2004; Mathews et al., 2010; Niu et al., 1999; Paez et al., 2004). However the resolution of these methods can be about 1 kbps depending on the size of the probe, the target locus and type of FISH being used. A subsequent development, also based on hybridization, known as comparative genomic hybridization (CGH) (Kallioniemi et al., 1992) uses tumor and normal control samples to compare copy number variation between normal and tumor tissue and even detect gene fusions (Przybytkowski, Ferrario, \& Basik, 2011). These methods have higher resolution of 500 bps to even 50 bps but can be
extremely sensitive to the quality of tumor samples and therefore have a high rate of errors (Fragouli et al., 2011). These methods are thus dependent on the quality of the probes and the genomic distance between the probes on the array. Furthermore, they can only detect aberrations that alter copy number; copy-number neutral events are missed.

With the advent of massively parallel high throughput sequencing technologies gaining ground due to reduction in cost, mutation definition is getting higher and higher resolution with the number of implicated oncogenes growing from 291 in 2004 (Futreal et al., 2004) to 384 in 2010 (Santarius, Shipley, Brewer, Stratton, \& Cooper, 2010) and the most recent database of cancer mutations, 'Mitelmans Database of Gene Chromosome Aberrations and Gene Fusions in cancer', a National Cancer Institute ( NCI ) resource showing 2038 gene fusions(Mitelman, Johansson, \& Mertens, 2014). These next generation sequencing methods have provided enormous amount of data and with it daunting bioinformatics challenges. Even with obvious computational challenges of analyzing large data, sequencing has identified novel causative mutations that were experimentally validated in diseases like mental retardation where there was no familial history (Vissers et al., 2010). Similar success was also seen in autism which is characterized as a multi spectrum disease. Exomesequencing of 928 autistic individuals identified 279 novel coding mutations (Sanders et al., 2012), a feat if not unachievable would be highly labor intensive in an experimental procedure.

Sequencing methods currently produce paired-end, short reads (30-100 base-pairs) such that both ends of a sequence segment are read with an intervening region of unread sequence in the middle. The two reads are paired such that they have the same identifier linking the two reads. Sequencing technologies today have adopted
paired end short read sequencing as the preferred method due to the relative low cost. The bioinformatics challenge is to map these short sequences correctly to the reference human genome and be able to discern variants. The algorithm developed for detecting variants was first laid down by Tuzun et al. (Tuzun et al., 2005) as:

- mapping the subject genome to the reference genome
- defining signatures for structural variation using the distribution based on size of the short reads (inserts)
- identifying the reads which support the structural variant as described by the signature and clustering all reads which support the same variant together, and finally
- removing false positives based on percent identity of the variants read with the reference genomes

Current tools use this basic model for defining structural variation with statistical adjustments for calculating sensitivity and specificity, and sometimes expanding the signatures for calling structural variants. These work very well in normal genomes and reasonably well in cancer genomes. However, cancer genomes are more complex as samples represent an admixture of normal and abnormal cells, the latter of which may differ genomically. The ability to detect somatic mutations in a given sample depends on tissue type, type of mutation (i.e. germline versus somatic), amount of intercellular heterogeneity in the tumor, the sequencing technology itself, and most importantly the algorithm used to detect variants. The choice of sequencing technology/platform also depends on the goal of the study. The comparison of sequencing platforms in metagenomic studies revealed that the use of short-read generating technology was better due to significantly increased number of
reads and therefore greater coverage of genomic regions (Mende et al., 2012). In contrast, a comparison of these sequencing technologies for low-coverage experiments in clinical setting to detect major copy number changes as part of prenatal diagnosis found that increased number of reads produced in one technological platform also made it more prone to GC-content bias due to greater number of PCR-cycles (S. Chen et al., 2014). After choosing the right platform for the study, the next challenge is to determine the appropriate tools to analyze the data and call variants. All major bioinformatics development has been open source, and tools such as National Center for Biotechnology Information, NCBI's BLAST (Basic Local Alignment Search Tool), University of California Santa Cruz, UCSC's Genome Browser, and Burrow-Wheeler aligner (BWA) have become the de-facto standards for alignment and graphical viewing of the genome. However, new variant detection algorithms are constantly being designed to address specific problems. For example, a study comparing 12 algorithms for quantifying somatic copy number variation using whole genome sequencing data (WGS) found that there were significant differences in sensitivity and specificity of these algorithms (Alkodsi, Louhimo, \& Hautaniemi, 2014). Sensitivity depended on the size of the variant being detected, while breakpoint detection accuracy was determined by the algorithmic approach of the tools. These studies demonstrate the bioinformatics challenges of analyzing DNA sequencing data, specifically variant detection. The multitude of data being generated by next generation sequencing and the analogous growth in bioinformatics tools to deal with this data shows the nascent state of the field in terms of standardized methodologies to analyze this data.

### 1.4 Variables used in the study

Taking the computational modeling for detection of SVs first presented by Tuzun et al. (Tuzun et al., 2005) further, Lee et al. (S. Lee, Cheran, \& Brudno, 2008)
proposed a probabilistic method for variant calling and controlling the false discovery rates. Many currently popular tools use this methodology with modifications. For example, SVDetect (Zeitouni et al., 2010) uses a windowing strategy in the paired end library together with the clustering mechanism explained above to define variants. BreakDancer (K. Chen et al., 2009) uses a more stringent method for defining probabilities and therefore reducing the output of false positives. CREST (Wang et al., 2011) on the other hand used another signature called soft clipping ${ }^{1}$ to extract anomalous reads, build contigs ${ }^{2}$ of soft-clipped reads and subsequently define variants based on similar probability calculations discussed above. Thus, all of these tools work off the alignments produced by the mapping tool. However, this can be a problem when there are multiple mapping regions with similar identity in the BWA mapping tool algorithm (Li \& Durbin, 2009). During these instances the tool will randomly assign the read to any location. This is a characteristic of the human genome: there are regions with considerable sequence homology and repeats such that the mapping tool fails to align sequences at a unique region on the genome. The variant detection tools completely ignore the fact that the alignment reported by the mapping tool may be a random position.

These repeat regions comprise up to $50 \%$ of the human genome (Smit, Hubley, \& Green, 2014) and play a very significant role in development of these structural variants due to errors in the DNA repair process (G. McVean, 2010). McVean has reviewed these recombination events that lead to mutation and identified hotspots in the genome that were more prone to rearrangement breakages. Further, these

[^0]regions tended to be closer to promotor regions but not in the transcribed regions and were driven by a particular class of repair system called non-allelic homologous recombination (G. A. McVean et al., 2004). Non-allelic homologous recombination events (NAHR) are the result of errors in DNA repair involving large chromosomal regions characterized by low copy repeats and up to $95 \%$ sequence identity (Stankiewicz \& Lupski, 2002), (Colnaghi, Carpenter, Volker, \& O'Driscoll, 2011a). Thus, there was information within the genome with a specific signature that was leading to error prone repair. Further, Ou et al. (Ou et al., 2011) were able to specifically prove NAHR as the cause of same unbalanced translocation in four unrelated families. The group was also able to map these 'NAHR' regions on the human genome based on the signature of low copy repeat regions with greater than $94 \%$ sequence identity. They were also able to find validated translocations in the database in these predicted rearrangement hotspots.

Current tools ignore this very relevant information driving these structural variation events in the genome. There is a general consensus in the scientific community that variant calling tools fail in repetitive regions (non-unique regions) of the genome because the mapping algorithms cannot reliably map these regions. This leads to a decrease in the signal-to-noise ratio, and thus the number of reads supportive of an aberration. While this is true, there is enough information known about these regions to account for them algorithmically in a variant detection tool.

DNA damage is acquired primarily during replication process of the cell. Cells have checks and balances to control for errors during repair and a breakdown of these repair mechanism leads to non-allelic homologous regions of the genome undergoing recombination and going unchecked. A low copy repeat region will share homology to many regions on the genome and even share significant identity with these regions.

The 3-D structural organization of the genome within the cell, as studied by $\mathrm{Hi}-\mathrm{C}$ experiments showed that regions of chromosome within close proximity to each other were more likely to interact compared to regions that were far away (Lieberman-Aiden et al., 2009). Another study probing the effect of DNA replication timing on generation of copy number variations in the cancer genome also found the correlation of spatial organization of the chromosomes within the genome affecting the occurrence of mutations at specific locations (De \& Michor, 2011). They went further with the hypothesis proving that those regions that were closer to each other spatially in the cell were also likely to have similar replication timing during cell division.

### 1.5 Problem Statement

The purpose of this research was to use the known genomic context driving the formation of chromosomal aberrations, both in terms of repeat structure and 3D spatial organization, to design a biologically sound computational algorithm to detect translocations in the cancer genome.

### 1.6 Research Question and Hypothesis

1. Are the tools able to detect validated structural variants from a known dataset?
2. Will the tools fail in regions of repeats?
3. Is the ability of the tool to detect variants driven primarily by the presence of these variants in high complexity (unique) regions?

## Research Hypothesis

1. Tools fail when the variants fall in the repeat regions. Variants in unique regions are more likely to be detected by the tools compared to variants in repeat regions.
2. A novel algorithm designed to take into account the genomic context that drives the formation of chromosomal alterations (De \& Michor, 2011; Lieberman-Aiden et al., 2009) as well as the genomic architecture of cancers (regions of the genome more susceptible to structural variant formation as seen in recurrent cancers) is likely to perform better than current tools that do not include this information.

## 2. BACKGROUND

### 2.1 DNA Sequencing

DNA sequencing is the decoding of genetic information locked in the DNA and is the machine translation of the nucleotide sequence that makes up the three billion bases of human genetic code. Fred Sanger introduced the chain-termination method for base determination (Sanger, Nicklen, \& Coulson, 1977) (Figure 2.1 (Estevezj., 2012)) which gained wide acceptance as the preferred method for sequencing. Although another method by Maxam-Gilbert (Maxam \& Gilbert, 1977) was introduced at the same time and used base-specific chemical degradation, the Sanger method became more popular due to its ease of use (Nunnally, 2005).

DNA sequencing can be broadly divided into four steps (Nunnally, 2005):

- Reaction
- Separation
- Detection
- Data analysis

The reaction step is specific to the type of method being used. Broadly, doublestranded DNA is broken mechanically or chemically into single-stranded DNA, mixed with a DNA polymerase ${ }^{3}$, DNA primer ${ }^{4}$, the four deoxynucleotide bases (adenine, guanine, tyrosine and cytosine) and one dideoxynucleotide (ddNTP) corresponding each of the bases that when incorporated during the polymerase reaction stops the lengthening of the DNA chain. Thus, the reaction yields a multiplicity of different

[^1]sized fragments where the length of the fragment depends on the chance incorporation of a ddNTP into the chain. Each type of radioactively or fluorescently labeled ddNTP (ddATP, ddTTP, ddCTP, ddGTP) is aliquotted into a separate reaction, and all four reactions are required to generate a sequence.

The separation step involves separating the DNA fragments obtained from the reaction step based on size. Earlier methods used polyacrylamide gel electrophoresis, in which the DNA fragments travelled vertically through the gel under a steady current for a set period of time. The distance traveled was dependent on size, with smaller fragments migrating more quickly than larger ones. Thus sequence was read from the bottom of the gel up. Currently, this process is done by the capillary based system where the sample are run through a very fine capillary and can be read simultaneously by the detector as the DNA sample is travelling through the capillary.

The detection of the separated fragments involves exposing the separated sample to X-ray film for radioactive labeling. Once the film was developed, the sequence of the DNA could be read from the bottom up by recording in which lane the smallest fragment appeared, followed by the next smallest, and so on. However, this method is not routinely used anymore. It has been replaced by the use of fluorescently labeled ddNTPs exposed to laser light that is simultaneously detected by the detecting machine. In the analysis step, all the data is compiled into a single continuous sequence.


Figure 2.1. Sanger DNA sequencing method. This figure explains the steps in modern Sanger technology. Adapted from Sanger Sequencing, by Estevezj, Retrieved March 27, 2009, from http://commons.wikimedia.org/wiki/File: Sanger-sequencing.svg. Copyright 2012 by Estevezj. Reprinted under Creative Commons Attribution-Share Alike 3.0 Unported license.

The automation of the Sanger-sequencing method led to the human genome project with the aim of sequencing the entire human genome and was completed in 2001 (Venter et al., 2001). The project was completed much earlier than expected using the above method together with shotgun assembly. In shotgun assembly process the entire genome is broken into random smaller pieces, these pieces are amplified by first cloning in bacterial cell (plasmids/bacterial artificial clones) and then through PCR ${ }^{5}$. These amplified PCR products have known fragment sizes (also known as insert size). The major contribution of this group was the development of mate-pair

[^2]reading methods. Different sized libraries of these inserts were created such that multiple clonally amplified fragments of circularized DNA sequences from each library are cut into linear pieces and read from both ends, put together to form contigs with unread gaps and finally contigs assembled together bioinformatically into scaffolds in the analysis step (Istrail et al., 2004). The size of the gap can be estimated with reasonable confidence based on the segment length (i.e. insert size) of the DNA library. This mate-pair sequencing and assembling method was a major step in the direction of high throughput analysis. The completion of the human genome assembly was a significant achievement as the scientific community now had the entire human genome decoded and publicly available through NCBI.

Technology for faster sequencing has improved over the years with a corresponding decrease in the cost due to the development of these new sequencing methods. These newer-generation sequencing technologies are now collectively called nextgeneration (next-gen) sequencing. The most commonly used next-gen applications include Roche/454, Illumina/Solexa, Life/APG and Helicos BioSciences (Metzker, 2010). The flow of steps more or less is the same as Sanger sequencing, including template preparation, sequencing, viewing and data analysis, with major improvements in template preparation and sequencing.

Template preparation has seen major advances in next generation sequencing methods with shift from bacterial artificial chromosomes (BACs) (Monaco \& Larin, 1994) (Shizuya \& Kouros-Mehr, 2001) due to inherent problems with BAC procedures, including loss of genomic material in the BAC during cell replication and introduction of replication errors as human errors during the mapping process. The two major types of template preparation in NGS technology are:

1. Clonally amplified templates
2. Single-molecule templates (Metzker, 2010)

Clonally amplified templates use single stranded DNA molecule with universal primer attached to beads (Figure 2.2)


Figure 2.2. Emulsion bead template preparation for next generation sequencing technology. Adapted by permission from Macmillan Publishers Ltd, "Sequencing technologies - the next generation", by M.L.Metzker, 2010, Nature reviews.Genetics, 11, p.33, copyright 2010.

These PCR-amplified beads can then be placed on a glass slide for the NGS sequence reading to be performed.

Solid phase amplification or bridge amplification can also be performed (Figure 2.3), where the primers are attached to glass slide and DNA fragments along with polymerase are added to the glass slide to produce spatially-separated clones of amplified DNA fragments.


Figure 2.3. Bridge amplification template preparation for next generation sequencing technology. Adapted by permission from Macmillan Publishers Ltd, "Sequencing technologies - the next generation", by M.L.Metzker, 2010, Nature reviews.Genetics, 11, p.33, copyright 2010.

A single molecule template has a major advantage of not requiring PCR amplification and therefore reducing errors. Thus there is no need for clonal amplification in this method and it is believed to be more representative of the original sample. This is achieved by either attaching primers to the glass slide and adding single-stranded DNA molecules to these immobilized primers (Figure 2.4) or attaching the singlestranded DNA molecules to the glass slide and then adding the primers to these immobilized DNA strands (Figure 2.4). A new approach uses the DNA polymerase bound to a glass slide and the single-stranded DNA template can be introduced to this polymerase and read in real-time (Figure 2.5).


Figure 2.4. Single molecule template preparation for next generation sequencing technology. Adapted by permission from Macmillan Publishers Ltd, "Sequencing technologies - the next generation", by M.L.Metzker, 2010, Nature reviews.Genetics, 11, p.33, copyright 2010.


Figure 2.5. Real-time template preparation for next generation sequencing technology. Adapted by permission from Macmillan Publishers Ltd, "Sequencing technologies - the next generation", by M.L.Metzker, 2010, Nature reviews.Genetics, 11, p.33, copyright 2010.

The next step in the process is sequencing the single-stranded DNA and reading it using imaging of fluorescence probes attached to the nucleotides. Currently there are four methods used for sequencing (Metzker, 2010):

1. Cyclic reversible termination
2. Sequencing by ligation
3. Single-nucleotide addition/ pyrosequencing
4. Real-time sequencing

Cyclic reversible termination (Figure 2.6) uses the cyclical process of incorporating fluorescent nucleotides, imaging, and termination using ddNTPs. Each step of a single nucleotide incorporation, termination, imaging, and washing is repeated until the entire template is read. This is most commonly used in clonally amplified templates. The method relies on the use of modified ddNTPs with more efficient cleavage of the fluorescent labels compared to Sanger sequencing.


Figure 2.6. Cyclic reversible termination method in next generation sequencing technology. Adapted by permission from Macmillan Publishers Ltd, "Sequencing technologies - the next generation", by M.L.Metzker, 2010, Nature reviews.Genetics, 11, p.36, copyright 2010.

Sequencing by ligation uses DNA ligase ${ }^{6}$ to attach to a fluorescent dye-labeled probe, washing extra probes away and read by imaging in a cyclical way (Figure 2.7). Pyrosequencing uses sulphurylase and luciferase to detect bioluminescence (Figure 2.7) instead of fluorescently labeled nucleotides. Real-time sequencing is the

[^3]most advanced of these technologies not requiring any terminators to stop the process of sequencing. DNA polymerases are attached to the glass slides and therefore sequencing can be performed without the need to terminate and the release of fluorescence from the nucleotide read in real-time by imaging.


Figure 2.7. Sequencing by ligation and pyrosequencing method in next generation sequencing technology. Adapted by permission from Macmillan Publishers Ltd, "Sequencing technologies - the next generation", by M.L.Metzker, 2010, Nature reviews.Genetics, 11, p.36, copyright 2010.

Technological advances will continue to improve these NGS technologies with the purpose of reducing sequencing errors. One of the first challenges of parallel sequencing was developing efficient algorithms for fast and accurate mapping of this data to the reference genome. Big data however brought with itself bioinformatics analysis challenges, the first being mapping the target genome to the reference.

### 2.2 Mapping Algorithms

The earliest efforts in algorithm development for mapping was done for discovering sequence homology that examines the relatedness of two sequences (protein-coding DNA sequences) in different species in order to understand its function (Pevsner, 2009). Understanding homology was important to define relatedness of proteins through evolution and define changes that occurred through speciation. The same principles were applied to DNA sequence mapping. Reads to be compared are placed along the x -axis and y -axis and scored +1 for each match, -2 for each mismatch and gap. This method was first developed by Needleman and Wunsch and is known as the Needleman-Wunsch algorithm (Needleman \& Wunsch, 1970). At the heart of this algorithm is dynamic programming (Eddy, 2004) which starts with laying out the matrix, deciding on scoring method and finally recursively finding the best path with the best optimal score. This algorithm is dynamic because it finds the best optimal score and keeps it in memory in order to avoid recalculation of scores that happens in a recursive process (Figure 2.8).


Optimum alignment scores 11:

$$
\begin{array}{rccccccc}
\mathbf{T} & - & - & \mathbf{T} & \mathbf{C} & \mathbf{A} & \mathbf{T} & \mathbf{A} \\
\mathbf{T} & \mathbf{G} & \mathbf{C} & \mathbf{T} & \mathbf{C} & \mathbf{G} & \mathbf{T} & \mathbf{A} \\
+5 & -6 & -6 & +5 & +5 & -2 & +5 & +5 \\
\hline
\end{array}
$$

Figure 1 - The filled dynamic programming matrix for two DNA sequences, $x=$ TTCATA and $y=$ TGCTCGTA, for a scoring system of +5 for a match, -2 for a mismatch and -6 for each insertion or deletion.
The cells in the optimum path are shown in red. Arrowheads are 'traceback pointers,' indicating which of the three cases were optimal for reaching each cell. (Some cells can be reached by two or three different optimal paths of equal score: whenever two or more cases are equally optimal, dynamic programming implementations usually choose one case arbitrarily. In this example, though, the optimal path is unique.)

Figure 2.8. Dynamic programming matrix. Reprinted by permission from Macmillan Publishers Ltd, "What is dynamic programming", by S.R.Eddy, 2004, Nature biotechnology, 22, p.909, copyright 2004.

The Needleman-Wunsch algorithm was used for global alignment and when adapted to local alignment as done by Smith-Waterman algorithm (Smith, Waterman, \& Fitch, 1981), proved even more useful in finding DNA sequences of local identity (Pevsner, 2009). The Smith-Waterman method uses one more row (m+1) and column $(n+1)$ for two sequences of length $m$ and $n$ and each subsequent score in the matrix is incremented from the value in the preceding diagonal cell, with a match getting a score of $+1,-0.3$ for a mismatch, and -1.3 for a gap. The Smith-Waterman algorithm does not allow for any negative scores. The highest score in the matrix signifies the end of alignment. This method allows for getting local alignments from a long sequence instead of global alignment of Needleman-Wunsch that is computationally intensive due to its recursive calculations. Local alignment is often
used for quick database searches and serves as the foundation for most population database search mapping tools like BLAST and BLAT.

### 2.3 Structural Variant Detection Tools

Tuzun et al. were the first to describe a computational framework for structural variant (SV) detection (Tuzun et al., 2005). The most gain in information while defining SVs comes from alignment of paired end reads i.e. single fragment of DNA read from both ends with an unread segment in between. The length of the original fragment is known and can further be empirically determined based on the insert size distribution of the experiment. Their approach is still the basic framework for SV detection and is as follows: 1) define probability distribution of insert sizes (length of fragment of DNA sequenced), 2) define discordant reads as those which lie at least 2-4 standard deviations outside this distribution, 3) cluster all reads that support the same SV such that at least 2 discordant reads identify the same SV, 4) identify the SV at the location using percent identity of the discordant reads and number of supporting reads.

Lee et al. proposed a more robust probabilistic method for SV detection (S. Lee et al., 2008). The key methodology proposed was:

1. Defining a probabilistic framework: All reads generated (paired end reads: short fragments of DNA sequenced from both sides with an area of unsequenced portion in between) are assumed to be independent of each other. The set of same type of structural variants is represented by C. Using the independence assumption each mapped read $A$ and $B$ which are part of the same cluster and therefore explaining the same SV will have a joined probability of belonging to the same cluster $C$ defined as $P(A, B \mid C)=P(A \mid C)$. $P(B \mid C)$. The mean insert size length $s$ is known from library generation for
sequencing experiment and also can be derived computationally from experiment output. The probability distribution gives the standard deviation of insert size. The signature for insertions and deletions is as follows: a length of $r$ represents the length of insertion or deletion and $s$ represents the correct mapping length of the insert in the reference. Therefore an insertion event is $s+r$ and deletion event is $s-r$. Using this signature we define the probability of read $A$ and read $B$ being in the insert size distribution of $s+r$ (or $s-r$ ) pairedend reads by using a maximizing function of the joint probability of $A$ and $B$ belonging to the same variant cluster. It thus calculated the probability that a read is produced from the genome sequenced. Similarly, in defining the signature for inversion and translocation we determine the joint probability of observing the 2 mate pairs in the same cluster. The higher this probability the more likely a read belongs to the cluster.
2. Defining the SV: The structural variant is defined using three features
i. Percent similarity of the mate pair sequence to the reference
ii. Product of the probabilities calculated above for a cluster. The larger the probability of a cluster the more it is reliable.
iii. Number of mate pairs in a cluster

Each SV can map to various clusters and therefore using the above three features we need to find the configuration which maximizes each of the features mentioned above. Using this configuration we can assign each SV to just a single cluster.

The primary challenge of the probabilistic model is that the distributions are based on insert sizes and are dependent on intrinsic information output by the mapping tool. If the mapping tool does not know where to place the read, the read can be either disregarded or placed randomly based on all possible best matches. Lee et al.
assigned confidence scores ( $p$-values) to insertion/deletion signatures, but this could not be done to inversions and translocations. If there were systematic errors in sequencing or in alignment algorithms, these will not be picked up using the probability calculation for inversions and translocations and will therefore be ignored. However with the current NGS platforms these errors have been considerably reduced due to higher coverage. Current tools built up on this method with SVDetect, BreakDancer and CREST discussed below.

The BreakDancerMax (K. Chen et al., 2009) algorithm uses the distance and alignment of paired end reads along with the distribution of reads which map to a certain location on the reference to define a SV. Using the SV signature, the program then searches for regions on the reference genome which anchor more SV's than expected. These regions form the putative breakpoints. For a particular region, whichever is the 'dominant' SV signature is used to identify the SV at the location i.e. clustering of all SV's for a particular location and choosing the SV. Further, it uses a confidence score to assign the probability of observing a SV at the location which is higher than chance using a chi-squared statistic with the cutoff p-value $<0.0001$. This tool assumes a Poisson distribution for the clustered variants. Using an analytical model for detecting true positive SV rate in a simulated dataset, they estimated that with an average insert size of 200 bps, insertions and deletions shorter than 40 bps would be difficult to detect.

BreakDancerMini tries to overcome the dependence on the insert size by using a sliding window test. The algorithm defines a window for the reference that is the mean of the insert library of a confidently mapped region+3 s.d - 2(average length of read). Using this window, the frame is shifted 1 bp at a time and the probability that the read lies in this window is calculated. Once again this method works well for insertions and deletions. BreakDancer Max and Mini together did a better job at indel
detection (experimentally validated 110 of 167 indels called by BreakDancer) and to some extent inversions (validated 4 of 13 inversions) but did not do so well with intra-chromosomal events (validated 2 of 6 intra-chromosomal translocations). This is primarily because it uses only paired end data where both read ends map within a defined distance. If only one end of the pair maps then this read is not considered. Translocations will show this pattern where one end of a paired end read mapped to one chromosome and the other end mapped on another chromosome. This information may be lost in the alignment process as the aligner aligned just one end of the read and information on the paired end is lost. Therefore, this tool is most likely to miss this inter-chromosomal event. Thus BreakDancer performs well with indels but does not have a very strong method for detecting translocations. SVDetect (Zeitouni et al., 2010) on the other hand uses a simplified strategy for clustering which is non-probabilistic. Clusters are formed from paired end reads which had incorrect distance (2-3 s.d.) and/or orientation from the reference genome. SVDetect then divides the reference genome into overlapping windows of fixed size and groups all the paired end reads which map to the same overlapping reference. SVDetect also uses clustering parameters like minimum number of reads supporting a SV, filtering those reads whose orientation is different from the majority in a cluster. The filtering process is user defined and the thresholds for filtering can be changed.

Another important feature added to SVDetect is the ability to predict copy number variations/duplication events. The algorithm does this by finding the ratio of depth of coverage of sample to a control dataset in a sliding window along the genome, though they do not define the length of the window. However, the detection of duplications needs a control dataset that has not been well defined in the paper. Thus SVDetect is able to identify insertions, deletions, inter-chromosomal events,
duplications and translocations using a clustering pattern and filtering process. However, SVDetect will also miss reads which were the alignment tool failed to assign a read unambiguously.

CREST (Wang et al., 2011), unlike other applications, does not use the concept of discordant paired reads. Instead it uses direct mapping of reads to the reference to identify the breakpoints. The Burrow Wheeler alignment tool soft clips reads during alignment, i.e. when there is partial alignment or when one end of the pair aligns perfectly and the other end is partially aligned the tool clips the sequence at the point of partial alignment. CREST collects all these soft clipped reads together, creates longer contigs of these reads using an assembly algorithm (CAP3) and realigns these contigs using BLAT (BLAST Like Alignment Tool) alignment algorithm. Thus, the identification of first soft clipped region defines the first breakpoint at location 1. All the soft clipped reads are collected and assembled into contigs using CAP3 algorithm and aligned using BLAT. Wherever these soft clipped contigs map to the genome is the second breakpoint, which will have soft clipped regions mapping to location 1. CREST requires the second contig generated from soft clipped regions at location 1 to be within a certain distance (user-defined) of the second location. The final call of variants includes only those reads with $>97 \%$ sequence similarity and a BLAT score >30 and defines the probability of observing at least c soft clipped reads with sequence coverage $C$ (at that location) to be less than or equal to 0.05 The signatures defined by CREST use the same basic signature of all the SVs but do not use the mapping distance of 2-3 s.d. It uses direct alignment information and therefore is able to identify the SVs at breakpoint coordinate level. In experiments comparing the platform it performed better than BreakDancer in detecting insertions, deletions, and inversions. CREST is more apt for inter-chromosomal event detection and therefore performs better in translocation events seen in cancer genomes.

However due to the alignment heavy algorithm, it performs badly in regions of repeats as the tool will perform only as good as the assembly/alignment tool. All of the above methods use the mapping tool output as the input and derive possible anomalous reads based on either the insert size distribution or flags set by the mapping tool. Thus the ability of the tools to detect variants depends entirely on the robustness of the alignment. In regions where the aligner cannot reliably map sequence, such as repetitive regions, it is assumed that the ability of the tools to detect variants is compromised. This study aimed to test this theory and to design a more effective approach to variant detection by utilizing the known biology driving variant formation, focusing specifically on translocations.

## 3. TESTI NG CONTEXT: REPEAT STRUCTURE OF GENOME

### 3.1 Overview

The overall approach for this study was three pronged: first, test the currently available tools against an experimentally validated dataset; second, develop an algorithm for translocation detection based on known biological information; and finally test our algorithm against the known tools. The first step was essentially to identify and define the reason why tools would fail to detect experimentally validated variants. Our hypothesis was that the tools would fail to recognize variants that were in genomic regions of repeat due to low complexity of these regions and therefore the inability of the aligner to map these regions uniquely. Thus stated differently, variants in repeat regions were less likely to be detected by the tool compared to variants in unique regions.

| Test current tools to <br> understand "context" of <br> the structural variants <br> failing to be detected <br> by the tools |
| :---: | :---: |$\Rightarrow$| Develop novel <br> algorithm based on <br> this "context" |
| :---: |

### 3.2 Testing Current Tools

The first step for proof of concept was to test the available structural variant detection tool on an experimentally validated dataset. The 1000 -Genomes project was designed to use extensive sequencing of many individuals around the world with different ethnicities in order to characterize all types of variants found in these individuals and to relate it to the phenotype ( 1000 genomes project). This project is a major international collaboration between universities around the world providing samples, sequencing data, and bioinformatics analysis in order to map the entire spectrum of genetic variation in the human population. The data has been made
publicly available on the website (www.1000genomes.org) (1000 Genomes Project Consortium, Abecasis, Altshuler, Auton, Brooks, Durbin, Gibbs, Hurles, \& McVean, 2010b). The project had low-coverage whole genome sequencing, exome sequencing, and high-coverage sequencing of trio-subjects (mother-father-child) from different ethnic populations around the world. For purpose of testing we used deep sequencing data of the trio subject from Nigeria (YRI: NA19238, NA19239, NA19240).

### 3.2.1 FASTQ file format

Most sequencing machine output the data in a widely accepted format known as the Sanger-FASTQ (Cock, Fields, Goto, Heuer, \& Rice, 2010). This is a text file with information about the sequence read and quality score of each base of the sequence read. Each read starts with a ‘@' followed by the identifier and description of the sequence which may be platform specific. The next line is a string of 'ATCGs', which is the actual read from the sequencing machine. The next line is a + sign which may be followed by a repeat of the sequence identifier and description line. The last line represents quality scores for each of the bases in the read sequence (Figure 3.1).

```
@SRR014849.1 EIXKN4201CFU84 length=93
GGGGGGGGGGGGGGGGCTTTTTTTTGTTTGGAACCGAAAGG
GTTTTGAATTTCAAACCCTTTTCGGTTTCCAACCTTCCAA
AGCAATGCCAATA
+SRR014849.1 EIXKN4201CFU84 length=93
3+&$#"""""""""""7F@71,'";C?,B;?6B;:EA1EA
1EA5'9B:?:#9EA0D@2EA5':>5?:%A;A8A;?9B;D@
/=<?7=9<2A8==
@title and optional description
sequence line(s)
+optional repeat of title line
quality line(s)
```

Figure 3.1. Sanger FASTQ format. Reprinted by permission from Oxford University Press, "The Sanger FASTQ file format for sequences with quality scores, and the Solexa/Illumina FASTQ variants", by P.J.Cock et al., 2010, Nucleic Acids Research, 38, p.1769, copyright 2010.

The last line equates quality values for each base read by the sequencing machine. This is a PHRED-based quality score converted to ASCII characters. PHRED was a computation tool developed by Ewing et al. (Ewing \& Green, 1998; Ewing, Hillier, Wendl, \& Green, 1998) to automatically assign quality values to sequencing trace files such as chromatograms that are generated by the sequencing machines. This is defined in the equation below.

$$
\begin{aligned}
& q=-10 \times \log _{10}(p)-\cdots-----(\text { Ewing et al.) } \\
& q=\text { quality score of the base } \\
& p=\text { estimated error probability for the base }
\end{aligned}
$$

According to this formula, a base having a $1 / 1000$ probability of being erroneous will have a quality score q of 30 . The lower the probability score, the higher the quality.

### 3.2.2 1000-Genomes data analysis

The files obtained from 1000-Genomes website were in FASTQ format. This study used Yoruba subject data (YRI: NA19238, NA19239, NA19240). The FASTQ files were downloaded from ftp://ftp.1000genomes.ebi.ac.uk/vol1/ftp/pilot_data/data/.

Datasets for each subject had sequencing files of various insert sizes. To maintain consistency we only used deep sequencing high coverage paired end library of insert size 260. Briefly, experiment ID-SRX001106 with 8 billion reads was used for NA19238, experiment ID-SRX000654 with 7.4 billion reads was used for NA19239 and experiment ID-SRX001102 with 6.1 billion reads was used for NA19240. Appendix A lists the files used for this analysis.

Sequencing reads from these files listed in Appendix A Table 1.1 were aligned to the reference genome NCBI build 36.1/ UCSC hg18. The reference human genome is assembled and maintained by the Genome Reference Consortium (GRC) (http://www.ncbi.nlm.nih.gov/projects/genome/assembly/grc/human/) which includes NCBI, Wellcome Trust, Sanger Institute, Genome Institute at Washington University and European Bioinformatics Institute (EBI) at NCBI. The complete genome and chromosome level FASTA files are stored at UCSC website which can be downloaded at (http://hgdownload.cse.ucsc.edu/goldenPath/hg18/chromosomes). The GRC, which is an extension of the Human Genome project (http://www.genome.gov/10001772), produces overlapping segments of high quality, longer length DNA sequence from a group of volunteers. Using in-house tools they release a comprehensive and rigorous representative of a consensus normal human reference genome for use in the public domain. The builds are constantly getting updated as more data becomes available on previously unsequenced regions of the human genome or improved on currently sequences regions.

Data for these various builds are stored in the FASTA format, which is a text file with a header line starting with ' $>$ ' containing information about the sequence followed on the next line by the longest continuous available sequence data as shown in Figure 3.2. Standard codes to represent amino acid sequences and nucleotides are used and can also have lower case letters representing the same nucleotide/amino acid.

```
>gi| 12345|ref|NM_12345.01|Homo sapiens XYZ
ATTTCGATTAATCGAGAAAAAAAATATTTTAGGGGGCCATTTATATACCCCCCCC
TACACCCACAC
```

Figure 3.2. FASTA format example. First line starts with " $>$ " followed by information about the sequence and next line contains actual sequence

Fastq files for the same insert size and same library were downloaded from the ftp site $^{7}$ for each of the trio subjects; NA19238, NA19239, NA19240. These were aligned to the reference genome UCSC hg18 using a popular tool for short read alignment, BWA-0.5.9 (Li \& Durbin, 2009). ${ }^{8}$ Briefly, BWA uses the Burrow-Wheeler transformation algorithm to compresses the data such that repetitive information is stored in a compressed format in prefix and suffix arrays in order to perform searches on the entire human genome. The output of the mapping tool is in SAM format described in Section 4.2.1. SAM is then converted to BAM, which is the binary representation of a SAM file using ‘samtools’ suite. ${ }^{9}$

The output from BWA mapping was used to further detect structural variants using an open-source toolkit. SVDetect (Zeitouni et al., 2010) is one such toolkit which uses paired-end mapping data to identify reads which occur at a distance greater than expected insert size and/or are in the incorrect orientation with respect to each

[^4]other. It then breaks the genome to create overlapping windows. Each anomalous paired-end read can anchor to these windows such that the two windows where the two paired-end reads map form a link. A filtering process removes links that has less than a certain number of reads supporting a link and other features defined in the tool. Structural variants are called based on mapping signatures for a particular type of variant. This tool has a tendency to output many false positives as it only uses a clustering method to call variants. Thus, any variant with 2 or more reads supporting it is called as a variant without using the underlying variants distribution to calculate confidence scores of calling a variant as done in BreakDancer. We therefore used this tool which reports all possible variants in order to test our hypothesis that tools fail to detect variants because these variants lie in the repeat regions.

Experimentally validated structural variants detected in the same trio subjects were used as the validation standard. The variant dataset was downloaded from the ftp site. ${ }^{10}$

Variants were detected computationally using various algorithms developed in-house by 1000-Genomes team ( 1000 Genomes Project Consortium, Abecasis, Altshuler, Auton, Brooks, Durbin, Gibbs, Hurles, \& McVean, 2010a). The variants were defined as mobile element insertions, tandem duplication, deletions, and novel sequences. We extracted only those variants from the variant files that were experimentally validated as indicated in the Description field of the file.

[^5]
### 3.2.3 SVDetect Analysis

SVDetect defines its set of variants as described in the Table 3.1.

Table 3.1: SVDetect defined structural variants

| NORMAL_SENSE | Correct ends orientation using <mates_orientation> as reference |
| :--- | :--- |
| REVERSE_SENSE | One of the ends has an incorrect orientation |
| DELETION | Deletion (NORMAL_SENSE \& mean insert size > $\mu$ +threshold* $\sigma$ ) |
| INSERTION | Insertion (NORMAL_SENSE \& mean insert size < $\mu$-threshold* ) |
| INVERSION | Inversion (REVERSE_SENSE) |
| INV_FRAGMT | Inversion of a genomic fragment, defined by balanced signatures <br> (BAL) |
| INS_FRAGMT | Insertion of a genomic fragment, defined by balanced signatures <br> (BAL) |
| INV_INS_FRAGMT | Inverted INS_FRAGMT (BAL) |
| LARGE_DUPLI | Large duplication |
| DUPLICATION | Duplication, medium size |
| SMALL_DUPLI | Small duplication (mean insert size <br> between subgroups) $\mu$-threshold* |
| INV_DUPLI | Inverted duplication (REVERSE_SENSE \& mean insert size < $\mu-$ <br> threshold* \& UNBAL) |
| TRANSLOC | Translocation |
| INV_TRANSLOC | Inverted translocation |
| COAMPLICON | Co-amplicons, two different fragments repeated in the same <br> strand sense (BAL) |
| INV_COAMPLICON | Inverted co-amplicons, two different fragments repeated in the <br> opposite strand sense (BAL) |
| SINGLETON | Singleton (mean insert size < $\mu$-threshold* $\sigma$ ), for Illumina mate- <br> pairs only |
| UNDEFINED | Undefined inter/intra-chromosomal SV type |

The purpose of this analysis was to see if the variants detected by SVDetect had some overlap in the same genomic region as those in the 1000-Genomes experimentally-validated variant dataset. To further characterize these variants, we looked for overlap of the variants with mapped repeat structure in the human genome (Smit et al., 2014) ${ }^{11}$. We used the RepeatMasker mapped repeat regions on the hg18 build human genome to identify the repeat regions in the human genome. The repeat elements in RepeatMasker database have been classified as LINEs: Long

[^6]interspersed nuclear elements; SINEs: Short interspersed nuclear elements; LTR: Long terminal repeats; DNA repeats and other.

### 3.3 Extracting Overlapping Regions

SVDetect-0.7f was used on the mapped BAM files to detect variants. Each variant file from 1000-Genomes classified as ‘Deletion’, ‘Novel Sequences’, ‘Mobile Element Insertion’ and 'Tandem Duplications’ was used separately to extract overlapping variants detected by SVDetect. 1000-Genomes variant files did not have subject identifier in file and thus had structural variants detected in all three trio subjects combined. SVDetect was run on each BAM file separately and then combined. We were only looking for regions of the genome that showed an overlap irrespective of the type of variant as defined in SVDetect. Overlap was defined as at least 10 percent of the insert size. In this case the insert size library chosen for download was 260 and therefore 10 percent overlap was at least 26 basepairs on either side of the read as represented in the Figure 3.3. Ten percent of insert size was used as an arbitrary cutoff.


Figure 3.3: Overlap definition for SVDetect; a: left overlap; b: right overlap; c: SVDetect variant completely within; d: 1000-genome variant completely within

A perl script was used to do the overlap extraction. Similarly, since the purpose was to define the context of the genomic region that is causing the tools to fail, RepeatMasker mapped repeat regions was used to do a similar overlap of with 1000genome variants. However in order to identify the overlapping regions with repeat structure within, the entire repeat region had to be present within the variant as shown in Figure 1.1.


Figure 3.4: Overlap definition for repeat; Blue arrows: 1000 genomes, Black arrows: repeat elements in RepeatMasker.

### 3.3.1 Results: SVDetect performance

The result from extraction analysis showed that SVDetect did poorly in detecting insertion elements (17\%) while it did very well in identifying validated novel sequences high in detection percent ( $67 \%$ ) (Table 3.2).

Table 3.2: Overlap results for SVDetect

| Type of Event | $1000-$ <br> Genomes <br> events | Events detected by <br> SV-Detect: True <br> Positive | Events not detected <br> by SV-Detect: False <br> Negative |
| :--- | :--- | :--- | :--- |
| Deletion | 9695 | $4657(48 \%)$ | 5038 |
| Mobile Element <br> Insertion | 492 | $85(17 \%)$ | 407 |
| Tandem <br> Duplication | 65 | $23(35 \%)$ | 42 |
| Novel <br> Sequences | 66 | $42(67 \%)$ | 24 |

SVDetect is known to give many false positives and thus is a highly sensitive tool for detection of structural variants. However, even with such a highly sensitive tool, less than $50 \%$ of the variants detected by SVDetect were in the same region as deletions events identified by 1000 -Genomes. SVDetect was in the true variants regions anywhere from 17 to 67 percent of the time depending on the type of variant. Even a highly sensitive tool was clearly not able to detect many of the validated variants.

### 3.3.2 Understanding context: Repeat

We therefore tried to understand if the underlying repeat structure within these variants was driving the ability of the tool to detect a variant. If a variant is in a repeat region, the mapping tool will not be able to definitively place the read representing the variant in a specific region due to low complexity of the repeat region. We hypothesized that validated variants that were not detected by the tool were more likely to have repeat elements within them than variants that were detected by the tool. As shown in Figure 3.4, we extracted those reads that had the
repeat element completely contained within the variants or variant completely contained within the repeat. The results are shown in Table 3.3, 3.4.

Table 3.3: Repeat elements within undetected events

| Type of Event | $1000-$ <br> Genomes <br> events | Events not detected by <br> SV-Detect: <br> Negatives (a) | Events with repeat <br> structure within (b) <br> (a/b\%) |
| :--- | :--- | :--- | :--- |
| Deletion | 9695 | 5038 | $2678(53 \%)$ |
| Mobile Element <br> Insertion | 492 | 407 | $134(33 \%)$ |
| Tandem <br> Duplication | 65 | 42 | $11(26 \%)$ |
| Novel <br> Sequences | 66 | 24 | $13(54 \%)$ |

Table 3.4: Repeat elements within detected events

| Type of Event | $1000-$ <br> Genomes <br> events | Events detected by <br> SV-Detect: True- <br> Positives(a) | Events with any repeat <br> structure within (b) <br> (a/b\%) |
| :--- | :--- | :--- | :--- |
| Deletion | 9695 | 4657 | $3754(81 \%)$ |
| Mobile Element <br> Insertion | 492 | 85 | $36(42 \%)$ |
| Tandem <br> Duplication | 65 | 23 | $13(57 \%)$ |
| Novel Sequences | 66 | 42 | $26(62 \%)$ |

Although 53\% of deletion events that were not detected by the tool showed some form of repeat structure, $81 \%$ of deletion events which were detected also showed a repeat structure. A similar pattern was seen with all other variant types where the events which were detected by the tool had a greater percent of repeat elements compared with those which were not detected by the tool. A chi-squared analysis (Tables 3.5, 3.6, 3.7, 3.8) showed a significant association between repeat structure and the ability of the tool to detect the variant for deletions and tandem duplication events but not significant for novel sequences and mobile element insertion events at $p<0.05$ level. Thus if there was a repeat structure within the variant, the tool was
more likely to detect it, which was contrary to our hypothesis. We therefore failed to reject the null of no association between repeat structure and ability of the tool to detect the variant.

Table 3.5: Chi-squared test: Relation of deletion events with repeat

| Deletion Events |  | Repeat structure <br> within | No Repeat structure <br> within |
| :--- | :--- | :--- | :--- |
| Detected by tool <br> positives) | (True | 3754 | 903 |

Note: Table tests if ability of the tool to detect deletion events is influenced by the repeat structure within it. The Chi-square statistic is 816.873 with $p$-value of 0 . This result is significant at $\mathrm{p}<0.05$.

Table 3.6: Chi-squared test: Relation of insertion events with repeat

| Insertion Events | Repeat structure <br> within | No Repeat structure <br> within | Total |
| :--- | :--- | :--- | :--- |
| Detected by tool (True <br> positives) | 36 | 49 | 85 |
| Not detected by tool (False <br> Negatives) | 134 | 273 | 407 |

Note: Table tests if ability of the tool to detect insertion events is influenced by the repeat structure within it The Chi-square statistic is 2.764 with $p$-value of $p=0.096$. This result is not significant at $p<0.05$.

Table 3.7: Chi-squared test: Relation of duplication events with repeat

| Duplication Events | Repeat structure <br> within | No Repeat structure <br> within | Total |
| :--- | :--- | :--- | :--- |
| Detected by tool (True <br> positives) | 13 | 10 | 23 |
| Not detected by tool (False <br> Negatives) | 11 | 31 | 42 |

Note: Table tests if ability of the tool to detect duplication events is influenced by the repeat structure within it The Chi-square statistic is 5.871 with $p$-value of 0.015 . This result is significant at $\mathrm{p}<0.05$.

Table 3.8: Chi-squared test: Relation of novel sequences events with repeat

| Novel Sequence Events | Repeat structure <br> within | No Repeat structure <br> within | Total |
| :--- | :--- | :--- | :--- |
| Detected by tool (True <br> positives) | 26 | 16 | 42 |
| Not detected by tool (False <br> Negatives) | 13 | 11 | 24 |

Note: Table tests if ability of the tool to detect novel sequences events is influenced by the repeat structure within it The Chi-square statistic is 0.378 with $p$-value of 0.538 . This result is not significant at $p<0.05$.

Since defining genomic context was the goal of the analysis, we further broke down the detected and undetected events by type of repeat structure as classified by RepeatMasker. Note, however, that the same variants can have more than one type of repeat structure. This breakdown also showed similar results. While LINEs and SINEs were more abundant of the repeat structure, undetected events (falsenegatives) showed a lower percentage of these events than detected events (truepositives) (Tables 3.9, 3.10, 3.11, 3.12).

Table 3.9: Repeats structure by type: Deletion

| 1000 <br> genome <br> events | DNA_repeats |  |  | LINEs | Yes | No | Yes | No | Yes | No |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| YTR | yes | No | Yther |  | No |  |  |  |  |  |
| Un- <br> detected <br> $(5038)$ | 245 <br> $(5 \%)$ | 4793 | 825 <br> $(16 \%)$ | 4213 | 1096 <br> $(22 \%)$ | 3942 | 423 <br> $(8 \%)$ | 4615 | 1113 <br> $(22 \%)$ | 3925 |
| Detected <br> $(4657)$ | 1560 <br> $(34 \%)$ | 3097 | 2480 <br> $(53 \%)$ | 2177 | 2710 <br> $(58 \%)$ | 1947 | 1954 <br> $(42 \%)$ | 2703 | 2549 | 2108 |
| $(55 \%)$ |  |  |  |  |  |  |  |  |  |  |$|$

Note: LINEs: Long interspersed nuclear elements; SINEs: Short interspersed nuclear elements; LTR: Long terminal repeats; Undetected: False negative events; Detected: True positive events (1000-Genomes event)

Table 3.10: Repeats structure by type: Mobile element insertion

| 1000 <br> genome <br> events | DNA repeats |  |  | LINEs |  | SINEs | LTR |  | Other |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
|  | Yes | No | Yes | No | Yes | No | yes | No | Yes | No |
| Undetected <br> $(407)$ | 11 <br> $(3 \%)$ | 396 | 69 <br> $(17 \%)$ | 338 | 20 <br> $(5 \%)$ | 387 | 24 <br> $(6 \%)$ | 383 | 31 <br> $(8 \%)$ | 376 |
| Detected <br> $(85)$ | 12 <br> $(14 \%)$ | 73 | 15 <br> $(18 \%)$ | 70 | 18 <br> $(21 \%)$ | 67 | 12 <br> $(14 \%)$ | 73 | 14 <br> $(16 \%)$ | 71 |

Note: LINEs: Long interspersed nuclear elements; SINEs: Short interspersed nuclear elements; LTR: Long terminal repeats; Undetected: False negative events; Detected: True positive events (1000-Genomes event)

Table 3.11: Repeats structure by type: Tandem duplication

| 1000 <br> genome <br> events | DNA_repeats |  |  | LINES |  | SINES | LTR |  | Other |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
|  | Yes | No | Yes | No | Yes | No | yes | No | Yes | No |
| Undetected <br> (24) | 1 <br> $(4 \%)$ | 23 | 3 <br> $(13 \%)$ | 21 | 6 <br> $(25 \%)$ | 18 | 2 <br> $(8 \%)$ | 22 | 5 <br> $(21 \%)$ | 19 |
| Detected <br> $(42)$ | 3 <br> $(7 \%)$ | 39 | 7 <br> $(17 \%)$ | 35 | 9 <br> $(21 \%)$ | 33 | 4 <br> $(10 \%)$ | 38 | 13 <br> $(31 \%)$ | 29 |

Note: LINEs: Long interspersed nuclear elements; SINEs: Short interspersed nuclear elements; LTR: Long terminal repeats; Undetected: False negative events; Detected: True positive events (1000-Genomes event)

Table 3.12: Repeats structure by type: Novel sequences

| 1000 <br> genome <br> events | DNA_repeats |  | LINES | SINES | LTR | Other |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
|  | Yes | No | Yes | No | Yes | No | yes | No | Yes | No |
| Undetected <br> $(42)$ | 1 <br> $(2 \%)$ | 41 | 5 <br> $(12 \%)$ | 37 | 2 <br> $(5 \%)$ | 40 | 3 <br> $(7 \%)$ | 39 | 1 <br> $(2 \%)$ | 41 |
| Detected <br> $(23)$ | 6 <br> $(26 \%)$ | 17 | 10 <br> $(43 \%)$ | 13 | 8 <br> $(34 \%)$ | 15 | 9 <br> $(39 \%)$ | 14 | 9 <br> $(39 \%)$ | 14 |

Note: LINEs: Long interspersed nuclear elements; SINEs: Short interspersed nuclear elements; LTR: Long terminal repeats; Undetected: False negative events; Detected: True positive events (1000-Genomes event)

Since presence of repeat elements was not driving the ability of the tool to detect variants, we decided that we had to be even more specific in defining the "context" of these variants. Repeat element in general, (including every type of repeat) within or around these variants was not the "context" causing these variants to go undetected.

Our goal was to design a tool to detect translocations, and thus the design step started from defining the context for translocations occurring in the genome i.e. understand how translocations occur in the genome (Bunting \& Nussenzweig, 2013) and build the context using this information. When we looked at the overlap between the NAHR regions and the genomic location of the variants missed by SVDetect, we observed a pileup in the regions where the tools failed to detect variants and the NAHR regions (Figure 3.5). Thus, we used NAHR as one part of the biological
context for the novel algorithm. However, there are multiple substrates for NAHR in the human genome. The likelihood of any two regions of NAHR taking part in exchange of genomic material (translocation) depends in part on the probability of these regions being in close proximity in 3-dimensional space. Therefore, the second part of our biological context was to incorporate information about the distribution of the genome in 3-dimensional space as derived from $\mathrm{Hi}-\mathrm{C}$ data.


Figure 3.5: NAHR overlap with 1000-Genomes variants. Pile-up of false negative deletion ebents (middle grey bar plot) shows concordant spike with pile-up of NAHR
 and red dot plot shows count of number of LINE elements (long interspersed nuclear elements) in 1 Kbps windows.

## 4. ALGORITHM DESI GN AND DEVELOPMENT

### 4.1 A Context Definition for Translocations

The purpose of defining "context" for detecting translocation in the genome was to effectively use information on mechanism of translocation formation in designing the tool. Earlier held belief that cancer occurs due to random genomic events is being proven ineffective due to studies which show non-random patterns to breakage sites, also called 'hotspots' (Jeffreys, Kauppi, \& Neumann, 2001). Recombination event is a process by the cellular machinery to create diversity through evolution. These events occur through double stranded breaks and the aberrant repair of these breakpoints leads to structural variations. However these breakpoint regions have specific signatures of low copy number variations which share significant homology (Colnaghi, Carpenter, Volker, \& O'Driscoll, 2011b). The breakpoint acting as substrates for anomalous repair (non-allelic homologous recombination, NAHR) lead to translocations and cause phenotypic changes (Deininger \& Batzer, 1999), (Ou et al., 2011), (Gu, Zhang, \& Lupski, 2008). Using this information for the purpose of designing a structural variant has not been done so far by any of the current tools. There is sufficient evidence to define these regions bioinformatically (Ou et al., 2011) and thus be able to use the signature of NAHR programmatically in tool design which was attempted in this study. The novelty of this approach is the use of inherent biologically processes driving these variants formation and is the first such attempt to program this information in a tool. This study tried to capture 3-D packing of the genome within a cell bioinformatically, thus taking the concept of "context" one step further. Translocations occurring from double-stranded breaks also do not occur randomly. Chromosomes which are more close to each other and with similar replication timings are more likely to interact and exchange genomic material through NAHR (Yaffe \& Tanay, 2011), (De \& Michor, 2011), (Wijchers, 2011). Instead of using empirical distribution of reads and variants in the data to derive
confidence score, this study attempted to use probability distribution of two regions in the genome interacting due to physical proximity. This study has thus introduced the concept of context-based evaluation of structural variation in the human genome as a novel approach.

### 4.2 Designing the Algorithm

The algorithm for the tool is shown in Figure 4.1. Broadly the tool could be divided into 3 major steps:

1. Extraction of anomalous reads from the bam file
2. Re-alignment using BLAT
3. Identifying potential translocations and defining probabilities using Hi-C data


Figure 4.1: Novel algorithm flowchart

### 4.2.1 Read extraction overview

Sequence Alignment Format: There are many tools available for mapping reads to the reference genome like BWA, Bowtie, MAQ. These tools give the output in a common format known as SAM format (Sequence Alignment/Map Format), which had been widely accepted as the default format for generating output files. Briefly the alignment output has a header section and an alignment section. All headers start with @ (eg. @HD) and contain general information about the experimental process used to generate the sequencing reads. Alignment sections contain 11 mandatory fields as listed in Table 4.1.

Table 4.1: SAM format fields: alignment section

| Field | Type | Regexp/Range | Brief description |
| :--- | :--- | :--- | :--- |
| QNAME | String | $[!-? \mathrm{~A}-\sim]\{1,255\}$ | Query template NAME |
| FLAG | Int | $\left[0,2^{16}-1\right]$ | bitwise FLAG |
| RNAME | String | $\backslash^{*} \mid[!-()+-<>-\sim][!-\sim]^{*}$ | Reference sequence NAME |
| POS | Int | $\left[0,2^{29}-1\right]$ | 1-based leftmost mapping POSition |
| MAPQ | Int | $\left[0,2^{8}-1\right]$ | MAPping Quality |
| CI GAR | String | $\^{*} \mid([0-$ <br> $9]+[$ MIDNSHPX $=])+$ | CIGAR string |
| RNEXT | String | $\^{*}\|=\|[!-()+-<>-\sim][!-$ <br> $\sim]^{*}$ | Ref. name of the mate/next read |
| PNEXT | Int | $\left[0,2^{29}-1\right]$ | Position of the mate/next read |
| TLEN | Int | $\left[-2^{29}+1,229-1\right]$ | observed Template LENgth |
| SEQ | String | $\backslash^{*} \mid[$ A-Za-z=.]+ | segment SEQuence |
| QUAL | String | $[!-\sim]+$ | ASCII of Phred-scaled <br> QUALity +33 |

The purpose of the extraction algorithm is to retrieve reads which were assigned by the aligner as improperly mapped. This information is coded in the bitwise 'FLAG' field of the SAM file as shown in Table 4.2.

Table 4.2: Bitwise FLAG of SAM file

| Bit | Description |
| :--- | :--- |
| $0 \times 1$ | template having multiple segments in sequencing |
| $0 \times 2$ | each segment properly aligned according to the aligner |
| $0 \times 4$ | segment unmapped |
| $0 \times 8$ | next segment in the template unmapped |
| $0 \times 10$ | SEQ being reverse complemented |
| $0 \times 20$ | SEQ of the next segment in the template being reversed |
| $0 \times 40$ | the first segment in the template |
| $0 \times 80$ | the last segment in the template |
| $0 \times 100$ | secondary alignment |
| $0 \times 200$ | not passing quality controls |
| $0 \times 400$ | PCR or optical duplicate |
| $0 \times 800$ | supplementary alignment |

The above table is represented in the hexadecimal system. The corresponding binary and decimal conversion is shown in Table 4.3.

Table 4.3: Bitwise Flag of SAM FLAG- binary to decimal conversion

| Bit | Binary | Decimal |
| :--- | :--- | :--- |
| $0 \times 1$ | 1 | 1 |
| $0 \times 2$ | 10 | 2 |
| $0 \times 4$ | 100 | 4 |
| $0 \times 8$ | 1000 | 8 |
| $0 \times 10$ | 100000 | 16 |
| $0 \times 20$ | 1000000 | 32 |
| $0 \times 40$ | 1000000 | 64 |
| $0 \times 80$ | 100000000 | 128 |
| $0 \times 100$ | 1000000000 | 256 |
| $0 \times 200$ | 10000000000 | 512 |
| $0 \times 400$ | 100000000000 | 1024 |
| $0 \times 800$ | 1000000000000 | 2048 |

The value for each bit is set to 0 or 1 based on the description, with 0 being no and 1 being yes. Thus if a sequencing segment which has been mapped has a value of 64 , it is the first segment in the template (Table 4.1, Table 4.2). Each mapped segment
will have a value as represented in the decimal system in the SAM file. Thus a value of 99 equals $64+32+2+1$. The corresponding natural language interpretation based on the tables 4.1 and 4.2 will be as below:

| FLAG | Bit | Description |
| :--- | :--- | :--- |
| 64 | $0 \times 40$ | the first segment in the template |
| 32 | $0 \times 20$ | SEQ of the next segment in the template <br> being reversed |
| 2 | $0 \times 2$ | each segment properly aligned according to <br> the aligner |
| 1 | $0 \times 1$ | template having multiple segments in <br> sequencing |

Thus the read with a SAM flag of 99 meant it mapped correctly. The conversion of SAM flag values to the corresponding interpretation can be done using one of the utilities of PICARD tool suite. ${ }^{12}$

CIGAR String: The sixth field in a SAM file is represented by the 'CIGAR' string. Just as 'FLAG' is used to describe the overall information of sequence read in terms of mapping to the reference genome, 'CIGAR' string explains similar information in terms of mapping at the base-pair level. Thus each base in a sequence read will have any one of the CIGAR values listed in Table 4.4.

Table 4.4: CIGAR String

| Operation | BAM | Description |
| :--- | :--- | :--- |
| M | 0 | alignment match (can be a sequence match or <br> mismatch) |
| I | 1 | insertion to the reference |
| D | 2 | deletion from the reference |
| N | 3 | skipped region from the reference |
| S | 4 | soft clipping (clipped sequences present in SEQ) |
| H | 5 | hard clipping (clipped sequences NOT present in <br> SEQ) |
| P | 6 | padding (silent deletion from padded reference) |
| = | 7 | sequence match |
| X | 8 | sequence mismatch |

[^7]An example of a sequence read mapped to the reference genome can be represented as:

Reference Genome Sequence: ATTTGCATCCCG TATTGGCA Query Read Sequence: ATTTGC TCCCGATATTGGCA
CIGAR String: 6M1D5M1I5M 6 Match, 1 Deletion, 5 Match, 1 Insertion, 8 Match

Reference Genome Sequence: A T T T G C A T C C C GTATTGGCA Query Read Sequence: ATTT CATCCCGTAggggg CIGAR String: 4M1D9M5S 4 Match, 1 Deletion, 9 Match, 5 Soft clipped

The CIGAR string informs about how many insertions and deletions of base pairs is needed to be done in the query sequence for the mapping tool to align the read to a given location on the reference genome.

FLAGs for read extraction: All mapping tools will give the output in SAM format with a FLAG assigned. The purpose of read extraction is to identify those reads that did not have enough information contained in the sequence for the mapping tool to align it at the correct location. This information is extracted using the FLAG field and 'CIGAR' string explained in Section 4.2.1. The FLAG fields used for extraction are shown in Table 4.5.

Table 4.5: FLAGs for read extraction

| Flag | Fields set to derive Flag | Description of fields for Flag |
| :---: | :---: | :---: |
| UNMAPPED READS |  |  |
| 101 | 64+32+4+1 | first in pair, mate reverse strand, read unmapped, read paired |
| 89 | 64+16+8+1 | first in pair, read reverse strand, mate unmapped, read paired |
| 117 | $64+32+16+4+1$ | first in pair, mate reverse strand, read reverse strand, read unmapped, read paired |
| 121 | $64+32+16+8+1$ | first in pair, mate reverse strand, read reverse strand, mate unmapped, read paired |
| 165 | $128+32+4+1$ | second in pair, mate reverse strand, read unmapped, read paired |
| 153 | 128+16+8+1 | second in pair, read reverse strand, mate unmapped, read paired |
| 185 | $128+32+16+4+1$ | second in pair, mate reverse strand, read reverse strand, read unmapped, read paired |
| 181 | $128+32+16+8+1$ | second in pair, mate reverse strand, read reverse strand, mate unmapped, read paired |
| 69 | 64+4+1 | first in pair, read unmapped, read paired |
| 73 | 64+8+1 | first in pair, mate unmapped, read paired |
| 133 | 128+4+1 | second in pair, read unmapped, read paired |
| 137 | 128+8+1 | second in pair, mate unmapped, read paired |
| 77 | 64+8+4+1 | first in pair, mate unmapped, read unmapped, read paired |
| 141 | $128+8+4+1$ | second in pair, mate unmapped, read unmapped, read paired |
| READS WITH NO INFORMATI ON ON MAPPING |  |  |
| 65 | 64+1 | first in pair, read paired |
| 129 | 128+1 | second in pair, read paired |
| 81 | 64+16+1 | first in pair, read reverse strand, read paired |
| 97 | 64+32+1 | first in pair, mate reverse strand, read paired |
| 145 | 128+16+1 | second in pair, read reverse strand, read paired |
| 161 | 128+32+1 | second in pair, mate reverse strand, read paired |
| 113 | 64+32+16+1 | first in pair, mate reverse strand, read reverse strand, read paired |
| 177 | 128+32+16+1 | second in pair, mate reverse strand, read reverse strand, read paired |


| Flag | Fields set to derive <br> Flag | Description of fields for Flag |
| :--- | :--- | :--- |
| READS WITH IMPROPER ORIENTATI ON |  |  |
| 115 | $64+32+16+2+1$ | first in pair, mate reverse strand, read <br> reverse strand, read mapped in proper pair, <br> read paired |
| 179 | $128+32+16+2+1$ | second in pair, mate reverse strand, read <br> reverse strand, read mapped in proper pair, <br> read paired <br> first in pair, read mapped in proper pair, <br> read paired |
| 67 | $64+2+1$ | second in pair, read mapped in proper pair, <br> read paired |
| 131 | $128+2+1$ |  |

### 4.2.2 Read extraction algorithm

Perl script algorithm to extract the reads using the 'Flag' field and 'CIGAR' string was used as described in Table 4.6.

Table 4.6: Read_Extract Algorithm

| Assumption: The input file has to be in BAM/SAM <br> format. | Creating FASTA <br> file from BAM <br> output of mapping |
| :--- | :--- | :--- |
| tool |  |


| ```AND (currentReadIdentifier OR previousReadIdentifier == matchlist array)) { if ('N' pattern not found) {print to output FASTA file with <RID1>/ <RID2> as suffix to identify two reads of a pair and also entire line to SAM file} } } END OF WHILE LOOP``` <br> Close all files | array and extracts only those reads where both reads of a pair are present and did not have more than $30 \%$ ' N 's. It further adds identifier <RID1>, <RID2> to keep paired end reads information |
| :---: | :---: |

The output from the extraction process is saved in both FASTA and SAM format. The FASTA file serves as the input for BLAT mapping tool. Both reads of a pair now have a unique identifier (RID1, RID2) which is an important addition in this step.

### 4.3 Re-alignment Using BLAT

### 4.3.1 The need for BLAT

BLAT which stands for BLAST-like alignment tool was developed in 2001 by James Kent at University of California, Santa Cruz. The popular alignment tool BLAST (Altschul, Gish, Miller, Myers, \& Lipman, 1990) developed by the NCBI uses the Smith-Waterman algorithm (Smith et al., 1981) to do local alignment of a defined length. The extension is done such that further extension of the segment will not improve the score.

BLAT (Kent, 2002) was used primarily because of its speed and comparable output to BLAST. For example a sample output from the simulated data of translocations had FLAGs set as 117 and 153 such that the first read was unmapped and the second read mapped to a region with $100 \%$ match (Figure 4.2). The output from BLAST and BLAT are shown in Figures 1 and 2. BLAT is very similar to BLAST in terms of its search algorithm but differs in the fact that BLAST builds an index of the
query to search through its database, while BLAT used the database to build its index and search through the query which significantly affects search time.

```
chr12_12610644_12611092_6:0:0_3:0:0_23bee 117 chr4 191129790 0 * = 191129790 0
ACTCTCCCATGTCTACTTCTTTATACACAGTCACGGCAACCATCTGATTTATCAATCTTTTCCCCACCTGACCCCCCTTTCCATTCCACAAAACCGCCAT
2222222222222222222222222222222222222222222222222222222222222222222222222222222222222222222222222222222
RG:Z:transl_RG_F
chr12_12610644_12611092_6:0:0_3:0:0_23bee 153 chr4 191129790 2 100M = 191129790 0
CCTAGGAAAACCAGAGACCTTTGTTCACTTGTTTATCTGCTGACCTTCCCTACTCTATTGTCCTGTGACCCGGCCAAATCCCCCTCTGCGAGAAACACCC
2222222222222222222222222222222222222222222222222222222222222222222222222222222222222222222222222222222
RG:Z:transI_RG_F XT:A:U NM:i:2 SM:i:2 AM:i:0 X0:i:1 X1:i:137 XM:i:2 XO:i:0 XG:i:0 MD:Z:51C1A46
```

Figure 4.2: Sample paired-end read of simulated dataset: SAM format

BLAST output for the first segment of the pair (with FLAG=117) showed a match with 56 regions in the 'Human plus Transcript' database and MegaBLAST algorithm (Figure 4.3). The corresponding mate segment of the pair (with FLAG=153) showed match with 200 regions in the same database and same algorithm (Figures 4.4).


Figure 4.3. MegaBLAST output with defaults for query sequence pair1


Figure 4.4: MegaBLAST output with defaults for query sequence pair2

For the same two query sequences, BLAT output is represented below (Figures 4.5, 4.6).

## BLAT Search Results

| ACTIONS | QUERY | SCORE START |  | END QSIZE |  | DENTITY | CHRO | STRAND | ND START | END | SPAN |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| browser details | YourSeq | 88 | 1 | 100 | 100 | 94.0\% | 12 | - | 12610644 | 12610743 | 100 |
| browser details | YourSeq | 86 | 1 | 100 | 100 | 93.0\% | X | - | 41357364 | 41357463 | 100 |
| browser details | YourSeq | 86 | 1 | 100 | 100 | 93.0\% | 2 | - | 143176874 | 143176973 | 100 |
| browser details | YourSeq | 86 | 1 | 100 | 100 | 93.0\% | 17 | + | 26087253 | 26087352 | 100 |
| browser details | YourSeq | 84 | 1 | 100 | 100 | 92.0\% | X | - | 73304225 | 73304324 | 100 |
| browser details | YourSeq | 84 | 1 | 100 | 100 | 92.0\% | X | - | 47967119 | 47967218 | 100 |
| browser details | YourSeq | 84 | 1 | 100 | 100 | 92.0\% | 7 | - | 139778194 | 139778293 | 100 |
| browser details | YourSeq | 84 | 1 | 100 | 100 | $92.0 \%$ | 7 | - | 26268700 | 26268799 | 100 |
| browser details | YourSeq | 84 | 1 | 100 | 100 | $92.0 \%$ | 6 | - | 74368418 | 74368517 | 100 |
| browser details | YourSeq | 84 | 1 | 100 | 100 | $92.0 \%$ | 4 | - | 65448670 | 65448769 | 100 |
| browser details | YourSeq | 84 | 1 | 100 | 100 | 92.0\% | 3 | - | 48747410 | 48747509 | 100 |

Figure 4.5: BLAT output with defaults for query sequence pair1

| BLAT Search Results |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ACTIONS QUERY | SCORE | START | END | QSIZE | IDENTITY | CHRO | STR | AND START | END | SPAN |
| browser details YourSeq | 96 | 1 | 100 | 100 | 98.0\% | 4 | + | 191129790 | 191129889 | 100 |
| browser details YourSeq | 94 | 1 | 100 | 100 | 97.0\% | X | - | 119493641 | 119493740 | 100 |
| browser details YourSeq | 94 | 1 | 100 | 100 | 97.0\% | X | - | 7536075 | 7536174 | 100 |
| browser details YourSeq | 94 | 1 | 100 | 100 | 97.0\% | 9 | - | 118642009 | 118642108 | 100 |
| browser details YourSeq | 94 | 1 | 100 | 100 | 97.0\% | 9 | - | 111899086 | 111899185 | 100 |
| browser details YourSeq | 94 | 1 | 100 | 100 | $97.0 \%$ | 9 | - | 94028941 | 94029040 | 100 |
| browser details YourSeq | 94 | 1 | 100 | 100 | 97.0\% | 9 | - | 94632200 | 94632299 | 100 |
| browser details YourSeq | 94 | 1 | 100 | 100 | $97.0 \%$ | 8 | - | 145357732 | 145357831 | 100 |
| browser details YourSeq | 94 | 1 | 100 | 100 | 97.0\% | 8 | - | 129199276 | 129199375 | 100 |

Figure 4.6: BLAT output with defaults for query sequence pair2
The BLAT output was comparable with BLAST and took less time even with a web interface. BLAT had 200+ entries for each query pair and showed significant identity with chromosome 12 for sequence pair 1 (Figure 1 and 2) and chromosome 4 for sequence pair 2. More importantly, BLAST did not show any output if the option for filtering in repeat regions was turned on. BLAT does not have such an option and gives all possible hits.

### 4.3.2 BLAT alignment method used in tool

We used the FASTA output from the read extraction as the input for BLAT tool. BLAT can be run as server-client interface where the user creates a server (gfServer) that
keeps the genome index in memory while the client (gfClient) is used to input query sequences that are sent to the gfServer. The server-client program has also been combined into a single standalone program such that the database and query are sent as arguments to the program. Instructions on setting up the client and options used by BLAT are provided in Appendix B. A perl script was written that allows the user to set up the server-client interface in a single step. A pre-requisite for this was that the 2-bit file for the genome needs to be made using a BLAT-faToTwoBit program. This program converts a FASTA file into ". 2 bit" format that is then used by gfServer to create the reference genome index file. In order to speed up the process, the reference genome was split by chromosome, and the query was run against each database in parallel.

We used BLAT-version 35.1 to run the queries with minimum tile match of 4 and minimum identity of $95 \%$. These queries were run in parallel, one for each chromosome. Output from BLAT is presented in the psl format (Appendix B, BLAT specification) which has the fields as defined in Table 1.1. The output from each chromosome mapping was finally merged together and sorted by 'Identifier' (field 10), 'Chromosome' (field 14) and 'Start position' (field 16) to get the sorted psl output file.

Table 4.7: BLAT output file format: psl

| Field \# | Field <br> (Type)$\quad$ Names | Description |
| :---: | :---: | :---: |
| 1 | Matches (int unsigned) | Number of bases that match that aren't repeats |
| 2 | Mismatches (int unsigned) | Number of bases that don't match |
| 3 | repMatches (int unsigned) | Number of bases that match but are part of repeats |
| 4 | nCount (int unsigned) | Number of 'N' bases |
| 5 | qNumInsert (int unsigned) | Number of inserts in query |
| 6 | qBaselnsert (int unsigned) | Number of bases inserted in query |
| 7 | tNumInsert (int unsigned) | Number of inserts in target |
| 8 | tBasel nsert (int unsigned) | Number of bases inserted in target |
| 9 | Strand ( $\operatorname{char}(2)$ ) | + or - for query strand optionally followed by + or for target strand |
| 10 | qName (varchar(255)) | Query sequence name |
| 11 | qSize <br> (int unsigned) | Query sequence size |
| 12 | qStart (int unsigned) | Alignment start position in query |
| 13 | qEnd (int unsigned) | Alignment end position in query |
| 14 | tName (varchar(255)) | Target sequence name |
| 15 | tSize (int unsigned) | Target sequence size |
| 16 | tStart (int unsigned) | Alignment start position in target |
| 17 | tEnd (int unsigned) | Alignment end position in target |
| 18 | blockCount (int unsigned) | Number of blocks in alignment. A block contains no gaps. |
| 19 | blockSizes (longblob) | Size of each block in a comma separated list |
| 20 | qStarts (longblob) | Start of each block in query in a comma separated list |
| 21 | tStarts (longblob) | Start of each block in target in a comma separated list |

### 4.4 De-Duplication Algorithm

The de-duplication algorithm essentially checks for rows with the same 'Identifier' and 'Chromosome' and bins it into buckets of 1000 basepairs windows. The output for this algorithm produces an SML file, an in-house defined format (Table 4.8).

Table 4.8: De-Duplication Algorithm

| Assumption: The input file is in psl format which has been sorted by 'Read identifier', 'Chromosome' and 'Start Position' | Creating sml file from psl file |
| :---: | :---: |
| ```Input Arguments for file name; if (file is in sorted psl format) { Create output file with sml extension;} Else {output error to screen}``` | Creates output file with .sml extension |
| While \{ <br> Read each line of psl file; Split into fields based on delimiter; Set fields 9:Identifier <br> 13: Chromosome <br> 15: Start position <br> 16: End position <br> 17: BlockCount as variables; Define variables current-range and previousrange using GetRange() function return value | Sets variables and defines new variables currentrange and previous-range |
| ```If (Prev Identifier not equal Curr Identifier) {Insert Read(to output file)} Else if (Prev Chr not equal Curr Chr) {if (BlockSize==1) {Insert Read(to output file)} Else if (current-range not equals previous-range) {if (BlockSize==1) {Insert Read(to output file)} } } } END OF WHILE LOOP``` Close all files | Bins data in 1000 basepair regions. If identifiers and chromosomes are same, compress the rows into 1000 basepair bins, finally filtered on BlockSize; deduplication of rows |
| ```Sub GetRange{ rangeSpan=1000; midpoint=(StartPosition+EndPosition)/2 if (midpoint < rangeSpan) {``` | GetRange function defines the 1000 base-pair blocks into which the consecutive rows |


| ```grange = 1; } else { grange = int(midpoint/rangespan) + 1; } return grange ; }``` | $\begin{array}{lr}\text { with } & \text { same } \\ \text { identifier } & \text { and }\end{array}$ chromosome are compressed i.e. de-duplicated. Mid-point of length of mapping region is used to check for range span. The return value of the function is the offset for each bin starting from 1 for 1-1000 bp bin |
| :---: | :---: |
| ```Sub InsertRead{ print 'Identifier’, 'Chromosome', 'Start position', 'Start position', 'current-range' }``` | Output format for the sml file defined |

Table 4.9: Example sml file from de-duplication algorithm

| Identifier | Chromo some | Start Position | End Position | Offset for Range |
| :---: | :---: | :---: | :---: | :---: |
| $\begin{aligned} & \text { 4_103112562_103113055 } \\ & -5: 1: 0 \_2: 0: 0 \text { dedbba<RID>1 } \end{aligned}$ | 4 | 950 | 1010 | 1 |
| $\begin{aligned} & \text { 4_14019666_14020194 } \\ & \overline{6}: 0: 0 \_1: 0: 0 \_17 \mathrm{f} 93 \mathrm{~b} 6<\mathrm{RID}>1 \end{aligned}$ | 8 | 999 | 1010 | 2 |
| $\begin{aligned} & \text { 4_4086095_4086615 } \\ & 3: 0: 0 \_2: 0: 0 \_ \text {a6acbd<RID>1 } \end{aligned}$ | 4 | 4086515 | 4086615 | 4087 |
| $\begin{aligned} & \text { 4_4086106_4086642 } \\ & \overline{3}: 0: 0 \_2: 0: 0 \_5 c d f a 1<\text { RID }>1 \end{aligned}$ | 4 | 4086105 | 4086205 | 4087 |
| $\begin{gathered} 4-4086106=4086642 \\ 3: 0: 0 \quad 2: 0: 0 \quad 5 \mathrm{cdfal}<\mathrm{RID}>2 \end{gathered}$ | 4 | 4086542 | 4086642 | 4087 |

### 4.5 Create_Matrix Algorithm

The SML file created from de-duplication algorithm serves as input file for matrix creation. The SML file essentially contains all the possible mapping locations for the queries in PSL file on the reference genome based on BLAT mapping algorithm in 1000 base pair windows.

The Create_Matrix algorithm creates all possible combinations of the read pair (RID1 and RID2), such that each row in output file is the location where readpair1 maps and the corresponding location where readpair2 maps on the human genome. The algorithm is defined in Table 4.10. The Create_Matrix algorithm essentially defines how to capture the NAHR regions and can be represented in the Figure 4.7a, 4.7b. Output is a MAT file with an example in Table 4.11.

Table 4.10: Create_Matrix Algorithm

| Assumption: The input file is in sml format and is sorted by 'Read Identifier' and 'Chromosome' | Creating mat file from sml file |
| :---: | :---: |
| ```Input Arguments for file name; if (file is in sorted sml format) { Create output files with mat extension; } Else {output error to screen}``` | Creates output files with .mat extension |
| While \{ <br> Read each line of psl file; <br> Split into fields based on delimiter; <br> Set fields 1:Identifier <br> 2: Chromosome <br> 3: Start position <br> 4: End position <br> 5: Range offset; | Sets variables |
| ```If (previous Identifier==current Identifier) {Insert Read(to output file)} Else { Create_Matrix() Insert Read() } If first line of file {Insert Read(to output file} } END OF WHILE LOOP``` |  |


| Close all files <br> Do UNIX sort on Read1,Read2, Range Offset1, Range Offset2 |  |
| :---: | :---: |
| ```Sub Create_Matrix{ Foreach read1 { Foreach read2 { If (Chromosome and Range offset is same) {do not print} Else { If (swapFunction==False) {print to outfile read1, read2} Else {print to outfile read2, read1} } }``` | Matrix is created based on conditions; if both read pairs map to the same region we do not print the output as these do not suggest a translocation. Else if read1 and read2 go on different locations then print to outfile. |
| ```For (read1) { For (next read1) { If (Chromosome and Range offset is same) {do not print} Else { If (swapFunction==False) {print to outfile read1, next read1} Else {print to outfile nextread1, read1} }``` | Second condition checks for readl going to different locations i.e. the same read mapping to different locations and creating a matrix of these combinations. Printing of values is such that Chromosome field1 is always less than Chromosome field2 and Chromosome $X$ is always less than |
| $\}_{\}}^{\}}$ | Chromosome Y |
| ```For (read2) { For (next read2) { If (Chromosome and Range offset is same) {do not print} Else { If (swapFunction==False) {print to outfile read2, next read2} Else {print to outfile nextread2, read2} } }``` | Third condition checks for read2 going to different locations i.e. the same read mapping to different locations and creating a matrix of these combinations. Printing of values is such that Chromosome field1 is always less than Chromosome field2 and Chromosome $X$ is always less than Chromosome Y . |


| $\}^{\}}$ |  |
| :---: | :---: |
| ```Sub InsertRead { If (Identifier==Read1) { write array for Read1;Chromosome and Offset} If (Identifier==Read2) { write array for Read2;Chromosome and Offset} }``` | Creates array for read1 and read 2 and pushes each read into the respective array based on Identifier |
| ```Sub SwapFunction\{ Set default SwapFunction return-value==No; Create arg1: Chromosome read1 arg2: Chromosome read2 arg3: Offset read1 arg4: Offset read2 If (Chromosome read \(1==Y\) and Chromosome read2==X ) \{ swap return value \(==\) Yes \(\}\) Else If (Chromosomes read1 and 2 are integers and Chromosome read1>Chromosome read2) \{ swap return value==Yes \} Return swap return-value; \}``` | This function returns a value of "Yes" if Chromosome on field1 is greater than Chromosome in field2. Y is greater than X as defined. |

Table 4.11: Example 'mat' file from create_matrix algorithm

| Chromosome1 | Offset1 | Chromosome2 | Offset2 |
| :--- | :--- | :--- | :--- |
| 1 | 1001 | 4 | 2021 |
| 1 | 1004 | 4 | 5036 |
| 1 | 2385 | 8 | 6236 |
| 1 | 5321 | 8 | 8288 |
| 4 | 4087 | 8 | 6974 |
| 4 | 4087 | 8 | 6974 |
| 4 | 4087 | 8 | 6974 |


$\times$
Not printed since both reads go to the same chromosome

Figure 4.7a: Create_Matrix Algorithm: Matrix creation with all possible combinations of read1 and read2 excluding reads where both go to same chromosome.


Figure 4.7b: Create_Matrix Algorithm: Swapping reads such that lower number chromosome is always on the left which makes the counting process efficient.

Non-allelic homologous recombination regions are defined by their sequence identity ( $>95 \%$ ) with different regions of the genome and greater than 5 kb in size such that these form substrates for anomalous pairing during repair. BLAT identifies all possible regions such that read-pair1 maps to one region and read-pair2 maps to another region which will therefore identify potential NAHR regions. The Create_Matrix algorithm, in creating the matrix, identifies potential NAHR pairs and therefore potential translocation breakpoints. Further support is added if read1 also maps to both locations on different chromosomes identified as potential breakpoint regions (e.g. Read1 mapping to chromosome 4 and chromosome 8) suggesting that these regions could be homologous and share significant identity. The purpose of
swapping the reads is to sort the reads for easier counting. This finally gives the number of reads supporting a translocation identified by the matrix.

### 4.6 Write_Count Algorithm

Output from matrix creation is a sorted MAT file and is the input for the write_count algorithm. The output for the write_count algorithm is a CHR file as shown in example Table 4.12.

Table 4.12: Write_count Algorithm

| Assumption: The input file is in mat format and is sorted by field1, field3, field2 and field4 in that order, i.e. Chromosome1, Chromosome2, Offset1, Offset2 | Creating chr file from mat file |
| :---: | :---: |
| ```Input Arguments for file name; if (file is in sorted .mat format) { Create output files with .chr extension; } Else {output error to screen}``` | Creates output files with .chr extension |
| ```While (Read each line of mat file) { If (previous record==current record) {Count=1; Increment count} Else {WriteCount(to output file)} } END OF WHILE LOOP``` Close all files | Counts number of records which are same |
| Sub WriteCount \{ Print each unique row with count; \} | Creates '.chr' output file format |

Table 4.13: Example 'chr' file from write_count algorithm

| Chromosome1 | Offset1 | Chromosome2 | Offset2 | Count |
| :--- | :--- | :--- | :--- | :--- |
| 1 | 1001 | 4 | 2021 | 1 |
| 1 | 1004 | 4 | 5036 | 1 |
| 1 | 2385 | 8 | 6236 | 1 |
| 1 | 5321 | 8 | 8288 | 1 |
| 4 | 4087 | 8 | 6974 | 3 |

### 4.7 Get_HiC-Score Algorithm

The last step in the process was to define the probability of calling a translocation as not false. We did not used distribution of insert size within the dataset to define this probability as is done by other variant calling programs (SVDetect, BreakDancer, PEMer). The cancer genome is highly heterogeneous with various clonal populations showing different types of variation signatures. Even with a 30X coverage of sequencing experiment, the likelihood of picking a variant in the heterogeneous cancer sample cannot be defined based on the distribution of reads in the current sample. We therefore wanted to capture as many variants as we could define using our algorithm and assign the probability based on an informative prior. This was determined using Hi-C experiments (Lieberman-Aiden et al., 2009) which takes into account the three dimensional positioning of the genomes within a cell and defines the probability of two regions of the genome interacting based on their physical proximity to each other, as determined experimentally. For the purpose of this analysis we used the database created by Lieberman-Aiden's groups at BROAD/MIT. ${ }^{13}$ The datasets contain Pearson's correlation coefficient for each combination of chromosome in one million base-pair (1Mbps) windows. This is the highest level of resolution for these experiments, and we had our variants defined in 1000 basepair windows. We had to accept this as a limitation of the study.

[^8]Table 4.14: Get_HiC-Score Algorithm

| Assumption: The input file is in chr format and is sorted by field1(chr1), field3(chr2), field2(Offset1) and field4(Offset2) in that order | Creating FASTA file from BAM output of mapping tool |
| :---: | :---: |
| ```Input Arguments for file name, path to Hi-C file; if (file is in sorted .chr format) { Create output files with .score extension; } Else {output error to screen}``` | Creates output files with .score extension |
| While (Read each line of chr file) <br> \{ <br> Split into fields based on delimiter; <br> Set fields 1: Chromosome1 <br> 2: Offset1 <br> 3: Chromosome2 <br> 4: Offset2 <br> 5: Chromosome Count <br> 6: Current $\mathrm{Hi}-\mathrm{C}$ filename <br> Create $\mathrm{Hi}-\mathrm{C}$ filename by dynamically setting values of chromosome1 and chromosome2 on $\mathrm{Hi}-\mathrm{C}$ filename template |  |
| ```If (current Hi-C filename not equals previous Hi-C filename) { Open Hi-C file () } row-midrange= CalcMidPosition(offset1) column-midrange=CalcMidPosition(offset2) rowposition=GetFilePosition(row-midrange); colposition=GetFilePosition(column-midrange); score = GetHi-CScore() CalcRange(Offset1,Offset2) If (score>0) { Print to outfile chr1,start1,end1,chr2,start2,end2 numberOfReads, scorefrom Hi-C } } END OF WHILE LOOP Close all files``` | Open each Hi-C file only once and calculate midrange of row and column offsets. Use these offsets to get the exact column and row position in $\mathrm{Hi}-\mathrm{C}$ file, open $\mathrm{Hi}-\mathrm{C}$ file and extract the score for $x$ row and y-cloumn. <br> Only Hi -C scores with Pearson's probability correlation greater than zero printed to output |
| Sub Open Hi-C file\{ <br> Concatenates $\mathrm{Hi}-\mathrm{C}$ filepath variable to dynamically created current $\mathrm{Hi}-\mathrm{C}$ filename and open corresponding file. <br> Puts entire file into @filearray; \} | Opens $\mathrm{Hi}-\mathrm{C}$ file bases on chromosomel and chromosome2 values and their offsets. Puts the entire file into file array which is scanned in the GetHi-CScore function for getting score. |

$\left.\begin{array}{|l|l|}\hline & \begin{array}{l}\text { Each Hi-C file is opened } \\ \text { only once since chr' file } \\ \text { is sorted by } \\ \text { chromosomes and }\end{array} \\ \text { therefore all } \\ \text { corresponding for Hi-C } \\ \text { scores for each } \\ \text { combination of row and } \\ \text { column offset values } \\ \text { can be obtained in one } \\ \text { open file function call. }\end{array}\right\}$

| r2start = r2end-(chrspan-1); |  |
| :--- | :--- |
| $\}$ |  |

The output from this program is as represented in Table 4.15 that is presented as the final output to the user. Further, the algorithm filtered out any $\mathrm{Hi}-\mathrm{C}$ scores below zero and only reported translocations with positive $\mathrm{Hi}-\mathrm{C}$ probability scores.

Table 4.15: Example 'score' file from get_HiC-score algorithm

| Chromo <br> some1 | Start | End | Chromo <br> some2 | Start | End | No. <br> of <br> reads | Hi-C <br> Score |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| chrX | 0 | 999 | chrX | 1000 | 1999 | 2 | 1 |
| chrX | 4000 | 4999 | chrX | 9000 | 9999 | 3 | 1 |
| chr1 | 5170000 | 5170999 | chrX | 4620000 | 4620999 | 2 | 0.050044 |
| chr1 | 5193000 | 5193999 | $\operatorname{chrX}$ | 17311000 | 17311999 | 2 | 0.018999 |
| chr1 | 5236000 | 5236999 | chrX | 2850000 | 2850999 | 8 | 0.061457 |
| chr1 | 16000 | 16999 | chr20 | 45680000 | 45680999 | 3 | 0.089301 |
| chr1 | 23000 | 23999 | chr20 | 63000 | 63999 | 2 | 0.359193 |

### 4.8 Proof of Concept

BLAT provides various input options to the user that can change the output significantly based on the user requirements. In our case, since we were looking for regions which share $95 \%$ sequence identity and can therefore give more than one significant hit, we wanted to keep the minldentity option as $95 \%$. As a proof of concept we created a single translocation in chromosomes 4 and 8 (t4;8)(p16.3;p23.1) with breakpoint of chromosome 4 at chr4: 4086365 and on chromosome 8 at chr8: 6973436 (Figure 4.8).


Figure 4.8: Translocation $(4 ; 8)$; derivative chromosomes created using (Hiller, Bradtke, Balz, \& Rieder, 2005).

The simulation data had derivative chromosomes 4 and 8 as well as normal chromosome 4 and 8 . This was created using the perl script provided by Dr Hayes for inserting manufactured translocations in the normal human genome. This mini FASTA file containing normal chromosome 4 and 8 and the derivative chromosome 4 and 8 was mapped to reference genome hg19 (UCSC) using BWA-0.5.9. The BAM file from this alignment was used as input for extraction of anomalous reads (Section 4.1). Since this simulated data only had chromosomes 4 and 8 , the database file for

BLAT was also only created for chromosome 4 and 8. BLAT was run on FASTA output file obtained from the extraction algorithm applied to the BWA-alignment BAM file.

There were 2,441,572 anomalous reads identified from the extraction step. We used two BLAT parameters, minMatch and minldentity to test for BLAT output. These are defined in the BLAT manual (Appendix B) as below:
-tileSize $=N$ sets the size of match that triggers an alignment, usually between 8 and 12. Default is 11 for DNA and 5 for protein.
-minMatch=N Sets the number of tile matches. Usually set from 2 to 4. Default is 2 for nucleotide, 1 for protein.
-minldentity $=\mathrm{N}$ Sets minimum sequence identity (in percent).Default is 90 for nucleotide searches, 25 for protein or translated protein searches.

We kept the default tile size of 11 for this analysis and adjusted the other two parameters. The output from BLAT is summarized in Table 4.16.

Table 4.16: BLAT parameter adjustment results

| Parameter | Number of <br> alignments | Number of <br> translocation <br> reported | Number of <br> reads <br> supporting |
| :--- | :--- | :--- | :--- |
| Defaults: <br> minIdentity=90 <br> minMatch=2 | 15111989 | 339887 | 3 |
| minIdentity=95 <br> minMatch=2 | 7482820 | 141527 | 2 |
| minIdentity=99 <br> minMatch=2 | 1073276 | 8733 | 0 |
| minIdentity=95 <br> minMatch=4 | 6620511 | 118282 | 0 |
| minIdentity=99 <br> minMatch=4 | 934711 | 8166 | 0 |

Testing the algorithm with a simulated single translocation $(t(4 ; 8))$ dataset showed that adjusting BLAT parameters to $95 \%$ identity and minimum tile match of two, captured the positive translocations while reducing false positive by almost a third (Table 4.16). Therefore these BLAT options were used in the final simulated dataset analysis. When tightening the alignment option to report $99 \%$ identity many false positives are reduced though at the cost of true positives. Since the aim was to detect these events which are most likely to be missed by BWA alignment tool and since $95 \%$ identity is the criteria for a non-allelic homologous recombination substrate, we used this option at the cost of reporting many false positives.

## 5. SIMULATED DATA ANALYSIS

### 5.1 Creating Simulated Dataset

The idea for creating a simulated dataset was to include known documented translocations that were biologically plausible and be able to test the algorithm on this dataset. An attempt to define drivers of translocation based on genomic architecture was done by Ou et al., (Ou et al., 2011) whereby they demonstrated recurrent translocation driven by non-allelic homologous recombination in unrelated families. They also mapped the regions which could be involved in potential translocation computationally using low-copy repeat regions in the genome which shares $>94 \%$ sequence identity and more than 5 kb in length. They were able to experimentally verify the computationally predicted translocations for the three regions identified in Table 5.1.

Burrow et al. analyzed recurrent translocation in cancers from various databases and tried to define the characteristics of these translocations. They found that over 50\% of the recurrent translocations mapped to fragile sites, defined as regions on the genome that show multiple gaps (Burrow, Williams, Pierce, \& Wang, 2009). We tried to derive our list of translocations from this comprehensive list (Appendix C) that could be near NAHR regions and found seven listed in Table 5.1. This list of translocation breakpoints was used to create derivative chromosomes using the translocation perl script mentioned in section 4.8.

Table 5.1: Translocation list used in creating simulated dataset

| Translocation from Burrow et al paper | Breakpoint First chr. | Breakpoint Second chr. | Gene First chr | Gene Second chr. |
| :---: | :---: | :---: | :---: | :---: |
| t(1;22)(1q21;22q11) | 142749690 | 19819306 | BCL9 | IGL@ |
| t(12;13)(12q14;13q13) | 50787223 | 40914700 | HMGA2 | LHFP |
| t(19;22)(19q13;22q11) | 63789868 | 49571663 | IGL@ | BCL3 |
| t(12;16)(12p13;16p13) | 8270437 | 5069858 | LAG3 | MYH11 |
| t(16;21)(16q24.3; 21q22.12) | 88815835 | 46924874 | RUNX1 | CBFA2T3 |
| t(9;14)(9p21;14q11) | 27286005 | 42868270 | TRA@ | CDKN2A |
| t(7;11)(7q34;11p15) | 128044540 | 124564 | TRB@ | LMO1 |
| Translocation from Ou et al. paper (Ou et al., 2011) | Breakpoint First chr. | Breakpoint Second chr. | Gene First chr | Gene Second chr. |
| t(4;8)(4p16.3;8p23.1) | 4088911 | 6992273 | NA | NA |
| t(4:11)(4p16.2;11p15.4) | 3852863 | 3569449 | NA | NA |
| t(8;12)(8p23.1;12p13.31) | 6992273 | 8367384 | NA | NA |

This simulated dataset included normal chromosomes and the derivative translocated chromosomes. The FASTA file was used to create simulated data using wgsim program that is a part of SAMTOOLS suite. A description of the options used for this program is presented in Table 1.1.

Table 5.2: Simulation read creating program- 'wgsim’ options

| Options | Type | Description |
| :--- | :--- | :--- |
| $e$ | FLOAT | base error rate [0.020] |
| d | INT | outer distance between the two ends [500] |
| s | INT | standard deviation [50] |
| N | INT | number of read pairs [1000000] |
| 1 | INT | Iength of the first read [70] |
| 2 | INT | Iength of the second read [70] |
| r | FLOAT | rate of mutations [0.0010] |
| R | FLOAT | fraction of indels [0.10] |
| X | FLOAT | probability an indel is extended [0.30] |
| c | NA | generate reads in color space (SOLiD read |
| C | NA | show mismatch info in comment rather than |
| h | NA | haplotype mode |

At the very least, it requires input FASTA file (translocation FASTA file) and length of the reads. In order to get 30X coverage of the genome we needed to define the number of read pairs needed ( -N , Table 1.1). Coverage is calculated using equation 1 below:

$$
\text { (1) } \text { Coverage }=\frac{\text { Length of read } \times \text { Number of reads }}{\text { Haploid Genome Length }}
$$

Thus for 30X coverage and read length of 100 basepairs, we calculated the number of read required for the haploid genome of length $3 \times 10^{9}$ to be 900 million reads. Output from this program gives two FASTQ files, one each for a paired end library. Further, a random subsample of this dataset was produced to simulate 15 X coverage. BWA mapping tool was used to produce the alignment BAM files. The BAM file was then run against current tools, namely SVDetect-0.7f and BreakDancer-1.1. The process flow for our algorithm was as described in Figure 5.1.


Figure 5.1. Process flow of novel algorithm

### 5.2 Simulated Data Analysis Results

SVDetect could correctly identify four out of the ten translocations created (4/10) and BreakDancer detected six out of ten $(6 / 10)$ briefly summarized in Table 5.3.

Table 5.3: Comparison of current tools with novel algorithm

| Translocation from Burrow et al paper | SVDetect | BreakDancer | Novel Algorithm |
| :---: | :---: | :---: | :---: |
| t(1;22)(1q21;22q11) | Detected | No | Detected |
| t(12;13)(12q14;13q13) | Detected | Detected | Detected |
| t(19;22)(19q13;22q11) | No | No | Detected |
| t(12;16)(12p13;16p13) | Detected | Detected | Detected |
| t(16;21)(16q24.3;21q22.12) | No | No | Detected |
| t(9;14)(9p21;14q11) | Detected | Detected | Detected |
| t(7;11)(7q34;11p15) | No | No | Detected |
| Translocation from Ou et al. paper (Ou et al., 2011) | SVDetect | BreakDancer | Novel Algorithm |
| t(4;8)(4p16.3;8p23.1) | No | Detected | Detected |
| t(4:11)(4p16.2;11p15.4) | No | Detected | Detected |
| t(8;12)(8p23.1;12p13.31) | No | Detected | Detected |

The novel algorithm was able to detect all the translocations at a 1000 base pair resolution. The number of reads supporting a translocation varied from 3 to 50. The three translocations that were not detected by either SVDetect or BreakDancer, $\mathrm{t}(19 ; 22), \mathrm{t}(16 ; 21)$ and $\mathrm{t}(7 ; 11)$, all had less than 10 reads supporting the translocation in the novel algorithm output. The score file output is represented in Table 5.4.

Table 5.4: Final score file output of novel algorithm 30X coverage

| Chr_1 | Start | End | Chr_2 | Start | End | No_ <br> of <br> reads | Pearsons <br> Corr- <br> elation <br> coeff <br> icient |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| chr1 | 142749000 | 142749999 | chr22 | 19819000 | 19819999 | 3 | 0.01779 |
| chr12 | 50786000 | 50786999 | chr13 | 40914000 | 40914999 | 3 | 0.00739 |
| chr12 | 50787000 | 50787999 | chr13 | 40914000 | 40914999 | 17 | 0.00739 |
| chr19 | 63789000 | 63789999 | chr22 | 49571000 | 49571999 | 4 | 0.30154 |
| chr12 | 8270000 | 8270999 | chr16 | 5069000 | 5069999 | 29 | 0.28795 |
| chr16 | 88815000 | 88815999 | chr21 | 46924000 | 46924999 | 9 | 0.42250 |
| chr9 | 42867000 | 42867999 | chr14 | 27285000 | 27285999 | 15 | 0.25198 |
| chr9 | 42868000 | 42868999 | chr14 | 27285000 | 27285999 | 23 | 0.25198 |
| chr9 | 42868000 | 42868999 | chr14 | 27286000 | 27286999 | 23 | 0.25198 |
| chr7 | 128044000 | 128044999 | chr11 | 124000 | 124999 | 6 | 0.25539 |
| chr4 | 4088000 | 4088999 | chr8 | 6991000 | 6991999 | 2 | 0.21970 |
| chr4 | 4088000 | 4088999 | chr8 | 6992000 | 6992999 | 15 | 0.21970 |
| chr4 | 3852000 | 3852999 | chr11 | 3569000 | 3569999 | 23 | 0.1985 |
| chr8 | 6991000 | 6991999 | chr12 | 8367000 | 8367999 | 7 | 0.24114 |
| chr8 | 6992000 | 6992999 | chr12 | 8367000 | 8367999 | 51 | 0.24114 |

SVDetect, although correctly identified chromosome 1 breakpoint for translocation $\mathrm{t}(1 ; 22)$ did not identify chromosome 2 breakpoint as precisely and had very few reads supporting this translocation (Table 5.5). BreakDancer did a very good job at breakpoint resolution but could only detect six out of ten (Table 5.6) while SVDetect picked four (Table 5.5).

Table 5.5: SVDetect output (trimmed): simulated data 30X coverage

| SV_type | BAL_ <br> type | chrom <br> o <br> some1 | chrom <br> o <br> some2 | no <br> _pai <br> rs | final <br> sco <br> re | break <br> point1 | break <br> point2 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| INV_ <br> TRANSLOC | UNBAL | chr1 | chr22 | 3 | 1 | $14274984-$ <br> 142750363 | $19960253-$ <br> 19960750 |
| INV- <br> TRANSLOC | UNBAL | chr1 | chr22 | 1 | 1 | $14274900-$ <br> 142749572 | $19958766-$ <br> 19959329 |
| TRANSLOC | UNBAL | chr1 | chr22 | 1 | 1 | $14274981-$ <br> 142750378 | $22001986-$ <br> 22002549 |
| INV <br> TRANSLOC | UNBAL | chr12 | chr13 | 1 | 1 | $50786416-$ <br> 50786979 | $40903686-$ <br> 40904249 |
| TRANSLOC | UNBAL | chr12 | chr16 | 24 | 0.83 | $8270917-$ <br> 8271106 | $5069215-$ <br> 5069396 |
| INV_- <br> TRANSLOC | UNBAL | chr9 | chr14 | 36 | 0.88 | $42867615-$ <br> 42867796 | $27285344-$ <br> 27285586 |

Table 5.6: BreakDancer output: simulated data 30X coverage

| Chrom <br> osome1 | Pos1 | Chrom <br> osome2 | Pos2 | Type | Size | Score | Number <br> of Reads |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| chr4 | 3852787 | chr11 | 3569702 | CTX | -499 | 99 | 4 |
| chr4 | 3853180 | chr11 | 3569260 | CTX | -499 | 99 | 6 |
| chr4 | 4088565 | chr8 | 6992324 | CTX | -499 | 99 | 2 |
| chr4 | 4088649 | chr8 | 6992415 | CTX | -499 | 99 | 2 |
| chr8 | 6992535 | chr12 | 8367239 | CTX | -499 | 99 | 2 |
| chr8 | 6992629 | chr12 | 8367239 | CTX | -499 | 99 | 3 |
| chr8 | 6991970 | chr12 | 8367520 | CTX | -499 | 99 | 3 |
| chr8 | 6992195 | chr12 | 8367520 | CTX | -499 | 99 | 3 |
| chr8 | 6992195 | chr12 | 8367640 | CTX | -499 | 99 | 5 |
| chr9 | 42868057 | chr14 | 27285682 | CTX | -499 | 99 | 19 |
| chr12 | 8270277 | chr16 | 5069881 | CTX | -499 | 99 | 3 |
| chr12 | 8270277 | chr16 | 5069996 | CTX | -499 | 99 | 5 |
| chr12 | 8270277 | chr16 | 5070133 | CTX | -499 | 95 | 2 |
| chr12 | 8270357 | chr16 | 5070133 | CTX | -499 | 99 | 3 |
| chr12 | 8270471 | chr16 | 5069481 | CTX | -499 | 99 | 3 |
| chr12 | 8270742 | chr16 | 5069481 | CTX | -499 | 98 | 2 |
| chr12 | 8270742 | chr16 | 5069585 | CTX | -499 | 99 | 11 |
| chr12 | 50787289 | chr13 | 40914358 | CTX | -499 | 99 | 3 |
| chr12 | 50787452 | chr13 | 40914358 | CTX | -499 | 99 | 8 |

As described, cancer genome complexity is overburdened further by factors like sample collection, tumor heterogeneity (Ulahannan, Kovac, Mulholland, Cazier, \& Tomlinson, 2013), and platform specific issues like AT-rich and GC-rich bias in the Illumina platform (Metzker, 2010). Even with a 30X coverage of the genome and a $100 \%$ representation of the alternative allele, the tools did not detect $40 \%$ to $60 \%$ of the simulated translocations. For the next step the dataset was randomly sampled so that $50 \%$ of the reads were picked to do analysis on a 15 X coverage dataset. This is more realistic for current whole-genome sequencing strategies of cancer. 30X or greater coverage is usually reserved for exome-sequencing mainly due to cost constraints. Even if a cancer genome is sequenced at 30 X or greater, the inherent intercellular heterogeneity coupled with the aneuploidy that typifies most tumors means that a given translocation may be represented by even fewer supporting reads for the regions of interest.

Even with 15 X coverage, our novel algorithm tool was able to detect all ten simulated translocations (Table 5.7) albeit with fewer reads supporting the translocations (compare with Table 5.4).

Table 5.7: Final score file output of novel algorithm 15X coverage

| Chr_1 | Start | End | Chr_2 | Start | End | No_ of reads | Pearsons Correlation coeff icient |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| chr1 | 142749000 | 142749999 | chr22 | 19819000 | 19819999 | 2 | 0.01779 |
| chr12 | 50786000 | 50786999 | chr13 | 40914000 | 40914999 | 2 | 0.00739 |
| chr12 | 50787000 | 50787999 | chr13 | 40914000 | 40914999 | 7 | 0.00739 |
| chr19 | 63789000 | 63789999 | chr22 | 49571000 | 49571999 | 3 | 0.30154 |
| chr12 | 8270000 | 8270999 | chr16 | 5069000 | 5069999 | 14 | 0.28795 |
| chr16 | 88815000 | 88815999 | chr21 | 46924000 | 46924999 | 5 | 0.42250 |
| chr9 | 42867000 | 42867999 | chr14 | 27285000 | 27285999 | 10 | 0.25198 |
| chr9 | 42868000 | 42868999 | chr14 | 27285000 | 27285999 | 8 | 0.25198 |
| chr9 | 42868000 | 42868999 | chr14 | 27286000 | 27286999 | 14 | 0.25198 |
| chr7 | 128044000 | 128044999 | chr11 | 124000 | 124999 | 3 | 0.25539 |
| chr4 | 4088000 | 4088999 | chr8 | 6992000 | 6992999 | 11 | 0.21970 |
| chr4 | 3852000 | 3852999 | chr11 | 3569000 | 3569999 | 9 | 0.1985 |
| chr8 | 6991000 | 6991999 | chr12 | 8367000 | 8367999 | 4 | 0.24114 |
| chr8 | 6992000 | 6992999 | chr12 | 8367000 | 8367999 | 28 | 0.24114 |

BreakDancer was only able to detect four (4/10) translocations with 15X coverage (Table 5.8).

Table 5.8: BreakDancer output: simulated data 15X coverage

| Chrom <br> osome1 | Pos1 | Chrom <br> osome2 | Pos2 | Type | Size | Score | Number <br> of Reads |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| chr4 | 3852745 | chr11 | 3569702 | CTX | -499 | 99 | 3 |
| chr4 | 3853180 | chr11 | 3569294 | CTX | -499 | 99 | 4 |
| chr8 | 6992049 | chr12 | 8367520 | CTX | -499 | 99 | 2 |
| chr9 | 42867966 | chr14 | 27285748 | CTX | -499 | 99 | 5 |
| chr9 | 42868047 | chr14 | 27285748 | CTX | -499 | 99 | 2 |
| chr12 | 8270471 | chr16 | 5069481 | CTX | -499 | 99 | 2 |
| chr12 | 8270566 | chr16 | 5069585 | CTX | -499 | 99 | 2 |

SVDetect was only able to detect two (2/10) variants (Table 5.9).

Table 5.9: SVDetect output (trimmed): simulated data 15X coverage

| SV_type | BAL_type | Chr_1 | Chr_ <br> 2 | nb <br> _pair <br> s | final <br> sco <br> re | break <br> point1 | break <br> point2 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| INV_- <br> TRANSLOC | UNBAL | chr1 | chr22 | 1 | 1 | $142749572-$ <br> 142749671 | $19959329-$ <br> 19959428 |
| TRANSLOC | UNBAL | chr12 | chr13 | 4 | 0.8 | 50786565 <br> -50786795 | $40915026-$ <br> 40915369 |
| TRANSLOC | UNBAL | chr12 | chr13 | 4 | 0.8 | 50787552 <br> -50788028 | $40913925-$ <br> 40914365 |

One of the major issues with the novel algorithm is that it gives large number of false positives despite filtering out variants with negative $\mathrm{Hi}-\mathrm{C}$ scores. Therefore distribution-based filtering is another approach to reduce false positives. Frequency distribution of $\mathrm{Hi}-\mathrm{C}$ data for chromosome 1 and chromosome 22 showed most of the regions ( 1 Mbps windows) with Pearson's correlation coefficient between 0.03 and 0.09 (Figure 5.2).


Figure 5.2: Hi-C score frequency distribution for chromosome 1 and 22

Setting the filtering cutoff at 0.09 would make the tool miss this translocation while setting it at the tail end at 0.249 will capture it (Table 5.4).

Frequency distribution of chromosome 12 and chromosome 13 (Figure 5.3) with cutoff above zero will include the translocation (Table 5.4).


Figure 5.3: Hi-C score frequency distribution for chromosome 12 and 13
Similarly to capture $\mathrm{t}(12 ; 16), \mathrm{t}(19 ; 22), \mathrm{t}(16 ; 21), \mathrm{t}(9 ; 14), \mathrm{t}(7 ; 11), \mathrm{t}(4 ; 8), \mathrm{t}(4 ; 11)$, and $t(8 ; 12)$ we need to include the far right end of the distribution (Figure 5.4, Figure 5.5, Figure 5.6, Figure 5.7, Figure 5.8, Figure 5.9, Figure 5.10, Figure 5.11). Thus, although these translocations had a strong Pearson's correlation value, filtering based on distribution will lose these translocations. A filtering method to filter out variants which are one standard deviation away from the mean in the positive direction will be a better approach to capture these events.


Figure 5.4: Hi-C score frequency distribution for chromosome 12 and 16


Figure 5.5: Hi-C score frequency distribution for chromosome 19 and 22


Figure 5.6: Hi-C score frequency distribution for chromosome 16 and 21


Figure 5.7: Hi-C score frequency distribution for chromosome 9 and 14


Figure 5.8: Hi-C score frequency distribution for chromosome 7 and 11


Figure 5.9: Hi-C score frequency distribution for chromosome 4 and 8


Figure 5.10: Hi-C score frequency distribution for chromosome 4 and 11


Figure 5.11: Hi-C score frequency distribution for chromosome 8 and 12

Thus in order to reduce noise, filtering variants with negative correlation coefficients is an essential first step. Further, arriving at a cutoff by looking at the distribution of $\mathrm{Hi}-\mathrm{C}$ scores to filter variants with strong positive correlation can further reduce noise and fine-tune the tool.

## 6. DISCUSSION

The aim of this study was to understand if characteristics of the structural variant made it more or less conducive to detection by current computational methods in use. We wanted to understand the structural variant in relation to its position in the genome and determine if this genomic context made some variants more difficult to be detected by current tools. The human genome is made up of approximately 20,000-25,0000 genes and targeted sequencing of only these known coding regions covers about 1\% of the 3 billion bases of the human genome, i.e. approximately 30 million (Brunham \& Hayden, 2013). Applying genome linkage analysis to single-gene Mendelian disease has met with considerable success as in cystic fibrosis (Kerem et al., 1989; Rommens et al., 1988) and Huntington's disease (Fox, Bloch, Fahy, \& Hayden, 1989). Application of next generation sequencing efforts to Mendelian disorders has led to discovery of more than 3000 genes associated with a phenotypically visible trait and more than 5000 phenotypic traits with known molecular basis as recorded to date in the Online Mendelian Inheritance in Man (OMIM database) (McKusick, 2014). Discovery of these specific genes has tremendous value in the predictive capability of genetic diagnosis through biomarkers as in Alzheimer as well as targeted gene therapy as in cystic fibrosis with the mean survival rate of cystic fibrosis going from 6 months in 1930's to 40 years in 1990's (Ikpa, Bijvelds, \& de Jonge, 2014) due to novel therapy. The impact of genotype-to-phenotype translation, while more obvious in Mendelian diseases, are not so easily translated in complex disease with possible multi-gene etiology. Translational bioinformatics tries to fill up this gap using high throughput data analysis techniques. Cancers are even more complicated due to presence of somatic variants that are not inherited and therefore more difficult to define. Further, tumors are highly heterogeneous, and therefore the presence of a variant will be highly
dependent on sample extraction and preparation. Thus, the aim is to identify the best approach to detect novel variants in a heterogeneous sample with highly mutated genome. The currently available tools were designed around the normal genome architecture assuming $100 \%$ allele frequency. Cancer genome does not conform to this assumption.

### 6.1 Repeat Analysis

The study analysis started with the hypothesis that the tools were failing to detect variants due to presence of these variants in repeat regions. If the mapping tools are unable to map reads uniquely, the tools will not be able to pick these variants. The first analysis for this study tried to define the characteristics of the validated variants in 1000 genomes trio dataset using current variant detection algorithms. The purpose was to understand the reason these tools were missing to detect the validated variants. SVDetect was chosen for this analysis due to its high false positive rate and thus be highly sensitive. Since the idea was to define context, true positive was defined as those variants detected by SVDetect with at least $10 \%$ of insert size overlap with variants in the 1000-genome validated file. Although less than $50 \%$ of the true variants had some overlap with the variants detected by the tools, their representation in the repeat regions was not as hypothesized. The hypothesis that structural variants missed by the tools were more likely to be in repeat regions was rejected by the analysis. A greater percent of the variants detected by the tools had repeats compared to variants not detected by the tools. This was consistent across all four variant types, deletion, mobile element insertion, tandem duplication, and novel sequences classified in the 1000-Genomes dataset. Repeat structure was not driving the tool's inability to pick less than $50 \%$ of true variants. Since $50 \%$ of the human genome is made up of repeats, this very broad classification of context did not prove very useful. Since the study was trying to
design a tool specific for translocation, the context was now focused towards mechanisms driving translocation which can be captured informatically.

### 6.2 Algorithm Development and Simulated Data Analysis

Translocations lead to genetic imbalances and are a precursor to cancers. Detection of the same recurrent translocation in four unrelated families by Ou et al. (Ou et al., 2011) led to implication of non-allelic homologous recombination (NAHR) as the driver for these variants in all these subjects. NAHR occurs due to aberrant DNA repair mechanism between regions that share considerable homology, also known as low copy repeat regions (LCRs). Unlike repeat elements, these LCRs are several thousand basepairs long and share greater than $95 \%$ identity. Using these characteristics, Ou et al. computationally mapped the NAHR regions on the human genome and predicted validated translocations in their database. This mapping was based on the segmental duplication map created using comparative genomic hybridization which identified novel structural variants in these regions of LCRs by Sharp et al. (Sharp et al., 2006). Segmental duplications as possible hotspots for structural variation events were first hypothesized and mapped by Bailey et al. (Bailey et al., 2002) identifying 169 such regions in the human genome. However, the physical co-location of these LCR regions in cell is also an important contributing factor to the actual interaction between these regions. Regions of chromosomes that are physically close to each other in 3-dimensional space are more likely to interact with each other as proved by chromosomal conformation capture experiments known as (3C or $\mathrm{Hi}-\mathrm{C}$ ) which reveal three-dimensional architecture of genome packing in the cell (Wijchers, 2011). The mobility of different regions in the genome is limited by the location of these regions in the genome (Chubb, Boyle, Perry, \& Bickmore, 2002). While accounting for these two types of important contextual based information about translocation etiology, the study was able to design and test the
algorithm based on this context and proved to do better than current methodologies for translocation detection.

The algorithm was designed to capture all reads that would be ambiguously aligned by the mapping tool. The most popular mapping tool currently in use, namely BWA (Li \& Durbin, 2009), was used for generating initial mapping. BWA's popularity is due to its speed in mapping billions of shorts reads in hours, achieved by its effective use of cache memory in indexing the reads and also wide acceptance in the bioinformatics community as a preferred mapping tool. The output from BWA is also in the SAM/BAM format (Li et al., 2009) accepted as the standard output format for alignment. BWA randomly assigns reads which map to more than one region on the genome. In their simulated read alignment, BWA mapped 11/1,569,108 incorrectly (Li \& Durbin, 2009), which still is a very low error rate but in a sequencing experiment with for example 50 million reads, approximately 500 reads would be wrongly assigned. While this may be sufficient for a normal genome, in a cancer genome sequencing project aimed at finding novel variants and showing heterogeneity, this number might make a difference. Analysis in this study showed that even with $100 \%$ allele frequency, which is not the case in tumors; the tools were missing variants in the simulated dataset. The tools pick up anomalous reads based on these reads mapping at greater/shorter distance than the normal distribution of insert size or incorrect orientation. Our novel algorithm does not use the probability distribution of insert size to pick reads. This is because reads with NAHR characteristics, i.e. reads which share greater than $95 \%$ identity, could have been placed at a location selected randomly by the mapping tool since these regions are so similar, and thus their imputed insert size is suspect. The algorithm accounted for this information by extracting reads which could have multiple mappings on the genome while extracting the partner paired read irrespective of its mapping score.

Re-mapping of these ambiguous reads with a local alignment tool like BLAT (W. J. Kent, 2002) further helps define context, since all possible genomic regions of identity are now reported. Further, instead of using read distribution of the data and number of reads supporting a type of variant to assign probability of calling a true variant, the algorithm again used context to define the probability of two regions being involved in translocation based on their known proximity to each other in 3-D cellular space and therefore the probability of interacting regions (Lieberman-Aiden et al., 2009). Even with 15 X coverage, the novel tool was able to detect all ten simulated translocations.

### 6.3 Conclusion

Designing a novel context based approach to detect translocations, the study showed a very effective way to detect these variants using a biologically derived context-based approach which has not been used so far to effectively mine structural variants. The study also rejected the hypothesis that repeat structure within the variant was driving the inability of current tools to detect true positive events.

The output from this novel algorithm could help discover many de-novo variants in cancers and provide a starting point for mining variant information from sequencing data. The purpose of this tool was not to define a few variants, but to give as many possible variants that could then be teased out by the user through experimental validation. Bioinformatics data analysis of such big volume data does suffer from copious output of false positives, but at the same time is the first step in moving towards more comprehensive follow-up using laboratory tools. Providing the user with biological context-based algorithm instills more confidence in the output, which was the purpose of this analysis.

### 6.4 Limitations

The major limitation of this study is the number of false positive reported in output. Ideally the user would like to see only the most relevant information that is currently embedded in a lot of noise. Noise can be reduced by narrowing BLAT's extraction parameters for stricter re-alignment. How this would play out in a real dataset was not explored in this analysis.

Since re-alignment with BLAT is a computationally intensive process, the access to high performance computing environment is a pre-requisite. Without highthroughput computing the analysis can get prohibitively time consuming. Running BLAT alignments in parallel greatly reduces the time, and we assume that users would have access to parallel, high performance computing resources.

The tool was also specifically designed for detecting translocations, and users would prefer getting the entire spectrum of structural variation in a single tool, which is another limitation for this study.

### 6.5 Future Direction and Research

The tool was specifically designed to obtain as much information from mapping as possible to be able to derive de-novo variants that it achieved at the cost of reporting a lot of noise. I would like to explore new methods to reduce noise in the data without compromising on the mining ability of the tool. Noise reduction parameters could also include evolutionary information of conserved versus nonconserved regions to remove implausible variants. I would also like to expand the tool capability to detect all type of other structural variants including deletion, insertion and inversion.

This is the first context-based tool designed to date and can prove useful for helping lay the framework for further algorithm development along these lines which take other biological context into account while designing bioinformatics tools.

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APPENDIX A
[LIST OF 1000-GENOMES FILES USED IN ANALYSIS]

## NA19238

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## NA19239

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## NA19238

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## NA19239

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## NA19240

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## NA19240

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| SRR004484 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002352 |
| SRR004484 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002352 |
| SRR004485 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002353 |
| SRR004485 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002353 |
| SRR004783 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002629 |
| SRR004783 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002629 |
| SRR004784 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002630 |
| SRR004784 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002630 |
| SRR004785 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002631 |
| SRR004785 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002631 |
| SRR004786 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002632 |
| SRR004786 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002632 |
| SRR004788 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002636 |
| SRR004788 | SRP000032 | 1000Genomes Project Pilot 2 | WUGSC | SRA002636 |


| NA19238 | NA19238 | NA19238 | NA19238 | NA19238 |
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| 9/24/2008 0:00 | SRS000212 | NA19238 | YRI | SRX001106 |
| 9/24/2008 0:00 | SRS000212 | NA19238 | YRI | SRX001106 |
| 9/24/2008 0:00 | SRS000212 | NA19238 | YRI | SRX001106 |
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| 8/15/2008 0:00 | SRS000213 | NA19239 | YRI | SRX000654 |
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| NA19238 | NA19238 |
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| _PLATFORM |  |
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| ILLUMINA | Illumina Genome Analyzer II |
| ILLUMINA | Illumina Genome Analyzer II |
| ILLUMINA | Illumina Genome Analyzer II |
| ILLUMINA | Illumina Genome Analyzer II |
| ILLUMINA | Illumina Genome Analyzer II |
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| ILLUMINA | Illumina Genome Analyzer II |
| ILLUMINA | Illumina Genome Analyzer II |
| ILLLUMINA | ILllumina Genome Analyzer |
| ILLUMINA | Illumina Genome Analyzer |
| ILLLUMMINA | Illumina Genome Analyzer II |
| ILLUMINA | Illumina Genome Analyzer II |
| ILLUMINA | Illumina Genome Analyzer II |
| ILLUMINA | Illumina Gena Gena |


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| :---: | :---: | :---: |
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| 2675169269 | 7592 | 260 |
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| 2675169269 | 7612 | 260 |
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| LIBRARY_NAME | RUN_NAME | INSERT_SIZE |
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| 2485443314 | 5686 | 260 |
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| 2485443314 | 5688 | 260 |


| ILLUMINA | Illumina Genome Analyzer | 2485443314 | 5688 | 260 |
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| ILLUMINA | Illumina Genome Analyzer | 2485443314 | 5895 | 260 |
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| ILLUMINA | Illumina Genome Analyzer | 2485443314 | 6430 | 260 |
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| INSTRUMENT_L | INSTRUMENT_MODEL | LIBRARY_NAME | RUN_NAME | INSERT_SIZE |
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| ILLUMINA | Illumina Genome Analyzer II | 2675080346 | 7224 | 260 |
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| ILLUMINA | Illumina Genome Analyzer II | 2675080346 | 7522 | 260 |
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| ILLUMINA | Illumina Genome Analyzer II | 2675080346 | 7526 | 260 |
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| PAIRED | data/NA19238/sequence_read/SRR005193_2.filt.fastq.gz | 11173375 |
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| NA19239 | NA19239 | NA19239 |
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| PAIRED | data/NA19239/sequence_read/SRR007422_1.filt.fastq.gz | 6497852 |
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| 405730548 high coverage |
| 402241500 high coverage |
| 402241500 high coverage |
| 370827504 high coverage |
| 370827504 high coverage |
| 321516576 high coverage |
| 321516576 high coverage |
| 290735100 high coverage |
| 290735100 high coverage |
| 265695444 high coverage |
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| 208535076 high coverage |
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| 162360360 high coverage |
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| 179863092 high coverage |
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| 231213708 high coverage |
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| 303081048 high coverage |
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| NA19239 |
| 266904936 high coverage |
| 26004289872 high coverage |
| 261289872 high coverage |
| 258137748 high coverage |
| 258137748 high coverage |
| 260820432 high coverage |
| 260820432 high coverage |

```
    266904936 high coverage
    292139748 high coverage
    292139748 high coverage
    290855916 high coverage
    290855916 high coverage
    286548336 high coverage
    286548336 high coverage
    186753852 high coverage
    186753852 high coverage
    222255108 high coverage
    222255108 high coverage
    218009520 high coverage
    218009520 high coverage
    242846460 high coverage
    2 4 2 8 4 6 4 6 0 ~ h i g h ~ c o v e r a g e ~
    238362984 high coverage
    238362984 high coverage
    241434144 high coverage
    241434144 high coverage
    209771064 high coverage
    209771064 high coverage
    233922672 high coverage
    233922672 high coverage
NA19240 NA19240
BASE_COUNT ANALYSIS_GROUP
    369375895 high coverage
    369375895 high coverage
    404750535 high coverage
    4 0 4 7 5 0 5 3 5 \text { high coverage}
    358825670 high coverage
    358825670 high coverage
    347263840 high coverage
    347263840 high coverage
    385727440 high coverage
    385727440 high coverage
    398260765 high coverage
    398260765 high coverage
    4 0 5 3 3 5 8 7 5 ~ h i g h ~ c o v e r a g e ~
    4 0 5 3 3 5 8 7 5 ~ h i g h ~ c o v e r a g e ~
    391054895 high coverage
    391054895 high coverage
```


## APPENDIX B

[BLAT SUITE PROGRAM SPECIFICATIONS AND USER'S GUIDE]

## BLAT Suite Program Specifications and User Guide

## General:

Blat produces two major classes of alignments: at the DNA level between two sequences that are of $95 \%$ or greater identity, but which may include large inserts, and at the protein or translated DNA level between sequences that are of $80 \%$ or greater identity and may also include large inserts. The output of BLAT is flexible. By default it is a simple tab-delimited file which describes the alignment, but which does not include the sequence of the alignment itself. Optionally it can produce BLAST and WU-BLAST compatable output as well as a number of other formats.

The main programs in the blat suite are:

- gfServer - a server that maintains an index of the genome in memory and uses the index to quickly find regions with high levels of sequence similarity to a query sequence.
- gfClient - a program that queries gfServer over the network, and then does a detailed alignment of the query sequence with regions found by gfServer.
- blat -combines client and server into a single program, first building the index, then using the index, and then exiting.
- webBlat - a web based version of gfClient that presents the alignments in an interactive fashion.

Building an index of the genome typically takes 10 or 15 minutes. Typically for interactive applications one uses gfServer to build a whole genome index. At that point gfClient or webBlat can align a single query within few seconds. If one is aligning a lot of sequences in a batch mode then blat can be more efficient, particularly if run on a cluster of computers. Each blat run is typically done against a single chromosome, but with a large number of query sequences.

Other programs in the blat suite are:

- pslSort - combines and sorts the output of multiple blat runs. (The blat default output format is .psl).
- pslReps - selects the best alignments for a particular query sequence, using a 'near best in genome' approach.
- pslPretty - converts alignments from the psl format, which is tabdelimited format and does not include the bases themselves, to a more readable alignment format.
- faToTwoBit - convert Fasta format sequence files to a dense randomlyaccessable .2bit format that gfClient can use.
- twoBitToFa - convert from the .2bit format back to fasta
- faToNib - convert from Fasta to a somewhat less dense randomly accessible format that predates .2bit. Note each .nib file can only contain a single sequence.
- nibFrag - convert portions of a nib file back to fasta.

In addition you may be interested in the following programs which are not part of the BLAT
suite:

- In Silico PCR - given two primers quickly find the sequence between them. Available from Kent Informatics. This includes webPCR, an interface similar to webBlat.
- The Genome Browser - display annotations as a series of tracks on top of the genome.

Available from the University of California
Santa Cruz. See
http://genome.ucsc.edu/license/.

## Running the Programs:

The command line options of each of the programs is described below. Similar summaries of usage are printed when a command is run with no arguments. See the next section for info on installing webBlat.

```
blat
blat - Standalone BLAT sequence search
command line tool usage:
```

blat database query [-ooc=11.ooc]
output.psl where:

```
    database and query are each either a .fa , .nib or .2bit file,
    or a list these files one file name per line.
    -ooc=11.ooc tells the program to load over-occurring
        11-mers from and external file. This
        will increase the speed
        by a factor of 40 in many cases, but is
    not required output.psl is where to put the output.
    Subranges of nib and .2bit files may specified using the syntax:
        /path/file.nib:seqid:start-end
    or
        /path/file.2bit:seqid:start-end
```

    or
        /path/file.nib:start-end
    With the second form, a sequence id of file:start-end will
    be used. options:
-t=type Database type. Type is one of:
dna - DNA sequence
prot - protein sequence
dnax - DNA sequence translated in six frames to protein
The default is dna
-q=type Query type. Type is one of:
dna - DNA
sequence
rna - RNA
sequence

```
    prot - protein sequence
    dnax - DNA sequence translated in six frames to
    protein rnax - DNA sequence translated in three
    frames to protein
    The default is dna
    -prot Synonymous with -t=prot -q=prot
    -ooc=N.ooc Use overused tile file N.ooc. N should
    correspond to the tileSize
    -tileSize=N sets the size of match that triggers an alignment.
    Usually between 8 and 12
    Default is }11\mathrm{ for DNA and 5 for protein.
    -stepSize=N spacing between tiles. Default is tileSize.
    -oneOff=N If set to 1 this allows one mismatch in tile and still
        triggers an alignments. Default
        is 0.
-minMatch=N sets the number of tile matches. Usually set from 2 to 4
    Default is 2 for nucleotide, 1 for
    protein.
-minScore=N sets minimum score. This is the matches minus the
    mismatches minus some sort of gap penalty. Default
    is 30
-minIdentity=N Sets minimum sequence identity (in percent). Default is
    90 for nucleotide searches, }25\mathrm{ for protein or
    translated protein searches.
-maxGap=N sets the size of maximum gap between tiles in a clump.
                                    Usually set from 0 to 3.
                                    Default is 2. Only relevent for
        minMatch > 1.
```

```
-noHead suppress .psl header (so it's just a tab-separated file)
-makeOoc=N.ooc Make overused tile file. Target needs to be complete
genome.
-repMatch=N sets the number of repetitions of a tile allowed before
    it is marked as overused. Typically this is 256 for
    tileSize
    12, 1024 for tile size 11, 4096 for tile
    size 10.
    Default is 1024. Typically only comes into play with
    makeOoc. Also affected by stepSize. When stepSize is
    halved repMatch is doubled to compensate.
-mask=type Mask out repeats. Alignments won't be started in masked region
    but may extend through it in nucleotide searches. Masked
    areas are ignored entirely in protein or translated
    searches. Types are
        lower - mask out lower cased sequence
        upper - mask out upper cased sequence
        out - mask according to database.out RepeatMasker
        .out file file.out - mask database according to
        RepeatMasker file.out
-qMask=type Mask out repeats in query sequence. Similar to -mask
    above but for query rather than target sequence.
-repeats=type Type is same as mask types above. Repeat bases will
    not be masked in any way, but matches in repeat areas
    will be reported separately from matches in other areas
    in the psl output.
```

```
-minRepDivergence=NN - minimum percent divergence of repeats to
    allow them to be unmasked. Default is 15. Only
    relevant for masking using RepeatMasker .out files.
-dots=N Output dot every N sequences to show program's progress
-trimT Trim leading poly-T
-noTrimA Don't trim trailing poly-A
-trimHardA Remove poly-A tail from qSize as well as
    alignments in psl output
    Run for fast DNA/DNA remapping - not allowing
    introns, requiring high %ID
-out=type Controls output file format. Type is one of:
    psl - Default. Tab separated format, no
    sequence pslx - Tab separated format with
    sequence
    axt - blastz-associated axt
    format maf - multiz-
    associated maf format sim4 -
    similar to sim4 format
    wublast - similar to wublast
    format
    blast - similar to NCBI blast
    format blast8- NCBI blast
    tabular format
    blast9 - NCBI blast tabular format with comments
-fine For high quality mRNAs look harder for small
    initial and terminal exons.Not recommended for
    ESTs
```

```
-maxIntron=N Sets maximum intron size. Default is 750000
-extendThroughN - Allows extension of alignment through large blocks of
N's
```

Here are some blat settings for common usage scenarios:

1) Mapping ESTs to the genome within the same species
$-00 c=11.00 c$
2) Mapping full length mRNAs to the genome in the same species - ooc=11.ooc -fine -q=rna
3) Mapping ESTs to the genome across species
```
-q=dnax -t=dnax
```

4) Mapping mRNA to the genome across species -q=rnax -t=dnax
5) Mapping proteins to the genome
$-q=p r o t-t=d n a x$
6) Mapping DNA to DNA in the same species
```
-ooc=11.ooc -fastMap
```

7) Mapping DNA from one species to another species
```
-q=dnax -t=dnax
```

When mapping DNA from one species to another the
query side of the alignment should be cut up
into chunks of 25 kb or less for best
performance.
gfServer
gfServer - Make a server to quickly find where DNA occurs
in genome. To set up a server:
gfServer start host port file(s)
Where the files are in .nib or .2bit format

To remove a server:
gfServer stop host port
To query a server with DNA sequence:
gfServer query host port probe.fa

To query a server with protein sequence:
gfServer protQuery host port probe.fa

To query a server with translated dna sequence:
gfServer transQuery host port probe.fa
To process one probe fa file against a .nib format genome (not starting server):
gfServer direct probe.fa file(s).nib
To figure out usage
level gfServer
status host port
To get input file
list gfServer
files host port

Options:
-tileSize=N size of $n$-mers to index. Default is 11 for nucleotides, 4 for proteins (or translated nucleotides).
-stepSize=N spacing between tiles. Default is tileSize.
-minMatch=N Number of n -mer matches that trigger detailed alignment

```
        Default is 2 for nucleotides, 3 for
        protiens.
    -maxGap=N Number of insertions or deletions allowed between n-mers.
        Default is 2 for nucleotides, 0 for
        protiens.
    -trans Translate database to protein in 6 frames. Note:
        it is best to run this on RepeatMasked data in this
        case.
    -log=logFile keep a log file that records server requests.
    -seqLog Include sequences in log file (not logged with -syslog)
    -syslog Log to syslog
    -logFacility=facility log to the specified syslog facility - default
local0.
-mask Use masking from nib file.
-repMatch=N Number of occurrences of a tile (nmer)
        that trigger repeat masking the tile.
        Default is 1024.
-maxDnaHits=N Maximum number of hits for a dna query that are sent from
the
    server. Default is 100.
-maxTransHits=N Maximum number of hits for a translated query
    that are sent from the server. Default is 200.
-maxNtSize=N Maximum size of untranslated DNA query sequence
    Default is 40000
-maxAsSize=N Maximum size of protein or translated DNA queries
        Default is 8000
-canStop If set then a quit message will actually take
        down the server
```

gfClient

```
gfClient - A client for the genomic
finding program usage:
    gfClient host port nibDir in.fa out.psl
where
    host is the name of the machine running the
    gfServer port is the same as you started
    the gfServer with
    nibDir is the path of the nib files relative to the current dir
        (note these are needed by the client as well as the server)
    in.fa a fasta format file. May contain
    multiple records out.psl where to put the
    output
options:
    -t=type Database type. Type is one of:
    dna - DNA sequence
    prot - protein sequence
    dnax - DNA sequence translated in six frames to protein
    The
        default
    is dna
    -q=type Query type. Type is one of:
    dna - DNA
    sequence
    rna - RNA
    sequence
    prot - protein sequence
```

```
        dnax - DNA sequence translated in six frames to
        protein rnax - DNA sequence translated in three
        frames to protein
    -dots=N Output a dot every N query sequences
    -nohead Suppresses psl five line header
    -minScore=N sets minimum score. This is twice the matches
        minus the mismatches minus some sort of gap
        penalty. Default is 30
-minIdentity=N Sets minimum sequence identity (in percent). Default is
    90 for nucleotide searches, 25 for protein or
    translated protein searches.
-out=type Controls output file format. Type is one of:
        psl - Default. Tab separated format without actual
        sequence pslx - Tab separated format with sequence
        axt - blastz-associated axt
        format maf - multiz-
        associated maf format
        wublast - similar to
        wublast format blast -
        similar to NCBI blast
        format
-maxIntron=N Sets maximum intron size. Default is }75000
```


## webBlat

webBlat generally is not run from the command line. See 'Setting Up webBlat instructions below` for information on this program.

## faToTwoBit

```
faToTwoBit - Convert DNA from fasta to
2bit format usage:
    faToTwoBit in.fa [in2.fa in3.fa ...]
out.2bit options:
    -noMask - Ignore lower-case masking in fa file.
```


## twoBitToFa

```
twoBitToFa - Convert all or part of .2bit file
to fasta usage:
    twoBitToFa input.2bit
output.fa options:
    -seq=name - restrict this to just one sequence
    -start=X - start at given position in sequence (zero-based)
    -end=X - end at given position in sequence (non-inclusive)
```


## faToNib

```
faToNib - Convert from .fa
```

to .nib format usage:
faToNib in.fa out.nib

## nibFrag

```
nibFrag - Extract part of a
```

nib file as .fa usage:
nibFrag file.nib start end strand out.fa

```
pslPretty
pslPretty - Convert PSL to human
readable output usage:
    pslPretty in.psl target.lst
query.lst pretty.out options:
    -axt - save in Scott Schwartz's axt format
    -dot=N Put out a dot every N records
    -long - Don't abbreviate long inserts
```

It's a really good idea if the psl file is sorted by target if it contains multiple targets. Otherwise this will be very very slow. The target and query lists can either be fasta files, nib files, or a list of fasta and/or nib files one per line. Currently this only handles nucleotide based psl files.

```
pslSort
pslSort - merge and sort psCluster
.psl output files usage:
    pslSort dirs[1|2] outFile
            tempDir inDir(s)
This will sort all of the .psl files in the
directories inDirs in two stages - first into
temporary files in tempDir and second into outFile.
        The device on tempDir
needs to have
enough space (typically 15-20 gigabytes if processing whole genome)
    pslSort g2g[1|2] outFile tempDir inDir(s)
```

```
This will sort a genome to genome alignment,
reflecting the alignments across the diagonal.
```

Adding 1 or 2 after the dirs or g2g will limit the program to only the first or second pass repectively
of the sort
Options:
-verbose=N Set verbosity level, higher for more output. Default 1
Note for huge files pslSort will run out of memory. The unix
sort command sort -k 10 *.psl > sorted.psl
may be preferable in these situations, though the psl header
lines should be removed or avoided with the -noHead option to
blat.

## pslReps

pslReps - analyse repeats and generate genome
wide best alignments from a sorted set of
local alignments
usage:
pslReps in.psl out.psl out.psr
where in.psl is an alignment file generated by
psLayout and sorted by pslSort, out.psl is the
best alignment output
and out.psr contains repeat info

```
options:
    -nohead don't add PSL header
-ignoreSize Will not weigh in favor of larger alignments so much
-noIntrons Will not penalize for not having introns when
            calculating size factor
-singleHit Takes single best hit, not splitting into parts
-minCover=0.N minimum coverage to output. Default is 0.
-ignoreNs Ignore 'N's when calculating minCover.
-minAli=0.N minimum
    alignment ratio
    default is 0.93
-nearTop=0.N how much can deviate from top and be taken
    default is 0.01
-minNearTopSize=N Minimum size of alignment that is
    near top for aligmnent to be kept.
        Default 30.
-coverQSizes=file Tab-separate file with effective query sizes.
    When used with -minCover, this
    allows polyAs to be excluded from
    the coverage calculation
```


## Setting Up webBlat

INSTALLING WEBBLAT
Installing A Web-Based Blat Server involves four major steps:

1) Creating sequence databases.
2) Running the gfServer program to create in-memory indexes of the databases.
3) Editing the webBlat.cfg file to tell it what machine and port the gfServer(s)
are running on, and optionally customizing the webBlat appearance to users.
4) Copying the webBlat executable and webBlat.cfg to a directory where the web server can execute webBlat as a CGI.

CREATING SEQUENCE DATABASES

You create databases with the program faToTwoBit. Typically you'll create a separate database for each genome you are indexing. Each database can contain up to four billion bases of sequence in an unlimited number of records. The databases for webPcr and webBlat are identical.

The input to faToTwoBit is one or more fasta format files each of which can contain multiple records. If the sequence contains repeat sequences, as is the case with vertebrates and many plants, the repeat sequences can be represented in lower case and the other sequence in upper case. The gfServer program can be configured to ignore the repeat sequences. The output of faToTwoBit is a file which is designed for fast random access and efficient storage. The output files store four bases per byte. They use a small amount of additional space to store the case of the DNA and to keep track of runs of

N 's in the input. Non- N ambiguity codes such as Y and U in the input sequence will be converted to N .

Here's how a typical installation might create a mouse and a human genome database:
cd/data/genomes mkdir twoBit
faToTwoBit human/hg16/*.fa twoBit/hg16.2bit faToTwoBit mouse/mm4/*.fa twoBit/mm4.2bit
There's no need to put all of the databases in the same directory, but it can simplify bookkeeping.

The databases can also be in the .nib format which was used with blat and gfClient/gfServer until recently. The .nib format only packed 2 bases per byte, and could only handle one record per nib file. Recent versions of blat and related programs can use . 2 bit files as well.

## CREATING IN-MEMORY INDICES WITH GFSERVER

The gfServer program creates an in-memory index of a nucleotide sequence database. The index can either be for translated or untranslated searches. Translated indexes enable protein-based blat queries and use approximately two bytes per unmasked base in the database. Untranslated indexes are used nucleotide-based blat queries as well
as for In-silico PCR. An index for normal blat uses approximately $1 / 4$ byte per base. For blat on smaller (primer-sized) queries or for In-silico PCR a more thorough index that requires $1 / 2$ byte per base is recommended. The gfServer is memory intensive but typically doesn not require a lot of CPU power. Memory permitting multiple gfServers can be run on the same machine.

A typical installation might go:
ssh bigRamMachine
cd/data/genomes/twoBit
gfServer start bigRamMachine 17779 hg16.2bit \&
gfServer -trans -mask start bigRamMachine 17778 hg16.2bit \&
the -trans flag makes a translated index. It will take approximately
15 minutes to build an untranslated index, and 45 minutes to build a translate index. To build an untranslated index to be shared with In-silico PCR do
gfServer -stepSize=5 bigRamMachine 17779 hg16.2bit \&
This index will be slightly more sensitive, noticeably so for small query sequences, with blat.

## EDITING THE WEBBLAT.CFG FILE

The webBlat.cfg file tells the webBlat program where to look for gfServers and for sequence. The basic format of the .cfg file is line oriented with the first word of the line being a command. Blank lines and lines starting with \#
are ignored. The webBlat.cfg and webPcr.cfg files are similar. The webBlat.cfg commands are:
gfServer - defines host and port a (untranslated) gfServer is running on, the associated sequence directory, and the name of the database to display in the webPcr web page.
gfServerTrans - defines location of a translated server.
background - defines the background image if any to display on web page company - defines company name to display on web page
tempDir - where to put temporary files. This path is relative to where the web server executes CGI scripts. It is good to remove files that haven't been accessed for 24 hours from this directory periodically, via a cron job or similar mechanism.
The background and company commands are optional. The webBlat.cfg file must have at least one valid gfServer or gfServerTrans line, and a tempDir line. . Here is a webBlat.cfg file that you might find at a typical installation:
company Awesome Research Amalgamated background /images/dnaPaper.jpg gfServer bigRamMachine 17778 /data/genomes/2bit/hg16.2bit Human Genome gfServer bigRamMachine 17779 /data/genomes/2bit/hg16.2bit Human Genome gfServer mouseServer 17780 /data/genomes/2bit/mm4.2bit Mouse Genome gfServer mouseServer 17781 /data/genomes/2bit/mm4.2bit Mouse Genome tempDir ../trash

## PUTTING WEbBLAT WHERE THE WEB SERVER CAN EXECUTE IT

The details of this step vary highly from web server to web server. On a typical Apache installation it might be:
ssh webServer cd kent/webBlat
cp webBlat webBlat.cfg /usr/local/apache/cgi-bin mkdir/usr/local/apache/trash chmod 777 /usr/local/apache/trash
assuming that you've put the executable and config file in kent/webBlat. The program will create some files in the trash directory. It is good to periodically clean out old files from this directory. On Mac OS-X instead you might do: cp webBlat webBlat.cfg/Library/WebServer/CGI-Executables mkdir /Library/WebServer/trash
chmod 777 /Library/WebServer/trash
Unless you are administering your own computer you will likely need to ask your local system administrators for help with this part of the webBlat installation.

## File Formats

.psl files
A .psl file describes a series of alignments in a dense easily parsed text format. It begins with a five line header which describes each field. Following this is one line for each alignment with a tab between each field. The fields are describe below in a format suitable for many relational databases.
matches int unsigned, \# Number of bases that match that aren't repeats
misMatches int unsigned, \# Number of bases that don't match
repMatches int unsigned, \# Number of bases that match but are part of repeats nCount int unsigned, \# Number of 'N' bases
qNumlnsert int unsigned, \# Number of inserts in query
qBasel nsert int unsigned, \# Number of bases inserted in query
tNumI nsert int unsigned, \# Number of inserts in target
tBasel nsert int unsigned, \# Number of bases inserted in target
strand char(2), \# + or - for query strand, optionally followed by + or - for target strand
qName varchar(255), \# Query sequence name
qSize int unsigned, \# Query sequence size
qStart int unsigned qEnd int unsigned , tName varchar(255), tSize int unsigned ,

> \# Alignment start position in query
> \# Alignment end position in query
> \# Target sequence name
> \# Target sequence size
tStart int unsigned, \# Alignment start position in target
tEnd int unsigned, \# Alignment end position in target
blockCount int unsigned, \# Number of blocks in alignment. A block contains no gaps.
blockSizes longblob, \# Size of each block in a comma separated list qStarts longblob, \# Start of each block in query in a comma separated list tStarts longblob, \# Start of each block in target in a comma separated list

In general the coordinates in psl files are "zero based half open." The first base in a sequence is numbered zero rather than one. When representing a range the end coordinate is not included in the range. Thus the first 100 bases of a sequence are represented as $0-100$, and the second 100 bases are represented as 100-200. There is a another little unusual feature in the .psl format. It has to do with how coordinates are handled on the negative strand. In the qStart/qEnd fields the coordinates are where it matches from the point of view of the forward strand (even when the match is on the reverse strand). However on the qStarts[] list, the coordinates are reversed.
Here's an example of a 30 -mer that has 2 blocks that align on the minus strand and 2 blocks on the plus strand (this sort of stuff happens in real life in response to assembly errors sometimes).
$0 \quad 1 \quad 2 \quad 3$ tens position in query
0123456789012345678901234567890 ones position in query
+++++++++ plus strand alignment on query
-------- ---------- minus strand alignment on query
Plus strand:
qStart 12 qEnd 31 blockSizes 4,5 qStarts 12,26
Minus strand:
qStart 4 qEnd 26 blockSizes 10,8 qStarts 5,19
Essentially the minus strand blockSizes and qStarts are what you would get if you reverse complemented the query. However the qStart and qEnd are nonreversed. To get from one to the other:
qStart $=$ qSize - revQEnd qEnd $=q$ Size - revQStart

## .2bit files

A .2bit file can store multiple DNA sequence (up to 4 gig total) in a compact randomly accessible format. The two bit files contain masking information as well as the DNA itself. The file begins with a 16 byte header containing the following fields:

1) signature - the number $0 \times 1 \mathrm{~A} 412743$ in the architecture of the machine that created the file.
2) version - zero for now. Readers should abort if they see a version number higher than 0 .
3) sequenceCount - the number of sequences in the file
4) reserved - always zero for now.

All fields are 32 bits unless noted. If the signature value is not as given, the reader program should byte swap the signature and see if the swapped version matches. If so all multiple-byte entities in the file will need to be byte-swapped. This enables these binary files to be used unchanged on different architectures.

The header is followed by a file index. There is one entry in the index for each sequence. Each index entry contains three fields:

1) nameSize - a byte containing the length of the name field
2) name - this contains the sequence name itself, and is variable length depending on nameSize.
3) offset - 32 bit offset of the sequence data relative to the start of the file

The index is followed by the sequence records. These contain 9 fields:

1) dnaSize - number of bases of DNA in the sequence.
2) nBlockCount - the number of blocks of N's in the file (representing unknown sequence).
3) nBlockStarts - a starting position for each block of N's
4) nBlockSizes - the size of each block of N's
5) maskBlockCount - the number of masked (lower case) blocks
6) maskBlockStarts - starting position for each masked block
7) maskBlockSizes - the size of each masked block
8) packedDna - the dna packed to two bits per base as so: $00-\mathrm{T}, \mathrm{01}$ - C, 10-A, 11 - G.
The first base is in the most significant 2 bits byte, and the last base in the least significant 2 bits, so that the sequence TCAG would be represented as 00011011 . The packedDna field will be padded with 0 bits as necessary so that it takes an even multiple of 32 bit in the file, as this improves i/o performance on some machines.
.nib files
A .nib file describes a DNA sequence packing two bases into each byte. Each nib file contains only a single sequence. A nib file begins with a 32 bit signature which is $0 \times 6 B E 93 D 3 A$ in the architecture of the machine that created the file, and possibly a byte-swapped version of the same number on another machine. This is followed by a 32 bit number in the same format which describes the number of bases in the file. This is followed by the bases themselves packed two bases to the byte. The first base is packed in the high order 4 bits (nibble), the second base in the low order four bits. In C code:
byte $=($ base1 $\ll 4)+$ base2
The numerical values for the bases are:
$0-\mathrm{T}, 1$ - C, 2 - A, 3 - G, 4 - N (unknown)
The most significant bit in a nibble is set if the base is masked.

## Limits

The gfServer program requires approximately 1 byte for every 3 bases in the genome it is indexing in DNA mode, and 1.5 bytes for each unmasked base in translated mode. The blat program requires approximately two bytes for each base in the genome in DNA mode, and three bytes for each base in translated mode. The other programs use relatively little memory.

APPENDIX C
[SIMULATED TRANSLOCATION REFERENCE LIST]

| Additional file 1-Comprehensive list of gene pairs involved in cancer-specific recurrent translocations which result in fusion transcripts |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Translocation | Gene | Location | Fragile Site | Gene | Location | Fragile |
| t(7;12)(p22;q13) | АСТВ | 7p22.1 | FRA7B (common, apc) | GLII | 12q13.3 |  |
| inv(7)(q21q34) | AKAP9 | 7 q 21.2 | FRA7E (common, apc) | BRAF | 7q34 |  |
| t(X;17)(p11;q25) | ASPSCR1 | 17q25.3 |  | TFE3 | Xp11.23 |  |
| inv(2)(p23q35) | ATIC | 2 q 35 |  | ALK | 2p23.2-p23.1 |  |
| t(17;20)(q23;q13) | BCAS4 | 20q13.13 |  | BCAS3 | 17q23.2 |  |
| t(2;3)(p16;q26) | BCL11A | 2p16.1 |  | MDS1 | 3q26.2 |  |
| t(5;14)(q35;q32) | BCL11B | $14 q 32.2$ |  | NKX2E | 5 q 35.2 | FRA5G (rare, folic acid) |
| t(5;14)(935;q32) | BCL11B | $14 q 32.2$ |  | TLX3 | 5 q 3.1 | FRA5G (rare, folic acid) |
| inv(14)(q11q32) | BCL11B | $14 q 32.2$ |  | TRD@ | 14 q 11.2 |  |
| t(14;18)(q32;q21) | BCL2 | 18 q 21.33 | FRA18B (common, apc) | IGH@ | 14q32.33 |  |
| t(2;18)(p11;q21) | BCL2 | 18 q 21.33 | FRA18B (common, apc) | IGK@ | 2p11.2 | FRA2L (rare, folic acid) |
| t(18;22)(q21;q11) | BCL2 | 18q21.33 | FRA18B (common, apc) | IGL@ | 22q11.22 |  |
| (8;19)(q24;q13) | BCL3 | 19q13.31 | FRA19A (common, 5-aza) | MYC | 8 q 24.21 |  |
| t(3;16)(q27;p13) | BCL6 | 3q27.3 | FRA3C (common, apc) | CIITA | 16p13.13 |  |
| t(3;8)(q27;q24) | BCL6 | 3q27.3 | FRA3C (common, apc) | MYC | 8 q 24.21 |  |
| t(1;14)(q21;q32) | BCL9 | 1 q 21.1 | FRA1F (common, apc) | IGH@ | 14 q 32.33 |  |
| (1 $1 ; 22$ )(q21; q 11$)$ | BCL9 | 1 q 21.1 | FRA1F (common, apc) | IGL@ | 22 q 11.22 |  |
| t(9;22)(q34;q11) | BCR | 22q11.23 |  | ABL1 | 9 q 3.12 |  |
| (8;22)(p12;q11) | BCR | 22q11.23 |  | FGFR1 | 8p12 |  |
| t(9;22)(p24;q11) | BCR | 22q11.23 |  | JAK2 | 9 p 24.1 |  |
| (4;22)(q12;q11) | BCR | 22 q 11.23 |  | PDGFRA | 4q12 | FRA4B (common, BrdU) |
| t(11;18)(q22;q21) | BIRC3 | 11q22.2 |  | MALT1 | 18q21.32 | FRA18B (common, apc) |
| t(X;11)(q21; 223 ) | BRWD3 | Xq21.1 |  | ARHGAP20 | 11q22.3-q23.1 |  |
| t(8;12)(q21;q22) | BTG1 | 12q21.33 | FRA12B (common, apc) | MYC | 8 q 24.21 |  |
| t(7;15)(p21;q21) | C15ORF21 | 15q21.1 |  | ETV1 | 7p21.2 |  |
| t(3;3)(q21;q26) | C3ORF27 | 3q21.3 |  | EVI1 | 3q26.2 |  |
| t(2;11)(p23;p15) | CARS | 11p15.4 |  | ALK | 2p23.2-p23.1 |  |
| t(16;16)(p13;q22), inv(16)(p13q22) | CBFB | 16 q 22.1 | FRA16B (rare, dist A), FRA16C (common, apc) | MYH11 | 16p13.11 | FRA16A (rare, folic acid) |
| t(5;10)(q33;q21) | CCDC6 | 10q21.2 | FRA10C (common, BrdU) | PDGFRB | 5 q 33.1 |  |
| $\operatorname{inv}(10)(q 11 q 21)$ | CCDC6 | 10q21.2 | FRA10C (common, BrdU) | RET | 10q11.21 | FRA10G (common, apc) |
| t(5;14)(q33;q32) | CCDC88C | 14 q 32.12 |  | PDGFRB | 5q33.1 |  |
| t(11;19)(q13;p13) | CCND1 | 11q13.2 | FRA11H (common, apc) | FSTL3 | 19p13.3 | FRA19B (rare, folic acid) |
| t(5;6)(q32-33;q22) | CD74 | 5 q 3.1 |  | ROS1 | 6 q 22.2 |  |
| t(16;17)(q21;p13) | CDH11 | 16 q 21 |  | USP6 | 17p13.2 |  |
| t(7;11)(q21;q23) | CDK6 | 7 q 21.2 | FRA7E (common, apc) | MLL | 11q23.3 | FRA11B (rare, folic acid), FRA11G (common, |
| t(5;7)(q35;q21) | CDK6 | 7 q 21.2 | FRA7E (common, apc) | TLX3 | 5 q 35.1 | FRA5G (rare, folic acid) |
| t(5;11)(q12;q23) | CENPK | 5q12.3 |  | MLL | 11q23.3 | FRA11B (rare, folic acid), FRA11G (common, |
| t(8;9)(p12;q33) | CEP110 | 9 q 3.2 |  | FGFR1 | 8p12 |  |


| t(4;12)(q12;p13) | CHIC2 | 4 q 12 | FRA4B (common, BrdU) | ETV6 | 12p13.2 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| t(4;19)(q35;q13) | CIC | $19 \mathrm{q13.2}$ | FRA19A (common, 5-aza) | DUX4 | 4 q 35.2 |  |
| t(2;17)(p23;q23) | CLTC | 17 q 23.1 | FRA17B (common, apc) | ALK | 2p23.2-p23.1 |  |
| t(X;17)(p11;q23) | CLTC | 17 q 23.1 | FRA17B (common, apc) | TFE3 | Xp11.23 |  |
| (2;22)(p23;q11) | CLTCL1 | 22q11.21 |  | ALK | 2p23.2-p23.1 |  |
| t(3;17)(q21;p13) | CNBP | 3q21.3 |  | USP6 | 17p13.2 |  |
| t(17;22)(q21;q13) | COL1A1 | 17 q 21.33 |  | PDGFB | 22q13.1 | FRA22A (rare, folic acid) |
| t(17;17)(p13;q21) | COL1A1 | 17q21.33 |  | USP6 | 17p13.2 |  |
| t(7;8)(q21;q12) | COL1A2 | 7 q 21.3 |  | PLAG1 | 8 q 12.1 |  |
| t(X;6)(q22;q13-14) | COL4A5 | Xq22.3 |  | COL12A1 | 6q13-q14.1 | FRA6D (common, BrdU) |
| t(1;2)(p13;q37) | COL6A3 | 2q37.3 | FRA2J (common, apc) | CSF1 | 1p13.3 |  |
| t(8;12)(p12;q15) | CPSF6 | 12 q 15 |  | FGFR1 | 8 p 12 |  |
| t(11;19)(q21;p13) | CRTC1 | 19p13.11 | FRA19B (rare, folic acid) | MAML2 | 11q21 |  |
| t(11;15)(q21;q26) | CRTC3 | 15 q 26.1 |  | MAML2 | 11q21 |  |
| t(3;8)(p22;q12) | CTNNB1 | 3p22.1 |  | PLAG1 | 8q12.1 |  |
| t(3;9)(q27;p24) | DMRT1 | 9 p 24.3 |  | BCL6 | 3q27.3 | FRA3C (common, apc) |
| t(1;1)(p36;q41) | DUSP10 | 1 q 41 |  | PRDM16 | 1p36.32 | FRA1A (common, apc) |
| t(5;12)(q33;q14) | EBF1 | 5q33.3 |  | LOC204010 | 12q14.3 |  |
| $\mathrm{t}(\mathrm{X} ; 21)(\mathrm{q} 25 ; \mathrm{q} 22)$ | ELF4 | Xq25 |  | ERG | 21q22.2 |  |
| (9;14)(q34;q32) | EML1 | 14q32.2 |  | ABL1 | 9 q 4.12 |  |
| inv(2)(p21p23), del(2)(p21p23)* | EML4 | 2p21 |  | ALK | 2p23.2-p23.1 |  |
| t(5;12)(q33;p13) | ERC1 | 12p13.33 |  | PDGFRB | 5q33.1 |  |
| t(10;12)(q11;p13) | ERC1 | 12p13.33 |  | RET | 10q11.21 | FRA10G (common, apc) |
| t(9;12)(q34;p13) | ETV6 | 12p13.2 |  | ABL1 | 9q34.12 |  |
| t(1;12)(q25;p13) | ETV6 | 12 p 13.2 |  | ABL2 | 1q25.2 |  |
| t(5;12)(q31;p13) | ETV6 | 12 p 13.2 |  | ACSL6 | 5 q 31.1 | FRA5C (common, apc) |
| t(1;12)(q21;p13) | ETV6 | 12p13.2 |  | ARNT | 1q21.2 | FRA1F (common, apc) |
| t(12;12)(p13;q13) | ETV6 | 12 p 13.2 |  | BAZ2A | 12q13.3 |  |
| $\mathrm{t}(12 ; 13)(\mathrm{p} 13 ; q 12)$ | ETV6 | 12 p 13.2 |  | CDX2 | $13 q 12.2$ |  |
| t(3;12)(q26;p13) | ETV6 | 12 p 13.2 |  | EVI1 | 3q26.2 |  |
| t(4;12)(p16;p13) | ETV6 | 12 p 13.2 |  | FGFR3 | 4p16.3 |  |
| t(12;13)(p13;q12) | ETV6 | 12 p 13.2 |  | FLT3 | 13q12.2 |  |
| t(6;12)(q22;p13) | ETV6 | 12p13.2 |  | FRK | 6q22.1 |  |
| t(10;12)(q24;p13) | ETV6 | 12 p 13.2 |  | GOT1 | 10q24.2 | FRA10A (rare, folic acid) |
| t(9;12)(p24;p13) | ETV6 | 12p13.2 |  | JAK2 | 9p24.1 |  |
| t(3;12)(q26;p13) | ETV6 | 12 p 13.2 |  | MDS1 | 3q26.2 |  |
| t(1;12)(p36;p13) | ETV6 | 12 p 13.2 |  | MDS2 | 1p36.11 | FRA1A (common, apc) |
| t(12;15)(p13;q25) | ETV6 | 12p13.2 |  | NTRK3 | 15q25.3 |  |
| t(4;12)(q12;p13) | ETV6 | 12p13.2 |  | PDGFRA | 4q12 | FRA4B (common, BrdU) |


| t(5;12)(q33;p13) | ETV6 | 12p13.2 |  | PDGFRB | 5 q 3.1 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| t(12;17)(p13;p13) | ETV6 | 12p13.2 |  | PER1 | 17p13.1 |  |
| inv(12)(p13q15) | ETV6 | 12 p 13.2 |  | PTPRR | $12 q 15$ |  |
| t(12;21)(p13;q22) | ETV6 | 12 p 13.2 |  | RUNX1 | 21q22.12 |  |
| t(6;12)(q23;p13) | ETV6 | 12p13.2 |  | STL | 6 q 23 |  |
| t(9;12)(q22;p13) | ETV6 | 12 p 13.2 |  | SYK | 9 q 22.2 |  |
| t(12;22)(q13;q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | ATF1 | 12q13.13 | FRA12A (rare, folic acid) |
| t(2;22)(q33;q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | CREB1 | 2q33.3 | FRA21 (common,apc) |
| t(12;22)(q13;q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | DDIT3 | 12 q 13.3 |  |
| t(21;22)(q22;q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | ERG | 21q22.2 |  |
| t(7;22)(p21;q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | ETV1 | 7p21.2 |  |
| t(17;22)(q21;q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | ETV4 | 17 q 21.31 |  |
| t(2;22)(q35;q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | FEV | 2 q 35 |  |
| t(11;22)(q24;q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | FLII | 11q24.3 |  |
| t(9;22)(q31;q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | NR4A3 | 9q31.1 |  |
| inv(22)(q12q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | PATZ1 | 22q12.2 | FRA22B (common, apc) |
| t(6;22)(p21;q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | POU5F1 | 6p21.33 | FRA6H (common, apc) |
| ( $2 ; 22$ )(q31;q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | SP3 | 2q31.1 | FRA2G (common, apc) |
| t(11;22)(p13;q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | WT1 | 11 p 13 | FRA11E (common, apc) |
| t(12;22)(p13;q12) | EWSR1 | 22q12.2 | FRA22B (common, apc) | ZNF384 | 12 p 13.31 |  |
| t(5;7)(q31;q34) | FCHSD1 | 5 q 31.3 |  | BRAF | 7 7 4 |  |
| t(6;8)(q27;p12) | FGFR1OP | 6 q 27 |  | FGFR1 | 8p12 |  |
| del(4)(q12q12)* | FIP1L1 | 4q12 | FRA4B (common, BrdU) | PDGFRA | 4q12 | FRA4B (common, BrdU) |
| t(4;17)(q12;q21) | FIP1L1 | 4q12 | FRA4B (common, BrdU) | RARA | 17 q 21.2 |  |
| t(2;13)(q36;q14) | FOXO1A | 13q14.11 |  | PAX3 | 2q36.1 |  |
| t(X;11)(q13;q23) | FOXO4 | Xq13.1 |  | MLL | 11 q 23.3 | FRA11B (rare, folic acid), FRA11G (common, |
| t(12;16)(q13;p11) | FUS | 16p11.2 |  | ATF1 | 12q13.13 | FRA12A (rare, folic acid) |
| t(11;16)(p11;p11) | FUS | 16p11.2 |  | CREB3L1 | 11 p 11.2 |  |
| t(7;16)(q34;p11) | FUS | 16p11.2 |  | CREB3L2 | 7q33-q34 |  |
| t(12;16)(q13;p11) | FUS | 16p11.2 |  | DDIT3 | 12q13.3 |  |
| t(16;21)(p11;q22) | FUS | 16p11.2 |  | ERG | 21922.2 |  |
| t(2;16)(q35;p11) | FUS | 16p11.2 |  | FEV | 2 q 35 |  |
| t(3;12)(q27;p13) | GAPDH | 12 p 13.31 |  | BCL6 | 3q27.3 | FRA3C (common, apc) |
| t(5;12)(q33;q24) | GIT2 | 12 q 24.11 | FRA12E (common, apc) | PDGFRB | 5 q 3.1 |  |
| t(10;14)(q11;q32) | GOLGA5 | 14 q 32.12 |  | RET | 10q11.21 | FRA10G (common, apc) |
| del(6)(q21q22)* | GOPC | 6 q 22.2 |  | ROS1 | 6 q 22.2 |  |
| del(8)(q12q24)* | HAS2 | 8q24.13 | FRA8C (common, apc), FRA8E (rare, dist A) | PLAG1 | 8q12.1 |  |
| t(8;19)(p12;q13) | HERV-K (LOC113386) | 19q13.43 | FRA19A (common, 5-aza) | FGFR1 | 8p12 |  |
| t(5;7)(q33;q11) | HIP1 | 7 q 11.23 | FRA7J (common, apc) | PDGFRB | 5 q 33.1 |  |
| t(3;6)(q27;p22) | HIST1H4I | 6p22.1 |  | BCL6 | 3q27.3 | FRA3C (common, apc) |


| inv(6)(p21q21) | HMGA1 | 6p21.31 | FRA6H (common, apc) | LAMA4 | 6 q 21 | FRA6F (common, apc) |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| t(12;14)(q14;q11) | HMGA2 | 12q14.3 |  | CCNB1IP1 | $14 \mathrm{q11.2}$ |  |
| t(8;12)(q22;q14) | HMGA2 | 12q14.3 |  | COX6C | 8 q 22.2 |  |
| t(2;12)(q37;q14) | HMGA2 | 12q14.3 |  | CXCR7 | 2 q 37.3 | FRA2J (common, apc) |
| t(5;12)(q33;q14) | HMGA2 | 12q14.3 |  | EBF1 | 5 q 3.3 |  |
| t(12;13)(q14;q13) | HMGA2 | 12q14.3 |  | LHFP | $13 q 13.3$ |  |
| t(3;12)(q28;q14) | HMGA2 | 12q14.3 |  | LPP | 3 q 28 |  |
| t(9;12)(p23;q14) | HMGA2 | 12q14.3 |  | NFIB | 9p23-p22.3 |  |
| $\mathrm{t}(12 ; 14)(\mathrm{q} 14 ; \mathrm{q} 24)$ | HMGA2 | 12q14.3 |  | RAD51L1 | 14q24.1 | FRA14C (common, apc) |
| t(7;7)(p15;p21) | HNRPA2B1 | 7p15.2 |  | ETV1 | 7p21.2 |  |
| t(8;10)(p11;q11) | HOOK3 | 8p11.21 |  | RET | 10q11.21 | FRA10G (common, apc) |
| t(6;16)(p21;q22) | HP | 16 q 22.3 |  | MRPS10 | 6p21.1 | FRA6H (common, apc) |
| t(3;14)(q27;q32) | HSP90AA1 | 14q32.31 |  | BCL6 | 3q27.3 | FRA3C (common, apc) |
| t(3;6)(q27;p21) | HSP90AB1 | 6 p 21.1 | FRA6H (common, apc) | BCL6 | 3q27.3 | FRA3C (common, apc) |
| t(1;14)(p22;q32) | IGH@ | 14q32.33 |  | BCL10 | 1 p 22.3 | FRA1D (common, apc) |
| t(2;14)(p16;q32) | IGH@ | 14q32.33 |  | BCL11A | 2p16.1 |  |
| t(14;19)(q32;q13) | IGH@ | 14q32.33 |  | BCL3 | 19q13.31 | FRA19A (common, 5-aza) |
| t(3;14)(q27;q32) | IGH@ | 14q32.33 |  | BCL6 | 3q27.3 | FRA3C (common, apc) |
| $\mathrm{t}(14 ; 15)(\mathrm{q} 32 ; \mathrm{q} 11-13)$ | IGH@ | 14q32.33 |  | BCL8 | $15 q 11.2$ |  |
| t(11;14)(q13;q32) | IGH@ | 14q32.33 |  | CCND1 | $11 q 13.2$ | FRA11A (rare, folic acid), FRA11H (common, |
| t(12;14)(p13;q32) | IGH@ | 14q32.33 |  | CCND2 | 12p13.32 |  |
| t(6;14)(p21;q32) | IGH@ | 14q32.33 |  | CCND3 | 6p21.1 | FRA6H (common, apc) |
| t(7;14)(q21;q32) | IGH@ | 14q32.33 |  | CDK6 | 7 q 21.2 | FRA7E (common, apc) |
| t(14;19)(q32;q13) | IGH@ | 14q32.33 |  | CEBPA | 19q13.11 | FRA19A (common, 5-aza) |
| $\mathrm{t}(14 ; 20)(\mathrm{q} 32 ; \mathrm{q} 13)$ | IGH@ | 14q32.33 |  | CEBPB | 20q13.13 |  |
| t(8;14)(q11;q32) | IGH@ | 14q32.33 |  | CEBPD | 8 q 11.21 |  |
| t(14;14)(q11;q32) | IGH@ | 14q32.33 |  | CEBPE | 14 q 11.2 |  |
| t(14;19)(q32;q13) | IGH@ | 14q32.33 |  | CEBPG | 19q13.11 | FRA19A (common, 5-aza) |
| t(12;14)(q23;q32) | IGH@ | 14q32.33 |  | CHST11 | 12q23.3 |  |
| $\mathrm{t}(11 ; 14)(\mathrm{q} 23 ; \mathrm{q} 32)$ | IGH@ | 14q32.33 |  | DDX6 | 11 q 23.3 | FRA11B (rare, folic acid), FRA11G (common, |
| t(7;14)(q21;q32) | IGH@ | 14q32.33 |  | ERVWE1 | 7q21.2 | FRA7E (common, apc) |
| $\mathrm{t}(12 ; 14)(\mathrm{p} 13 ; \mathrm{q} 32)$ | IGH@ | 14q32.33 |  | ETV6 | 12p13.2 |  |
| t(1;14)(q23;q32) | IGH@ | 14q32.33 |  | FCGR2B | 1q23.3 |  |
| t(1;14)(q21;q32) | IGH@ | 14q32.33 |  | FCRL4 | 1q23.1 |  |
| t(4;14)(p16;q32) | IGH@ | 14q32.33 |  | FGFR3 | 4p16.3 |  |
| t(3;14)(p14;q32) | IGH@ | 14q32.33 |  | FOXP1 | 3p14.1 |  |
| t(6;14)(p22;q32) | IGH@ | 14q32.33 |  | ID4 | 6p22.3 |  |
| t(14;22)(q32;q11) | IGH@ | 14q32.33 |  | IGL@ | 22q11.22 |  |
| t(5;14)(q31;q32) | IGH@ | 14q32.33 |  | IL3 | 5q31.1 | FRA5C (common, apc) |


| t(6;14)(p25;q32) | IGH@ | $14 q 32.33$ |  | IRF4 | 6 p 25.3 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| t(1;14)(p35;q32) | IGH@ | 14q32.33 |  | LAPTM5 | 1p35.2 |  |
| t(1;14)(q25;q32) | IGH@ | 14q32.33 |  | LHX4 | 1 q 25.2 |  |
| t(14;16)(q32;q23) | IGH@ | 14q32.33 |  | MAF | 16 q 23.1 |  |
| t(14;20)(q32;q12) | IGH@ | 14q32.33 |  | MAFB | 20 q 12 |  |
| t(14;18)(q32;q21) | IGH@ | 14q32.33 |  | MALT1 | 18 q 21.32 | FRA18B (common, apc) |
| t(1;14)(q22;q32) | IGH@ | 14q32.33 |  | MUC1 | 1 q 22 |  |
| (8;14)(q24;q32) | IGH@ | 14q32.33 |  | MYC | 8q24.21 |  |
| t(10;14)(q24;q32) | IGH@ | 14q32.33 |  | NFKB2 | 10 q 24.32 |  |
| t(11;14)(q23;q32) | IGH@ | 14q32.33 |  | PAFAH1B2 | 11q23.3 | FRA11B (rare, folic acid), FRA11G (common, |
| t(9;14)(p13;q32) | IGH@ | 14q32.33 |  | PAX5 | 9 p 13.2 |  |
| t(11;14)(q23;q32) | IGH@ | 14q32.33 |  | PCSK7 | 11q23.3 | FRA11B (rare, folic acid), FRA11G (common, |
| t(4;14)(p14;q32) | IGH@ | 14q32.33 |  | RHOH | 4p14 |  |
| t(14;19)(q32;q13) | IGH@ | 14q32.33 |  | SPIB | 19q13.33 | FRA19A (common, 5-aza) |
| t(14;14)(q11;q32), inv(14)(q11q32) | IGH@ | 14q32.33 |  | TRA@ | 14q11.2 |  |
| inv(14)(q11q32) | IGH@ | 14q32.33 |  | TRD@ | 14q11.2 |  |
| t(4;14)(p16;q32) | IGH@ | 14q32.33 |  | WHSC1 | 4p16.3 |  |
| t(14;16)(q32;q23) | IGH@ | 14q32.33 |  | wwox | 16 q 23.1 | FRA16D (common, apc) |
| t(1;2)(p22;p11) | IGK@ | 2p11.2 | FRA2L (rare, folic acid) | BCL10 | 1 p 22.3 | FRA1D (common, apc) |
| t(2;19)(p11;q13) | IGK@ | 2 p 11.2 | FRA2L (rare, folic acid) | BCL3 | 19q13.31 | FRA19A (common, 5-aza) |
| t (2;3)(p11;q27) | IGK@ | 2p11.2 | FRA2L (rare, folic acid) | BCL6 | 3q27.3 | FRA3C (common, apc) |
| t(2;11)(p11;q13) | IGK@ | 2 p 11.2 | FRA2L (rare, folic acid) | CCND1 | 11q13.2 | FRA11A (rare, folic acid), FRA11H (common, |
| t(2;12)(p11;p13) | IGK@ | 2 p 11.2 | FRA2L (rare, folic acid) | CCND2 | 12p13.32 |  |
| t(2;7)(p11;q21) | IGK@ | 2 p 11.2 | FRA2L (rare, folic acid) | CDK6 | 7 q 21.2 | FRA7E (common, apc) |
| t(2;18)(p11;q21) | IGK@ | 2 p 11.2 | FRA2L (rare, folic acid) | FVT1 | 18 q 21.33 | FRA18B (common, apc) |
| $\mathrm{t}(2 ; 8)(\mathrm{p} 11 ; \mathrm{q} 24)$ | IGK@ | 2 p 11.2 | FRA2L (rare, folic acid) | MYC | 8q24.21 |  |
| t(2;8)(p11;q24) | IGK@ | 2p11.2 | FRA2L (rare, folic acid) | PVT1 | 8q24.21 |  |
| t(2;6)(p11;q25) | IGK@ | 2p11.2 | FRA2L (rare, folic acid) | ZC3H12D | 6 q 25.1 |  |
| t(19;22)(q13;q11) | IGL@ | 22q11.22-q11.23 |  | BCL3 | 19q13.31 | FRA19A (common, 5-aza) |
| t(3;22)(q27;q11) | IGL@ | 22q11.22-q11.23 |  | BCL6 | 3q27.3 | FRA3C (common, apc) |
| t(11;22)(q13;q11) | IGL@ | 22q11.22-q11.23 |  | CCND1 | 11q13.2 | FRA11A (rare, folic acid), FRA11H (common, |
| t(12;22)(p13;q11) | IGL@ | 22q11.22-q11.23 |  | CCND2 | 12p13.32 |  |
| t(6;22)(p21;q11) | IGL@ | 22q11.22-q11.23 |  | CCND3 | 6p21.1 | FRA6H (common, apc) |
| t(7;22)(q21;q11) | IGL@ | 22q11.22-q11.23 |  | CDK6 | 7 q 21.2 | FRA7E (common, apc) |
| t(16;22)(q23;q11) | IGL@ | 22q11.22-q11.23 |  | MAF | 16 q 23.1 |  |
| t(8;22)(q24;q11) | IGL@ | 22q11.22-q11.23 |  | MYC | 8q24.21 |  |
| t(8;22)(q24;q11) | IGL@ | 22q11.22-q11.23 |  | PVT1 | 8q24.21 |  |
| t(2;22)(p16;q11) | IGL@ | 22q11.22-q11.23 |  | REL | 2p16.1 |  |
| t(16;22)(q23;q11) | IGL@ | 22q11.22-q11.23 |  | wwox | 16 q 23.1 | FRA16D (common, apc) |
| t(4;16)(q27;p13) | IL2 | 4 q 27 |  | DEXI | 16p13.13 |  |


| t(4;16)(q27;p13) | IL2 | 4 q 27 |  | TNFRSF17 | 16p13.13 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (3;16)(q27;p12) | IL21R | 16 p 12.1 | FRA16E (rare, dist A) | BCL6 | 3 q 27.3 | FRA3C (common, apc) |
| t(5;9)(q33;q22) | ITK | 5 q 33.3 |  | SYK | 9q22.2 |  |
| t(6;7)(p21;p15) | JAZF1 | 7p15.2-p15.1 |  | PHF1 | 6p21.32 | FRA6H (common, apc) |
| (7; ;17)(p15;q11) | JAZF1 | 7p15.2-p15.1 |  | SUZ12 | $17 \mathrm{q11.2}$ |  |
| ( $2 ; 17$ )(p23;q25) | KIAA1618 | 17 q 25.3 |  | ALK | 2p23.2-p23.1 |  |
|  | KIF5B | 10p11.22 |  | PDGFRA | 4 q 12 | FRA4B (common, BrdU) |
| t(10;14)(q11;q22) | KTN1 | 14q22.3 |  | RET | 10q11.21 | FRA10G (common, apc) |
| t(12; 16)(p13;p13) | LAG3 | $12 \mathrm{p13.31}$ |  | MYH11 | 16 p 13.11 | FRA16A (rare, folic acid) |
| t(1;7)(p35;q34) | LCK | 1 p35.1 |  | TRB@ | 7934 |  |
| t(3;13)(q27;q14) | LCP1 | $13 \mathrm{q14.12}$ |  | BCL6 | 3 q 27.3 | FRA3C (common, apc) |
| t(5;8)(p13;q12) | LIFR | 5 p 13.1 | FRA5A (common, BrdU) | PLAG1 | 8912.1 |  |
| del(3)(q27q28)* | LPP | 3928 |  | BCL6 | 3 q 27.3 | FRA3C (common, apc) |
| t(7;19)(q34;p13) | LYL1 | 19p13.13 | FRA19B (rare, folic acid) | TRB@ | 7934 |  |
| t(11;19)(q13;q13.4) | MALAT1 | 11913.1 | FRA11H (common, apc) | MHLB1 | 19q13.4 | FRA19A (common, 5-aza) |
| t(6;11)(p21.1;q13) | MALAT1 | 11q13.1 | FRA11H (common, apc) | TFEB | 6 p 21.1 | FRA6H (common, apc) |
| t(3;18)(p21;q21) | MALT1 | 18921.32 | FRA18B (common, apc) | MAP4 | 3p21.31 |  |

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APPENDIX E
[TOOL SCRIPTS]

## Extraction algorithm

\#!/usr/bin/perl -w
use strict;
use warnings;
my \$infile_name=shift(@ARGV);
my \$outfile1_name=\$infile_name;
my \$outfile2_name=\$infile_name;
\# check file extension is .sam or .bam
if(\$infile_name=~/.(s|b)am\$/)\{
\$outfile1_name=~s/.(b|s)am\$/.extraction2.sam/; \# add .extraction.sam suffix to output file \$outfile2_name=~s/.(b|s)am\$/.extraction2.fa/; \# add .extraction.fq suffix to output file \}else\{
die "Error: input file needs to be in bam/sam format. \n"; \# if file extension not .sam and .bam file then exit
\}
my \$isbam=(\$infile_name =~ /.bam\$//? 1:0; \# if .bam file then return 1(true)

```
if($isbam){
    open(INFILEHDL, "samtools view $infile_name |") or die "$0: can't open ".$infile_name.":$!\n"; #
open bam file
}else{
    open INFILEHDL, "<".$infile_name or die "$0: can't open ".$$nfile_name.":$!\n"; #open sam file
}
open(OUTFILEHDL1, ">$outfile1_name") or die "$0: can't write in the output: $outfile1_name :$!\n";
# open file in overwite mode
open(OUTFILEHDL2, ">$outfile2_name") or die "$0: can't write in the output: $outfile2_name :$!\n";
# open file in overwite mode
my $readsum = 0;
my $readcount = 0;
my $prevline = "";
my $prevfld0 = "";
my $prevfld1 = "";
my $prevfld5 = "";
my $prevfld9 = "";
my $currentfld0 = "";
my $currentfld1 = "";
my $currentfld5 = "";
my $currentfld8 = "";
my $currentfld9 = "";
my $first = 0;
my $pattlen =0;
my $lpattern="";
my @meanlist =('99','163','147','83');
my @matchlist =
('73','133','89','121','165','181','101','117','153','185','69','137','77','141','67','131','115','179','81','161'
,'97','145','65','129','113','177');
while (my $LINE=<INFILEHDL>) # read line till EOF
```

\{
\# chomp my \$LINE; \# removes trailing whitespace
my @L=split(/\t+/,\$LINE); \# split on white space, + will merge multiple whitespace \#my @L=split;
if (\$L[0]=~/^@/) \# current line first character is @ then print that line and skip that line
\{
print OUTFILEHDL1 \$LINE;
next;
\}
\$currentfld0 = "";
\$currentfld1 = "";
\$currentfld5 = "";
\$currentfld8 = "";
\$currentfld9 = "";
if ( scalar(@L)>= 11) \# array length is atleast 11
\{
\$currentfld0 = \$L[0]; \#Identifier
\$currentfld1 = \$L[1]; \#FLAG
\$currentfld5 = \$L[5]; \#CIGAR
\$currentfld8 = \$L[8]; \#TLEN
\$currentfld9 = \$L[9]; \#SEQUENCE
if (\$first ==0) \# build search pattern only once
\{
\$pattlen = int(length(\$currentfld9)*.30); \#30\% or less based on decimal
point of identifier length
\$lpattern = "N" x \$pattlen;
\$first = 1;
print \$lpattern."\n"; \# for testing purpose
\}
if (grep\{\$currentfld1 eq \$_\} @meanlist)
\{
\$readsum = \$readsum + \$currentfld8;
\$readcount = \$readcount + 1 ;
\}
if ( (\$currentfld0 eq \$prevfld0) and ( (grep\{\$currentfld1 eq \$_\} @matchlist) || (grep\{\$prevfld1 eq
\$_\} @matchlist) || (\$currentfld5=~/S/ || \$prevfld5=~/S/)) )
\{
print OUTFILEHDL1 \$prevline."\n";
print OUTFILEHDL1 \$LINE."\n";
if (index(\$prevfld9,\$lpattern)==-1 and index(\$currentfld9,\$lpattern)==-1 )
\# if $N$ Pattern not found in Read $1 \&$ Read 2 then write to FA file
\{
print OUTFILEHDL2 ">".\$prevfld0."<RID>1"."\n";
print OUTFILEHDL2 \$prevfld9."\n";
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```
                                    print OUTFILEHDL2 ">".$currentfld0."<RID>2"."\n";
                                    print OUTFILEHDL2 $currentfld9."\n";
                                }
        }
    } # if scalar
    $prevline = $LINE ;
    $prevfld0 = $currentfld0;
    $prevfld1 = $currentfld1;
    $prevfld9 = $currentfld9;
} # while
close OUTFILEHDL1;
close OUTFILEHDL2;
close INFILEHDL;
```


## De-duplication algorithm

```
\#!/usr/bin/perl-w
\# Assumptions
\# psl file needs to be sorted by identifier,chromosome,start position
use strict;
use warnings;
my \$infile_name=shift(@ARGV);
my \$outfile1_name=\$infile_name;
if (not defined \$infile_name)
\{ die "Error: .psl filename missing \n";
\}
\# check file extension is psl
if(\$infile_name=~/.(p)sl\$/)\{
\$outfile1_name=~s/.(p)sl\$/.sml/; \# add .sml suffix to output file
\}else\{
die "Error: input file needs to be in psl format. \(\backslash n\) "; \# if file extension not .psl then exit
\}
open(OUTFILEHDL1, ">\$outfile1_name") or die "\$0: can't write in the output: \$outfile1_name :\$! \n";
\# open file in overwite mode
open INFILEHDL, "<".\$infile_name or die "\$0: can't open ".\$infile_name.":\$!\n"; \#open sam file
my @read1;
my @read2;
my \$prevline = "";
my \$prevfld9 = "";
my \$prevfld13 = "";
my \$prevfld15 = "";
my \$prevfld16 = "";
my \$prevfld17 = "";
my \$prevrange = 0;
```

```
my $currentfld9 = "";
my $currentfld13 = "";
my $currentfld15 = "";
my $currentfld16 = "";
my $currentfld17 = "";
my $currrange = 0;
```

my \$iline =1;
my \$processline = 0;
my $\$ \mathrm{i}=0$;
my $\$ \mathrm{j}=0$;
my \$lastmaxrange = 0;
my @cid;
my @pid;
while (my \$LINE=<INFILEHDL>) \# read line till EOF
\{
my @L=split(/\t/,\$LINE); \# split on tab space
\$processline +=1;
if (scalar(@L)>= 18) \# array length is at least 17
\{
\$currentfld9 = \$L[9]; \# Identifier
if(substr(\$L[13],0,3) eq "chr")
\{
\$currentfld13 = substr(\$L[13],3); \# Chromosome
\}
else
\{
\$currentfld13 = \$L[13]; \#Chromosome
\}
\$currentfld15 = \$L[15]; \# Start
\$currentfld16 = \$L[16]; \# End
\$currentfld17 = \$L[17]; \# Blockcount
\#\$currrange = GetRange\$currentfld15,\$currentfld16();
\$currrange = GetRange(\$currentfld15,\$currentfld16);
@cid = split("<RID>",\$currentfld9);
@pid = split("<RID>",\$prevfld9);
if ( (scalar(@cid) >1) )
\{
if (\$currentfld9 ne \$prevfld9) \# identifier not match then save
\{
\$lastmaxrange = 0; \# if identifier changes then reset last
max range

```
    if ($currentfld17 eq "1") # blockcount = 1 then add to array
    {
        InsertRead();
```

    \}
    \}
    ```
                                    else # if same identifier then check chr range and blockcount \{
    if ( $currentfld13 ne $prevfld13 ) # chromsome diff then save
    {
        $lastmaxrange = 0; # if chromosome changes then reset
last max range
    if ($currentfld17 eq "1") # blockcount = 1 then add to
array
    {
        InsertRead();
        }
        else # if chromosome same then check range
        {
        if ( $currrange != $prevrange) # if range diff then
save
        {
            if ($currentfld17 eq "1") # blockcount =
1 then add to array
                    {
                                    InsertRead();
                                    }
        }
        }
                            }
            }
            }
    } #scalar L
} # eof
close OUTFILEHDL1;
close INFILEHDL;
sub GetRange
{
    my $rangespan = 1000;
    my $grange = 0;
    my $st = 0;
    my $ed = 0;
    $st = $_[0];
    $ed = $_[1];
    my $midpoint = int(($ed+$st)/2);
    if ($midpoint < $rangespan)
    {
        $grange = 1;
    }
    else
    {
        $grange = int($midpoint/$rangespan) + 1;
            1 6 7
```

```
    }
    return $grange ;
}
sub InsertRead
{
    if ($currrange >$lastmaxrange)
    {
        print OUTFILEHDL1
$currentfld9."\t".$currentfld13."\t".$currentfld15."\t".$currentfld16."\t".$currrange."\n" ;
            $lastmaxrange = $currrange ;
            $prevline = my $LINE;
            $prevfld9 = $currentfld9;
            $prevfld13 = $currentfld13;
            $prevfld15 = $currentfld15;
            $prevfld16 = $currentfld16;
            $prevfld17 = $currentfld17;
            $prevrange = $currrange;
    }
}
```


## Create_matrix algorithm

```
#!/usr/bin/perl -w
```

```
#!/usr/bin/perl -w
```

```
# Assumption
```


# Assumption

# sml file is ordered by Idenitifier,chromosome

# sml file is ordered by Idenitifier,chromosome

use strict;
use strict;
use warnings;
use warnings;
my \$infile_name=shift(@ARGV);
my \$infile_name=shift(@ARGV);
my $outfile1_name=$infile_name;
my $outfile1_name=$infile_name;
if (not defined \$infile_name)
if (not defined \$infile_name)
{
{
die "Error: .sml filename missing \n";
die "Error: .sml filename missing \n";
}
}

# check file extension is sml

# check file extension is sml

if($infile_name=~/.(s)ml$/){
if($infile_name=~/.(s)ml$/){
$outfile1_name=~s/.(s)ml$/.unsort/; \# add .unsort suffix to output file
$outfile1_name=~s/.(s)ml$/.unsort/; \# add .unsort suffix to output file
}else{
}else{
die "Error: input file needs to be in sml format.\n"; \# if file extension not .sml then exit
die "Error: input file needs to be in sml format.\n"; \# if file extension not .sml then exit
}
}
open(OUTFILEHDL1, ">\$outfile1_name") or die "\$0: can't write in the output: $outfile1_name :$!\n";
open(OUTFILEHDL1, ">\$outfile1_name") or die "\$0: can't write in the output: $outfile1_name :$!\n";

# open file in overwite mode

# open file in overwite mode

open INFILEHDL, "<".\$infile_name or die "$0: can't open ".$infile_name.":$!\n"; #open sml file
open INFILEHDL, "<".$infile_name or die "$0: can't open ".$infile_name.":\$!\n"; \#open sml file
my @read1;
my @read1;
my @read2;
my @read2;
my \$prevline = "";

```
my $prevline = "";
```

```
my $prevfld1 = "";
my $prevfld2 = "";
my $prevfld3 = "";
my $prevfld4 = "";
my $prevfld5 = "";
my $currentfld1 = "";
my $currentfld2 = "";
my $currentfld3 = "";
my $currentfld4 = "";
my $currentfld5 = "";
my $matfilename = $outfile1_name;
my $processline = 0;
my $i=0;
my $j=0;
my @cid;
my @pid;
$matfilename = ~s/.(u)nsort$/.mat/;
while (my $LINE=<INFILEHDL>) # read line till EOF
{
    chomp($LINE);
    my @L=split(/\t/,$LINE); # split on tab space
    $processline += 1;
    if (scalar(@L)>= 5) # array length is at least 5
    {
```

```
$currentfld1 = $L[0]; # Identifier
```

\$currentfld1 = \$L[0]; \# Identifier
\$currentfld2 = \$L[1]; \# Chromosome
\$currentfld2 = \$L[1]; \# Chromosome
\$currentfld3 = \$L[2]; \# Start
\$currentfld3 = \$L[2]; \# Start
\$currentfld4 = \$L[3]; \# End
\$currentfld4 = \$L[3]; \# End
\$currentfld5 = \$L[4]; \# range
\$currentfld5 = $L[4]; # range
@cid = split("<RID>",$currentfld1);
@cid = split("<RID>",$currentfld1);
@pid = split("<RID>",$prevfld1);
@pid = split("<RID>",$prevfld1);
if ( (scalar(@cid) >1) and (scalar(@pid) >1) )
if ( (scalar(@cid) >1) and (scalar(@pid) >1) )
{
{
    if ($pid[0] eq $cid[0]) # identifier match
    if ($pid[0] eq \$cid[0]) \# identifier match
{
{
InsertRead();
InsertRead();
}
}
else \# if different match then save array to file
else \# if different match then save array to file
{
{
CreateMatrix();
CreateMatrix();
@read1 =(); \# \$\#read1 = -1 \# clear array
@read1 =(); \# $#read1 = -1 # clear array
        @read2 =(); # clear array
        @read2 =(); # clear array
            InsertRead();
            InsertRead();
}
}
}
}
if (($processline == 1) and (scalar(@cid) >1) ) \# save first line to array
if ((\$processline == 1) and (scalar(@cid) >1) ) \# save first line to array
{
{
InsertRead();

```
    InsertRead();
```

```
            }
    } #scalar L
    $prevline = $LINE;
    $prevfld1 = $currentfld1;
    $prevfld2 = $currentfld2;
    $prevfld3 = $currentfld3;
    $prevfld4 = $currentfld4;
    $prevfld5 = $currentfld5;
} # eof
CreateMatrix(); #save array to file remaining ones
close OUTFILEHDL1;
close INFILEHDL;
# do unix sort for field 1,Field3,Field2,Field4 here on the mat file
system ("sort -k1n,1 -k3n,3 -k2n,2 -k4n,4 $outfile1_name > $matfilename");
sub CreateMatrix
{
    my @ln1;
    my @ln2;
    my $cnt = 1;
    foreach my $r1(@read1) # r1 and r2 combination
    {
        foreach my $r2(@read2)
        {
        @ln1 = split(/\t/,$r1);
            @ln2 = split(/\t/,$r2);
            if ($\operatorname{ln}1[0] eq }$\operatorname{ln}2[0] and $ ln1[1] eq $ ln2[1]) # if chromosome and offset i
same then do not print
            {
                # do not print
            }
            else
            {
                if (PivotFile($ln1[0],$ln2[0],$ln1[1],$ln2[1]) == 2)
    {
        print OUTFILEHDL1 $r1."\t".$r2."\n";
    }
        else
    {
        print OUTFILEHDL1 $r2."\t".$r1."\n";
    }
            }
    }
    }
    for(my $x=0; $x < scalar(@read1); $x++) # r1 unique combination
```

    {
            for(my $j=$cnt; $j < scalar(@read1); $j++)
            {
        @ln1 = split(/\t/,$read1[$x]);
        @ln2 = split(/\t/,$read1[$j]);
        if ($\operatorname{ln}1[0] eq }$\operatorname{ln}2[0] and $ ln1[1] eq $ ln2[1]) # if chromosome and offset is same then d
    not print
{
\# do not print
}
else
{
if (PivotFile($ln1[0],$ln2[0],$\operatorname{ln}1[1],$\ln2[1])== 2)
print OUTFILEHDL1 $read1[$x]."\t".$read1[$j]."\n";
}
else
{
print OUTFILEHDL1 $read1[$j]."\t".$read1[$x]."\n";
}
}
}
\$cnt++;
}
\$cnt = 1;
for(my \$x=0; \$x < scalar(@read2); \$x++) \# r2 unique combination
{
for(my $j=$cnt; \$j < scalar(@read2); $j++)
    {
        @ln1 = split(/\t/,$read2[$x]);
        @ln2 = split(/\t/,$read2[$j]);
        if ($\operatorname{ln}1[0] eq }\$\operatorname{ln}2[0] and \$ ln1[1] eq \$ ln2[1]) \# if chromosome and offset is same then do
not print
{
\# do not print
}
else
{
if (PivotFile($ln1[0],$ln2[0],$ln1[1],$ln2[1]) == 2)
{
print OUTFILEHDL1 $read2[$x]."\t".$read2[$j]."\n";
}
else
{
print OUTFILEHDL1 $read2[$j]."\t".$read2[$x]."\n";
}
}
}
\$cnt++;
}

```
\}
sub InsertRead
\{
        if(\$cid[1] eq "1") \# read 1
        \{
            push @read1, \$currentfld2."\t".\$currentfld5;
        \}
        if(\$cid[1] eq "2") \# read 2
        \{
            push @read2, \$currentfld2."\t".\$currentfld5;
        \}
\}
sub PivotFile
\{
    my \$arg1 ="";
    my \$arg2 ="";
        my \$arg3 ="";
        my \$arg4 ="";
    \$arg1 = \$_[0]; \# chr1
    \$arg2 = \$_[1]; \# chr2
        \$arg3 = \$_[2]; \# offset1
        \$arg4 = \$_[3]; \# offset2
    my \$rw;
    my \$co;
    if (uc(\$arg1) eq 'X' or uc(\$arg1) eq 'Y')
    \{
        \(\$ r w=0 ;\)
    \}
    else
    \{
        \$rw = int(\$arg1);
    \}
    if (uc(\$arg2) eq 'X' or uc(\$arg2) eq 'Y')
    \{
        \$co=0;
    \}
    else
    \{
        \$co = int(\$arg2);
    \}
    my \(\$\) pvot \(=2\); \# default 2 = no 1 = yes
    if ( \(\$ \mathrm{rw}==0\) and \(\$ \mathrm{co}==0\) ) \# both row and col are x or y
    \{
        if (\$arg1 gt \$arg2) \# row greater than col then pivot
        \{
            \$pvot = 1;
        \}
    \}
```

    else
    {
    if ($rw == 0 and $co > 0) # row is x or y, col is number then pivot
    {
        $pvot = 1;
    }
    }
if (\$rw > 0 and $co > 0) # row and col both are numbers
{
    if ($rw > \$co) \# row is greater than col then pivot
{
$pvot = 1;
    }
}
    if ($arg1 eq $arg2) # if and row and col is same
    {
            if (int($arg3) > int(\$arg4)) \# if row offset is greater than col offset
{
\$pvot = 1;
}
}
return \$pvot;
}

```

\section*{Write_count algorithm}
```

\#!/usr/bin/perl -w

# Assumption

# mat file is ordered by field1,field3,field2,field4

use strict;
use warnings;
my \$infile_name=shift(@ARGV);
my $outfile1_name=$infile_name;
if (not defined \$infile_name)
{
die "Error: .sml filename missing \n";
}

# check file extension is mat

if($infile_name=~/.(m)at$/){
$outfile1_name=~s/.(m)at$/.chr/; \# add .chr suffix to output file
}else{
die "Error: input file needs to be in mat format.\n"; \# if file extension not .mat then exit
}

```
```

open(OUTFILEHDL1, ">\$outfile1_name") or die "\$0: can't write in the output: $outfile1_name :$!\n";

# open file in overwite mode

open INFILEHDL, "<".\$infile_name or die "$0: can't open ".$infile_name.":\$!\n"; \#open mat file
my \$prevline = "";
my \$iline = 1;
my \$cnt = 0;
my \$first = 1;
while (my $LINE=<INFILEHDL>) # read line till EOF
{
    chomp($LINE);
if (\$first == 1)
{
\$prevline = \$LINE;
$first = 0;
    }
    if ($prevline eq \$LINE)
{
\$cnt = \$cnt + 1;
}
else
{
WriteCount();
\$cnt = 1;
}
\$prevline = \$LINE;
} \# eof
WriteCount(); \#save remaining ones
close OUTFILEHDL1;
close INFILEHDL;
sub WriteCount
{
print OUTFILEHDL1 $prevline."\t".$cnt."\n";
}

```

\section*{Get Hi-C Score algorithm}
\#!/usr/bin/perl -w
use strict;
use warnings;
\# Assumptions chr file is sorted by field1,field3,field2,field4
\# Hi -c file range is 1000000 and they are tab delimited
my \$infile_name = shift(@ARGV);
my \$hicfile_path = shift(@ARGV);
my \$outfile1_name = \$infile_name;
\# check file extension is chr
if(\$infile_name=~/.(c)hr\$//)\{
\$outfile1_name=~s/.(c)hr\$/.score/; \# add .score suffix to output file \}else\{
die "Error: input file needs to be in chr format.\n"; \# if file extension not .chr then exit \}
if (not defined \$hicfile_path )
\{ die "Error: HIC file path not found. \(\ n\) ";
\}
open(OUTFILEHDL1, ">\$outfile1_name") or die "\$0: can't write in the output: \$outfile1_name :\$!\n"; \# open file in overwite mode open INFILEHDL, "<".\$infile_name or die "\$0: can't open ".\$infile_name.":\$!\n"; \#open MAT file my \$currR1fld9 = ""; my \$currR1fld13 = ""; my \$currR1fld15 = ""; my \$currR1fld16 = ""; my \$currR2fld9 = ""; my \$currR2fld13 = ""; my \$currR2fld15 = ""; my \$currR2fld16 = " "; my \$prevR1fld13 = ""; my \$prevR2fld13 = "";
my \$currhicfile = "";
my \$prevhicfile = "";
my \$hicfile = "";
my @filearray ;
my \(\$\) colmidrange \(=0\);
my \$rowmidrange = 0;
my \$rowpos = 0;
my \(\$\) colpos \(=0\);
my \(\$\) r1start \(=0\);
my \$r1end =0;
my \$ r 2 start =0;
my \$r2end =0;
my \$currCount="";
my \$score = "";
my \$pivot = 2;
while (my \$LINE=<INFILEHDL>) \# read line till EOF
\{
chomp(\$LINE);
my @L=split(//t/,\$LINE); \# split on tab space
\$currR1fld9 = "";
\$currR1fld13 = "";
\$currR1fld15 = "";
\$currR1fld16 = "";
\$currR2fld9 = "";
\$currR2fld13 = "";
\$currR2fld15 = "";
\$currR2fld16 = "";
```

            $currCount = "";
            $r1start=0;
            $r1end=0;
            $r2start=0;
            $r2end=0;
            $score = "";
            $pivot = 2;
            if (scalar(@L)>= 5) # array length is at least 7
    {
\$currR1fld9 = \$L[0]; \# Chromosome
\$currR1fld13 = \$L[1]; \# Column Offset
\$currR2fld9 = \$L[2]; \# Chromosome
\$currR2fld13 = \$L[3]; \# Column Offset
\$currCount = $L[4]; # Chromosome Count
            if ($pivot == 1) \# if row chr greater than col chr then swap chr position
HIC_gm06690_chr2_chr18_1000000_pearson.txt
{
$currhicfile =
"HIC_gm06690_chr".uc($currR2fld9)."_"."chr".uc(\$currR1fld9)."_1000000_pearson.txt"; \# file name
will chr21_chr22.hic
}
else
{
$currhicfile=
"HIC_gm06690_chr".uc($currR1fld9)."_"."chr".uc($currR2fld9)."_1000000_pearson.txt"; # file name
will chr21_chr22.hic
    }
    if ($currhicfile ne \$prevhicfile) \# if previous file not same as current file open file
{
openHiCFile(); \#open file
}
\$rowpos = 0;
\$colpos = 0;
$rowmidrange = CalcMidPosition($currR1fld13);
$colmidrange = CalcMidPosition($currR2fld13);
if (\$pivot == 1)
{
$rowpos = GetFilePosition($colmidrange);
$colpos = GetFilePosition($rowmidrange);
}
else
{
$rowpos = GetFilePosition($rowmidrange);
$colpos = GetFilePosition($colmidrange);
}

```
```

        $score = GetHiCScore(); # get hic score
        CalcRange($currR1fld13,$currR2fld13);
        if ($score gt 0) # score greater than 0 then print
        {
            if ($currCount > 1) # count is greater than 1 then print
            {
            print OUTFILEHDL1
    "chr".$currR1fld9."\t".$r1start."\t".$r1end."\t"."chr".$currR2fld9."\t".$r2start."\t".$r2end."\t".$curr
Count."\t".$score."\n";
}
}
}
\$prevR1fld13 = \$currR1fld13;
\$prevR2fld13 = \$currR2fld13;
\$prevhicfile = \$currhicfile;
}
close OUTFILEHDL1;
close INFILEHDL;

# open file and load into array

sub openHiCFile
{
\$hicfile = $hicfile_path.$currhicfile;
@filearray = ();
open INFILEHIC, "<".\$hicfile or die "$0: can't open ".$hicfile.":\$!\n"; \#open hi-c file
@filearray = <INFILEHIC>;
shift(@filearray);
close INFILEHIC;
}

# get HI-C Score base on row and col position

sub GetHiCScore
{
my @cline;
my \$scr = "Not found";
if (scalar(@filearray) > $rowpos) # check if that range exists or not
    {
        @cline = split(/\t/,$filearray[$rowpos]);
#print "score col 0 ".$cline[0]." col 1 ".\$cline[1]."\n";
if (scalar(@cline) > \$colpos) \# check if range exist or not
{
\$scr = $cline[$colpos];
}
else
{

```
print "column offset not found
```

".$currhicfile."\t".$rowpos."\t".$colpos."\n";
    }
    }
    else
    {
        print "row offset not found ".$currhicfile."\t".$rowpos."\t".$colpos."\n";
}
return \$scr;
}

# Calculate Position based on range (not used)

sub GetPosition
{
my \$readpos = \$_[0];
my \$start = 0;
my \$end = 999999;
my \$rfactor = 1000000; \# hi-c file range span by million
my \$pos = 0;
for (my \$i = 0; \$i < 1000; $i++)
    {
            if (($readpos >= $start) and ($readpos <= \$end))
{
$pos=$i;
last;
}
\$start += \$rfactor;
\$end += \$rfactor
}
return \$pos;
}

```
sub GetFilePosition
\{
    my \$readpos = \$_[0];
    my \$colspan =1000000; \#hi-c file span by million
    my \$pos = 0;
    if (\$readpos < \$colspan)
    \{
        \$pos =1;
    \}
    else
    \{
            \$pos = int(\$readpos/\$colspan) + 1;
        \}
    return \$pos;
\}
sub CalcMidPosition
\{
```

            my $coloff = $_[0];
            my $chrspan = 1000; # assuming chr file column offset span is 1000, change this if needed
            my $cpos = 0;
            my $mend = ($coloff * $chrspan)-1;
            my $mstart = $mend - ($chrspan-1);
    $cpos = int(($mend+$mstart)/2);
    return $cpos;
    }
sub CalcRange
{
my \$r1col = \$_[0];
my \$r2col = \$_[1];
my \$chrspan = 1000; \# assuming chr file column offset span is 1000, change this if needed
$r1end = ($r1col * \$chrspan)-1;
\$r1start = $r1end-($chrspan-1);
$r2end = ($r2col * \$chrspan)-1;
\$r2start = $r2end-($chrspan-1);
}
sub PivotFile
{
my \$arg1 ="";
my \$arg2 ="";
\$arg1 = \$_[0]; \# row
\$arg2 = \$_[1]; \# col
my \$rw;
my $co;
    if (uc($arg1) eq 'X' or uc(\$arg1) eq 'Y')
{
\$rw=0;
}
else
{
$rw = int($arg1);
}
if (uc($arg2) eq 'X' or uc($arg2) eq 'Y')
{
\$co=0;
}
else
{
$co = int($arg2);
}
my $pvot = 2; # default 2 = no 1 = yes
    if ($rw == 0 and \$co == 0) \# both row and col are x or y

```
\{
if (\$arg1 gt \$arg2) \# row greater than col then pivot \{ \(\$\) pvot \(=1 ;\) \}
\}
else
\{
        if ( \(\$ \mathrm{rw}==0\) and \(\$ \mathrm{co}>0\) ) \# row is x or y , col is number then pivot
        \{
        \$pvot = 1;
        \}
\}
if (\$rw > 0 and \(\$ c o>0\) ) \# row and col both are numbers
\{
        if (\$rw > \$co) \# row is greater than col then pivot
        \{
        \$pvot = 1;
        \}
\}
return \$pvot;
\}```


[^0]:    ${ }^{1}$ Soft clipping: Paired end read mapped to the reference genome where one end mapped globally an d the other mapped partially, may need the unaligned ends of the read 'clipped' to achieve mapping. The read is labeled as ' S ' in its CIGAR string (Section 4.2.1).
    ${ }^{2}$ Contig: Creating a longer sequence of DNA from overlapping smaller subsequences.

[^1]:    ${ }^{3}$ Polymerase: Enzyme present in the cells, used during DNA replication process for synthesizing a new strand of DNA from a copy/template.
    ${ }^{4}$ Primer: Short DNA segment of known sequence which attached to the DNA strand at its end

[^2]:    ${ }^{5}$ PCR: Polychromase chain reaction, a procedure to create multiple copies of the DNA using DNA polymerase

[^3]:    ${ }^{6}$ Ligase: cellular enzyme which catalyzes the formation of bonds between two DNA strands.

[^4]:    ${ }^{7} 1000$ genomes data download:ftp://ftp-trace.ncbi.nih.gov/1000genomes/ftp/data/
    ${ }^{8}$ BWA sourceforge page: http://bio-bwa.sourceforge.net/bwa.shtml
    ${ }^{9}$ SAMTools sourceforge page: http://samtools.sourceforge.net/

[^5]:    ${ }^{10}$ 1000-Genomes structural variation data page:
    ftp://ftp.1000genomes.ebi.ac.uk/vol1/ftp/pilot_data/paper_data_sets/a_map_of_hu man_variation/trio/sv/

[^6]:    ${ }^{11}$ RepeatMasker webpage: http://www.repeatmasker.org

[^7]:    ${ }^{12}$ PICARD tool webpage: http://picard. sourceforge.net/index.shtml, http://picard.sourceforge. net/explain-flags.html

[^8]:    ${ }^{13} \mathrm{Hi}-\mathrm{C}$ database webpage: http://hic.umassmed.edu/welcome/welcome.php

